

A quantitative Lovász criterion for Property B

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Abstract

A well known observation of Lovász is that if a hypergraph is not 2-colorable, then at least one pair of its edges intersect at a single vertex. In this short paper we consider the quantitative version of Lovász's criterion. That is, we ask how many pairs of edges intersecting at a single vertex, should belong to a non 2-colorable n -uniform hypergraph? Our main result is an *exact* answer to this question, which further characterizes all the extremal hypergraphs. The proof combines Bollobás's two families theorem with Pluhar's randomized coloring algorithm.

1 Introduction

A hypergraph $\mathcal{H} = (V, E)$ consists of a vertex set V and a set of edges E where each $X \in E$ is a subset of V . If all edges of \mathcal{H} have size n then \mathcal{H} is called an n -uniform hypergraph, or n -graph for short. A hypergraph is 2-colorable if one can assign each vertex $v \in V$ one of two colors, say *Red/Blue*, so that each $X \in E$ contains vertices of both colors. Miller [6], and later Erdős in various papers, referred to this property as *Property B*, after F. Bernstein [2] who introduced it in 1907. Since deciding if a hypergraph is 2-colorable is *NP*-hard one cannot hope to find a simple characterization of all 2-colorable hypergraphs. Instead, one looks for general sufficient/necessary conditions for having this property. For example, a famous result of Seymour [8] states that if \mathcal{H} is not 2-colorable then $|E| \geq |V|$. Probably the most well studied question of this type asks for the smallest number of edges in an n -graph that is not 2-colorable. The study of this quantity, denote $m(n)$, was popularized by Erdős, see [1] for a comprehensive treatment. Despite much effort by many researchers, even the asymptotic value of $m(n)$ has not been determined yet.

A pair of edges $X, Y \in E(\mathcal{H})$ is *simple* if $|X \cap Y| = 1$. Let $m_2(\mathcal{H})$ denote the number of ordered simple pairs of edges of \mathcal{H} . A well known observation of Lovász [5] states that if \mathcal{H} is not 2-colorable then $m_2(\mathcal{H}) > 0$. Despite its simplicity, this observation underlies the best known bounds for $m(n)$, see [4, 7]. It is natural to ask if one can obtain a quantitative version of Lovász's observation, that is, estimate how small can $m_2(\mathcal{H})$ be in an n -graph not satisfying property *B*? Our main result in this paper states that (somewhat surprisingly), one can give an exact answer to the above extremal question as well as characterize the extremal n -graphs.

Let K_{2n-1}^n denote the complete n -graph on $2n-1$ vertices. It is easy to see that K_{2n-1}^n is not 2-colorable and that $m_2(K_{2n-1}^n) = n \cdot \binom{2n-1}{n}$. We first observe that this simple upper bound is tight.

Proposition 1.1. *If an n -graph is not 2-colorable then $m_2(\mathcal{H}) \geq n \cdot \binom{2n-1}{n}$.*

As with any extremal problem, one would like to know which graphs or hypergraphs are extremal with respect to this problem. For example, Turán's theorem states that among all n -vertex graphs not containing a complete t -vertex subgraph, there is only one graph maximizing the number of edges. In the setting of our

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problem, it is easy to see that K_{2n-1}^n is not the only non 2-colorable n -graph satisfying $m_2(\mathcal{H}) = n \cdot \binom{2n-1}{n}$, since one can take a copy of K_{2n-1}^n and add to it more vertices and edges without increasing the number of simple pairs. Our main result in this paper characterizes the extremal n -graphs, by showing that this is in fact the only way to construct an n -graph meeting the bound of Proposition 1.1.

Theorem 1. *If a non 2-colorable n -graph \mathcal{H} satisfies $m_2(\mathcal{H}) = n \cdot \binom{2n-1}{n}$ then it contains a copy of K_{2n-1}^n .*

While the proof of Proposition 1.1 is implicit in Pluhar's [7] argument for bounding $m(n)$, the proof of Theorem 1 is more intricate, relying on Bollobás's two families theorem [3] as well as on a refined analysis of Pluhar's randomized algorithm for 2-coloring n -graphs.

2 Proof of Proposition 1.1

In this section we describe several preliminary observations regarding a coloring algorithm introduced in [7], and use them to derive Proposition 1.1. The algorithm is the following:

Algorithm $\text{Col}(\mathcal{H}, \pi)$. The input is a hypergraph $\mathcal{H} = (V, E)$ and an ordering $\pi : V \mapsto \{1, \dots, |V|\}$ (that is, π is a bijection). The output is a 2-coloring of V (not necessarily a proper one). The algorithm runs in $|V|$ steps, where in each time step $1 \leq i \leq |V|$, the vertex $\pi^{-1}(i)$ is being colored *Blue* if this does not form any monochromatic *Blue* edge. Otherwise, $\pi^{-1}(i)$ is colored *Red*.

We now state an important property of $\text{Col}(\mathcal{H}, \pi)$. For two disjoint subsets $X, Y \subseteq V$, we use the notation $\pi(X) < \pi(Y)$ whenever $\max_{x \in X} \pi(x) < \min_{y \in Y} \pi(y)$, that is, the elements of X precede all the elements of Y in the ordering π . Suppose (X, Y) is a simple pair of edges in \mathcal{H} with¹ $X \cap Y = y$. We say that π *separates* (X, Y) if $\pi(X \setminus y) < \pi(y) < \pi(Y \setminus y)$.

Claim 2.1. *If $\text{Col}(\mathcal{H}, \pi)$ fails to properly color \mathcal{H} then π separates at least one pair of simple edges.*

Proof. We first observe that (by definition) for every ordering π , the algorithm $\text{Col}(\mathcal{H}, \pi)$ does not produce monochromatic *Blue* edges. Suppose then it produced a *Red* edge $Y \in E$. Let y be the first vertex of Y according to the ordering π . If y was colored red, then there must have been an edge X so that $y \in X$, and all other vertices of X were already colored *Blue* (otherwise the algorithm would color y *Blue*). This means (X, Y) is simple and that π separates it. \square

Note that the claim above already shows that if \mathcal{H} is not 2-colorable then $m_2(\mathcal{H}) > 0$. For the proof of Proposition 1.1 we will also need the following simple fact.

Claim 2.2. *A random permutation separates any given simple pair with probability $1/n \binom{2n-1}{n}$.*

Proof. Let (X, Y) be a simple pair, and let $X \cap Y = y$. A permutation π separates (X, Y) if and only if $\pi(X \setminus y) < \pi(y) < \pi(Y \setminus y)$, and this happens with probability exactly

$$\frac{(n-1)!(n-1)!}{(2n-1)!} = \frac{1}{n \binom{2n-1}{n}}$$

as desired. \square

The above claims suffice for proving Proposition 1.1.

Proof (of Proposition 1.1): Assume $m_2(\mathcal{H}) < n \binom{2n-1}{n}$. Suppose we pick a uniformly random π . Then by the union bound and Claim 2.2, we infer that with positive probability π does not separate any simple pair edges. Hence, there is a π not separating any simple pair. Claim 2.1 then implies that $\text{Col}(\mathcal{H}, \pi)$ will produce a legal 2-coloring of \mathcal{H} . \square

¹Here, and in what follows, we slightly abuse notation by writing y instead of the more appropriate $\{y\}$.

3 Proof of Theorem 1

For the rest of this section fix some non 2-colorable n -graph $\mathcal{H} = (V, E)$ satisfying $m_2(\mathcal{H}) = n \binom{2n-1}{n}$. We need to show that \mathcal{H} contains a copy of K_{2n-1}^n . We start with a few preliminary claims regarding \mathcal{H} .

First, we show that no π separates more than one simple pair.

Claim 3.1. *Every ordering π separates at most one simple pair.*

Proof. Suppose π separates two simple pairs. By Claim 2.2, the assumption on $m_2(\mathcal{H})$, and by linearity of expectation, the expected number of simple pairs separated by a random permutation is exactly 1. Hence, if π separates 2 simple pairs, then there must exist a permutation σ which separates less than 1, and therefore 0, simple pairs. Therefore, by Claim 2.1 we obtain that $\text{Col}(\mathcal{H}, \sigma)$ produces a legal 2-coloring of \mathcal{H} , which is a contradiction to the assumption that \mathcal{H} is not 2-colorable. \square

Claim 3.2. *If (X, Y) and (X', Y) are simple pairs, then $X \cap Y \neq X' \cap Y$.*

Proof. We observe that if $X \cap Y = X' \cap Y = y$, then there is a π that separates both (X, Y) and (X', Y) , and this will contradict Claim 3.1. Indeed, if (X, Y) and (X', Y) are simple pairs and $X \cap Y = X' \cap Y = y$, then $(X \cup X') \setminus y$ and Y are disjoint. Therefore, any π satisfying

$$\pi((X \cup X') \setminus y) < \pi(y) < \pi(Y \setminus y)$$

separates (X, Y) and (X', Y) . This completes the proof. \square

In addition to the above observations about \mathcal{H} , the last ingredient we will need is the following theorem of Bollobás [3].

Lemma 3.3. *Let I be an index set. For all $i \in I$, let A_i and B_i be subsets of a set V of p elements satisfying the following conditions:*

- i. $A_i \cap B_i = \emptyset$ for all $i \in I$, and*
- ii. $A_j \not\subseteq A_i \cup B_i$ for all $i \neq j \in I$.*

Then, we have

$$\sum_{i \in I} \frac{1}{\binom{p-|B_i|}{|A_i|}} \leq 1,$$

with equality if and only if $B_i = B$ for all $i \in I$ and the sets A_i are all the q -tuples of the set $P \setminus B$ for some value of q .

Let us now show how to use Lemma 3.3 in order to derive Theorem 1. Recall that V is the vertex set of \mathcal{H} and set $p := |V|$. Let $M(\mathcal{H})$ be a collection of simple pairs (X, Y) defined as follows; out of all the simple pairs (X, Y) with the same ‘‘second’’ set Y , put in $M(\mathcal{H})$ one of these pairs. Observe that by Claim 3.2 each Y belongs to at most $|Y| = n$ simple pairs of the form (X, Y) (i.e, with Y as the second set), implying that $t := |M(\mathcal{H})| \geq \frac{1}{n} \cdot m_2(\mathcal{H}) = \binom{2n-1}{n}$. We now define a collection \mathcal{F} consisting of pairs of subsets of V as follows: For every simple pair $s := (X, Y) \in M(\mathcal{H})$, define $A_s = X \setminus Y$ and $B_s = V \setminus (X \cup Y)$, and let $\mathcal{F} = \{(A_s, B_s) : s \in M(\mathcal{H})\}$. For convenience, let us rename the pairs in \mathcal{F} as (A_i, B_i) with $1 \leq i \leq t$.

Now we wish to show that \mathcal{F} satisfies the conditions in Lemma 3.3. Observe that if it does, then since

$$\sum_{i=1}^t \frac{1}{\binom{p-|B_i|}{|A_i|}} = \sum_{i=1}^t \frac{1}{\binom{2n-1}{n-1}} \geq 1,$$

it follows by the first part of Lemma 3.3 that the last inequality is in fact an equality. Therefore, by the second part of Lemma 3.3, we conclude that all the B_i 's are the same set B , and the set of all the A_i 's

consists of all $n - 1$ subsets of a ground set of size $2n - 1$. That is, let $B = B_i$ and $U = V \setminus B$. Then we have that $|U| = 2n - 1$, and that the sets A_i are all the $n - 1$ subsets of U . Since by construction we have that $U \setminus A_i \in E(\mathcal{H})$ for all i , we conclude that \mathcal{H} restricted to the set U is a copy of K_{2n-1}^n as desired. It thus remains to show the following:

Claim 3.4. \mathcal{F} satisfies the conditions in Lemma 3.3

Proof. The first condition $A_i \cap B_i = \emptyset$ for all i is trivially satisfied by construction. For the second condition, let (A, B) and (A', B') be two elements in \mathcal{F} coming from simple pairs (X, Y) and (X', Y') belonging to $M(\mathcal{H})$, respectively. Recall that by the way we defined $M(\mathcal{H})$ and \mathcal{F} we have $Y \neq Y'$. Let us use y and y' to denote the unique elements in $X \cap Y$ and $X' \cap Y'$, respectively. We wish to show that $A \not\subseteq A' \cup B'$, which, by construction, is implied by $(X \setminus y) \cap Y' \neq \emptyset$. Assuming $(X \setminus y) \cap Y' = \emptyset$, we will derive a contradiction to Claim 3.1 by showing that there is a permutation π separating two distinct simple pairs.

Observe that it cannot be that $y \in Y'$. Indeed, if it was the case, then together with the assumption that $(X \setminus y) \cap Y' = \emptyset$ we would infer that (X, Y) and (X, Y') are both simple pairs intersecting at y (and distinct as $Y \neq Y'$), contradicting Claim 3.2. Assume then that $y \notin Y'$ (so in particular $y \neq y'$). We claim that we can find a π satisfying

$$\pi(X \setminus y) < \pi(y) < \pi((X' \setminus y') \setminus X) < \pi(y') < \pi((Y \cup Y') \setminus (X \cup X')).$$

Indeed, the only thing that needs to be justified is the ability to place y' as above, which follows from the fact that $y' \in Y'$ and the assumption $(X \setminus y) \cap Y' = \emptyset$ which together imply that $y' \notin X$. Observe that since π first places $X \setminus y$ and then y , the pair (X, Y) is separated by π . Such a π clearly places $X' \setminus y'$ before y' and the assumption $(X \setminus y) \cap Y' = \emptyset$ together with the fact that $y \notin Y'$ imply that such a π places all of $Y' \setminus y'$ after y' , so it separates (X', Y') as well, giving us the desired contradiction. \square

This completes the proof of Theorem 1.

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