

Abstract Syntax

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Outline

- The general idea
- CUP
- Motivating example
Interpreter for arithmetic expressions
- The need for abstract syntax
- Abstract syntax for arithmetic expressions

Semantic Analysis during Recursive Descent Parsing

- Scanner returns “semantic values” for some tokens
- The function of every non-terminal returns the “corresponding subtree value”
- When $A \rightarrow B C D$ is applied the function for A can use the values returned by B , C , and D
 - The function can also pass parameters, e.g., to $D()$, reflecting left contexts

$E \rightarrow \mathbf{num} E'$

$E' \rightarrow \varepsilon$

$E' \rightarrow + \mathbf{num} E'$

```
int E() {  
    switch (tok) {  
        case num : temp=tok.val; eat(num);  
                    return EP(temp);  
        default: error(...); } }
```

```
int EP(int left) {  
    switch (tok) {  
        case $: return left;  
        case + : eat(+);  
                    temp=tok.val; eat(num);  
                    return EP(left + temp);  
        default: error(...) ; } }
```

Semantic Analysis during Bottom-Up Parsing

- Scanner returns “semantic values” for some tokens
- Use parser stack to store the “corresponding subtree values”
- When $A \rightarrow B C D$ is reduced the function for A can use the values returned by B , C , and D
- No action in the middle of the rule

Example

$E \rightarrow E + \text{num}$

$E \rightarrow \text{num}$

num 5

+

E 7

E 12

terminal Integer NUMBER;
terminal PLUS,MINUS,MULT,DIV;
terminal LPAREN, RPAREN;
terminal UMINUS;
non terminal Integer expr;
precedence left PLUS, MINUS;
precedence left DIV, MULT;
Precedence left UMINUS;
%%

expr ::= expr:e1 PLUS expr:e2
 {: RESULT = new Integer(e1.intValue() + e2.intValue()); :}
 | expr:e1 MINUS expr:e2
 {: RESULT = new Integer(e1.intValue() - e2.intValue()); :}
 | expr:e1 MULT expr:e2
 {: RESULT = new Integer(e1.intValue() * e2.intValue()); :}
 | expr:e1 DIV expr:e2
 {: RESULT = new Integer(e1.intValue() / e2.intValue()); :}
 | MINUS expr:e1 %prec UMINUS
 {: RESULT = new Integer(0 - e1.intValue()); :}
 | LPAREN expr:e1 RPAREN
 {: RESULT = e1; :}
 | NUMBER:n
 {: RESULT = n; :}

stack	input	action
\$	7+11+17\$	shift
num 7 \$	+11+17\$	reduce e → num
e 7 \$	+11+17\$	shift
+ e 7 \$	11+17\$	shift
num 11 + e 7 \$	+17\$	reduce e → num

stack	input	action
e 11 + e 7 \$	+17\$	reduce $e \rightarrow e+e$

e 18 \$	+17\$	shift
------------	-------	-------

+ e 18 \$	17\$	shift
-----------------	------	-------

num 17 + e 18 \$	\$	reduce $e \rightarrow \text{num}$
----------------------------------	----	-----------------------------------

stack	input	action
e 17 + e 18 \$	\$	reduce $e ::= e+e$

e 35 \$	\$	accept
------------	----	--------

So why can't we write
all the compiler code
in Bison/CUP?

Historical Perspective

- Originally parsers were written w/o tools
- yacc, bison, cup... make tools acceptable
- But it is still difficult to write compilers in parser actions (top-down and bottom-up)
 - Natural grammars are ambiguous
 - No modularity principle
 - Many useful programming language features prevent code generation while parsing
 - Use before declaration
 - gotos

Abstract Syntax

- Intermediate program representation
- Defines a tree - Preserves program hierarchy
- Generated by the parser
- Declared using an (ambiguous) context free grammar (relatively flat)
 - Not meant for parsing
- Keywords and punctuation symbols are not stored (Not relevant once the tree exists)
- Big programs can be also handled (possibly via virtual memory)

Issues

- Concrete vs. Abstract syntax tree
- Need to store concrete source position
- Abstract syntax can be defined by:
 - Ambiguous context free grammar
 - C recursive data type
 - “Constructor” functions
- Debugging routines linearize the tree

Abstract Syntax for Arithmetic Expressions

$\text{Exp} \rightarrow \mathbf{id}$	(IdExp)
$\text{Exp} \rightarrow \mathbf{num}$	(NumExp)
$\text{Exp} \rightarrow \text{Exp Binop Exp}$	(BinOpExp)
$\text{Exp} \rightarrow \text{Unop Exp}$	(UnOpExp)
$\text{Binop} \rightarrow +$	(Plus)
$\text{Binop} \rightarrow -$	(Minus)
$\text{Binop} \rightarrow *$	(Times)
$\text{Binop} \rightarrow /$	(Div)
$\text{Unop} \rightarrow -$	(UnMin)

```
package Absyn;
abstract public class Absyn { public int pos ;}
Exp extends Absyn { } ;
class IdExp extends Exp { String rep ;
    IdExp(r) { rep = r ;}
}
class NumExp extends Exp { int number ;
    NumExp(int n) { number = n ;}
}
class OpExp {
    public final static int PLUS=1; public final static int Minus=2;
    public final static int Times=3; public final static int Div=4;
}
final static int OpExp.PLUS, OpExp.Minus, OpExp.Times, OpExp.Div;
class BinExp extends Exp {
    Exp left, right; OpExp op ;
    BinExp(Exp l, OpExp o, Bin Exp r) {
        left = l ; op = o; right = r ;
    }
}
}
```



```

terminal Integer NUMBER;
terminal PLUS,MINUS,MULT,DIV;
terminal LPAREN, RPAREN;
terminal UMINUS;
nonterminal Exp expr;
precedence left PLUS, MINUS;
precedence left DIV, MULT;
Precedence left UMINUS;
%%
expr ::= expr:e1 PLUS expr:e2
      | expr:e1 MINUS expr:e2
      | expr:e1 MULT expr:e2
      | expr:e1 DIV expr:e2
      | MINUS expr:e1 %prec UMINUS
      | LPAREN expr:e1 RPAREN
      | NUMBER:n
      | RESULT = new BinExp(e1, OpExp.PLUS, e2); :}
      | RESULT = new BinExp(e1, OpExp.MINUS, e2); :}
      | RESULT = new BinExp(e1, OpExp.PLUS, e2); :}
      | RESULT = new BinExp(e1, OpExp.DIV, e2); :}
      | RESULT = new BinExp(new NumExp(0), OpExp.MINUS, e1); :}
      | RESULT = e1; :}
      | RESULT = new NumExp(n.intValue()); :}

```

Summary

- Jflex(Jflex) and CUP simplify the task of writing compiler/interpreter front-ends
- Abstract syntax provides a clear interface with other compiler phases
 - Supports general programming languages
- But the design of an abstract syntax for a given PL may take some time