

Lecture notes for “Analysis of Algorithms”: Dynamic All-Pairs Shortest Paths

Lecturer: *Uri Zwick*

Slightly updated code

Function `apsp($G = (V, E, w)$)`

```

 $Q \leftarrow \text{make-heap}()$ 
foreach  $v \in V$  do
   $\pi[v] \leftarrow \text{make-path}(v)$ 
foreach  $u, v \in V$  do
   $P[u, v] \leftarrow \text{make-heap}()$ 
foreach  $v \in V$  do
   $\text{insert-in-edges}(G, v)$ 
build-paths}(G)

```

Function `build-paths(G)`

```

foreach  $u, v \in V$  do
   $sp[u, v] \leftarrow \text{null}$ 
  if  $P[u, v] \neq \phi$  then
     $\pi \leftarrow \text{find-min}(P[u, v])$ 
     $\text{heap-insert}(Q, \pi)$ 
while  $Q \neq \phi$  do
   $\pi \leftarrow \text{extract-min}(Q)$ 
   $u \leftarrow \text{start}[\pi]; v \leftarrow \text{end}[\pi]$ 
  if  $sp[u, v] = \text{null}$  then
     $sp[u, v] \leftarrow \pi$ 
    if  $\neg \text{sel}[\pi]$  then
       $\text{new-shortest-path}(\pi)$ 

```

Function `insert-in-edges(G, v)`

```

foreach  $e = (u, v) \in E_{in}[v]$  do
   $\pi \leftarrow \text{make-path}(e)$ 
   $\text{insert}(L[\pi[v]], \pi)$ 
   $\text{insert}(R[\pi[u]], \pi)$ 
   $\text{heap-insert}(P[u, v], \pi)$ 

```

Function `insert-out-edges(G, v)`

```

foreach  $e = (v, u) \in E_{out}[v]$  do
   $:$ 

```

Function `new-shortest-path(π)`

```

 $\text{sel}[\pi] \leftarrow \text{true}$ 
 $\text{insert}(L^*[r[\pi]], \pi)$ 
 $\text{insert}(R^*[l[\pi]], \pi)$ 
foreach  $\pi' \in L^*[l[\pi]]$  do
   $\pi'' \leftarrow \text{make-path}(\pi', \pi)$ 
   $\text{insert}(L[\pi], \pi'')$ 
   $\text{insert}(R[\pi'], \pi'')$ 
   $u \leftarrow \text{start}[\pi'']; v \leftarrow \text{end}[\pi'']$ 
   $\text{heap-insert}(P[u, v], \pi'')$ 
   $\text{heap-insert}(Q, \pi'')$ 
foreach  $\pi' \in R^*[r[\pi]]$  do
   $:$ 

```

Function remove-paths(G, v)

```
 $S \leftarrow \text{make-stack}()$ 
push( $S, \pi[v]$ )
while  $S \neq \phi$  do
   $\pi \leftarrow \text{pop}(S)$ 
  foreach  $\pi' \in L[\pi] \cup R[\pi]$  do
    push( $S, \pi'$ )
    delete( $L[r[\pi']], \pi'$ )
    delete( $R[l[\pi']], \pi'$ )
     $u \leftarrow \text{start}[\pi']$  ;  $v \leftarrow \text{end}[\pi']$ 
    heap-delete( $P[u, v], \pi'$ )
    if sel[ $\pi'$ ] then
      delete( $L^*[r[\pi']], \pi'$ )
      delete( $R^*[l[\pi']], \pi'$ )
```

Function update(G, v)

```
remove-paths( $G, v$ )
insert-in-edges( $G, v$ )
insert-out-edges( $G, v$ )
build-paths( $G$ )
```

Function full-update(G, v)

```
 $v[t] \leftarrow v$ 
update( $G, v$ )
 $j \leftarrow 1$ 
while  $j \leq t$  do
  update( $G, v[t - j]$ )
   $j \leftarrow 2j$ 
 $t \leftarrow t + 1$ 
```

References

- [DI04] C. Demetrescu and G.F. Italiano. A new approach to dynamic all pairs shortest paths. *Journal of the ACM*, 51(6):968–992, 2004.