

Some Advanced ML Features

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Content

- Recap last week
- Some benefits of no side-effects
- Some advanced features
 - Exceptions
 - Module System
 - Imperative Features

Class	Grammar/Meta-Variable	Examples
Identifiers	x, y	$a, x, y, x_y, \text{foo}1000$
Datatype Type constructors	X, Y	Nil, Const, List
Constants	c	2, 4.0, "ff", []
Expressions e	$e ::= x \mid c \mid e_1 e_2 \mid (e_1, e_2, \dots, e_n) \mid$ $\mid \text{if } e_1 \text{ then } e_2 \text{ else } e_3 \mid$ $\mid \text{let } [\text{rec}] d_1 \text{ and } \dots \text{ and } d_n \text{ in } e$ $\mid \text{match } e \text{ with } p_1 \rightarrow e_1 \mid \dots \mid p_n \rightarrow e_n$	
Patterns	$p ::= c \mid x [:t] \mid (p_1, \dots, p_n) \mid X \mid X(p)$	$a:\text{int}, (x:\text{int}, y:\text{int})$
Declarations	$d ::= p = e \mid y p [:t] = e \mid$ $\mid \text{datatype } Y = X_1 [\text{of } t_1] \mid \dots \mid X_n [\text{of } t_n]$	$\text{let one} = 1$ $\text{let sq}(x:\text{int}): \text{int}$
Types	$t ::= \text{int} \mid \text{float} \mid \text{string} \mid \text{char} \mid$ $\mid t_1 \rightarrow t_2 \mid t_1 * \dots * t_n \mid Y$	
Values	$v ::= c \mid (v_1, \dots, v_n) \mid X(v)$	

Factorial in ML

```
let rec fac n = if n = 0 then 1 else n * fac (n - 1)
```

```
// val fac : int -> int = <fun>
```

```
let rec fac n : int = if n = 0 then 1 else n * fac (n - 1)
```

```
let rec fac = function  
  | 0 -> 1  
  | n -> n * fac(n - 1)
```

```
let fac n =  
  let rec ifac n acc =  
    if n=0 then acc else ifac n-1, n * acc  
  in ifac n, 1
```

Plan

- Basic Programming in ML [last week, this recitation]
- Advanced ML Programming [This week]
- Javascript
- Runtime Management for PL
- Type Inference for ML
- Scala
- ...

Benefits of Functional Programming

- No side-effects
- Referential Transparency
 - The value of expression e depends only on its arguments
- Conceptual
- Commutativity
- Easier to show that the code is correct
- Easier to generate efficient implementations

Let-expressions

3 questions:

- Syntax: `let d_1 and ... and d_n in e`
 - Each d_i is any *binding* and e is any *expression*
- Type-checking: Type-check each d_i and e in a static environment that includes the previous bindings.
Type of whole let-expression is the type of e
- Evaluation: Evaluate each d_i and e in a dynamic environment that includes the previous bindings.
Result of whole let-expression is result of evaluating e .

Silly Examples

```
let silly1 (z : int) =  
  let x = if z > 0 then z else 34  
    and y = x+z+9  
  in  
    if x > y then x*2 else y*y  
  
let silly2 (z : int) =  
  let x = 1  
  in  
    (let x = 2 in x+1) +  
    (let y= x+2 in y+1)  
// val silly2 : int -> int = <fun>
```

silly2 is poor style but shows let-expressions are expressions

- Can also use them in function-call arguments, if branches, etc.
- Also notice shadowing

List Example

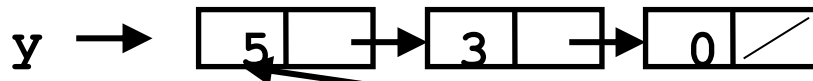
```
let rec append l1 l2 =  
  match l1 with  
  | [] -> l2  
  | hd :: tl -> hd :: append tl l2  
//val append : 'a list -> 'a list -> 'a list = <fun>
```

```
let x = [2;4] //val x : int list = [2; 4]
```

```
let y = [5;3;0] //val y : int list = [5; 3; 0]
```

```
let z = append x y
```

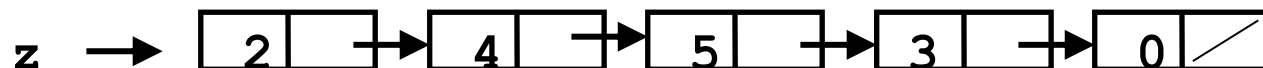
```
//val z : int list = [2; 4; 5; 3; 0]
```



or



*(can't tell,
but it's the
first one)*



Exceptions

- Captures abnormal behavior
 - Error handling
 - Integrated into the type system

```
exception Error // exception Error
let sqrt1 (x : float) : float =
  if x < 0.0 then raise Error
  else sqrt x
// val sqrt1 : float -> float = <fun>
```

```
exception FailWith of string
```

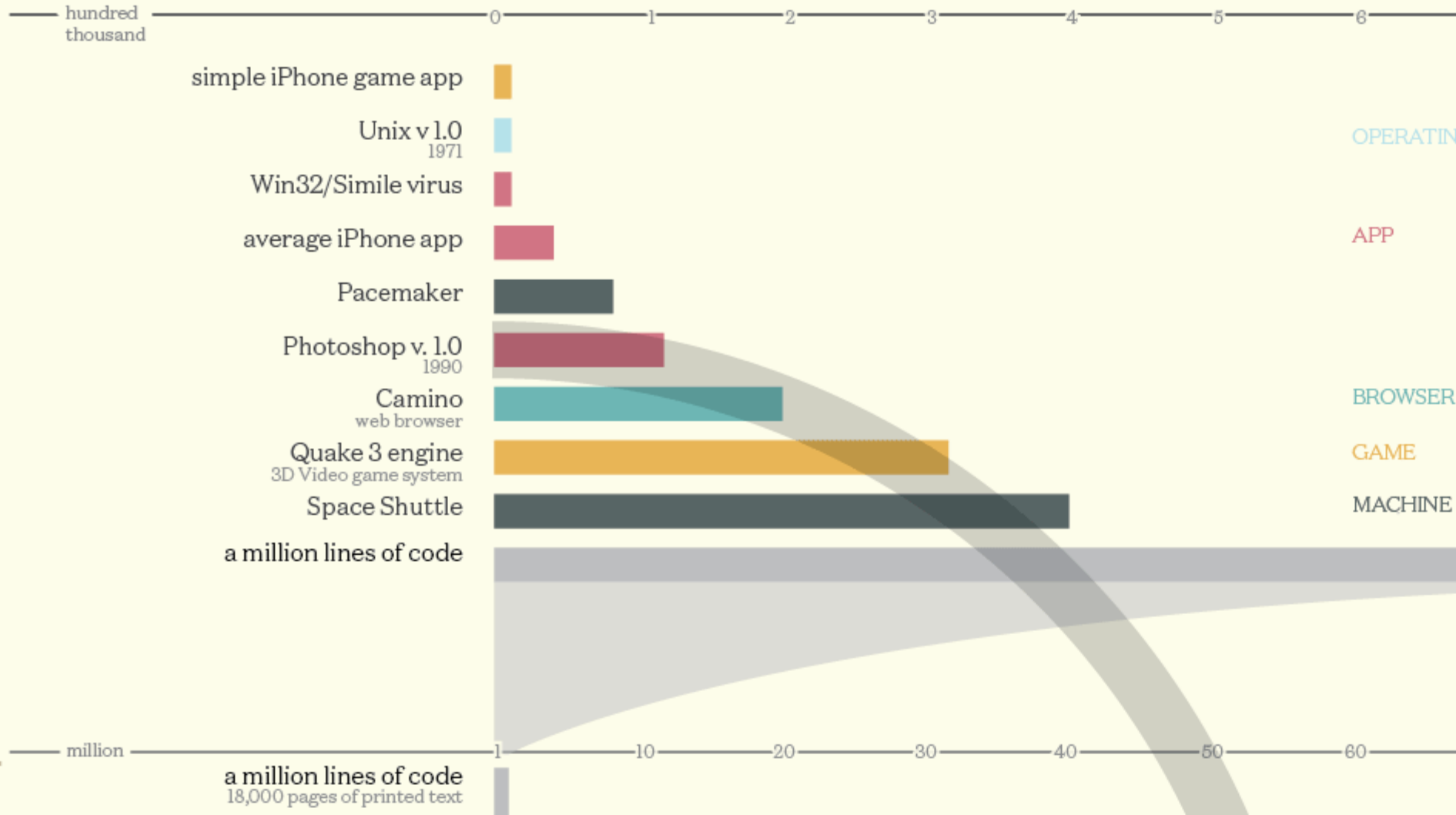
```
raise (FailWith "Some error message")
```

Question?

- What's the largest program you've ever worked on by yourself or as part of a team?
 - A. 10-100 LoC
 - B. 100-1,000 LoC
 - C. 1,000-10,000 LoC
 - D. 10,000-100,000 LoC
 - E. $\geq 100,000$ LoC

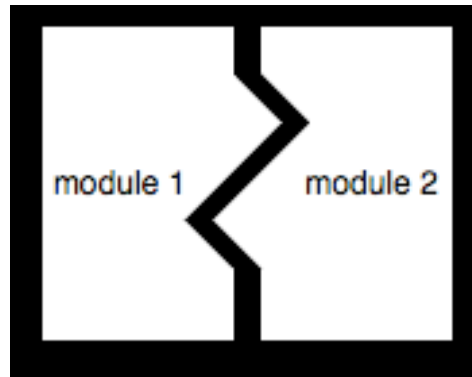
Codebases

Millions of lines of code



Modularity

- **Modular programming**: code comprises independent *modules*
 - developed separately
 - understand behavior of module in isolation
 - reason locally, not globally



Java features for modularity

- **classes, packages:**
 - organize identifiers (classes, methods, fields, etc.) into namespaces
- **Interfaces:**
 - describe related classes
- **public, protected, private:**
 - control what is visible outside a namespace

Ocaml Features for Modularity

- **modules:** organize identifiers (functions, values, etc.) into namespaces
- **signatures:**
 - describe related modules
- **abstract types:**
 - control what is visible outside a namespace

Ocaml Modules

- Syntax:
`module` ModuleName = `struct` *definitions* `end`
- The name must be capitalized
- The definitions can be any top-level definitions
 - `let`, `type`, `exception`
- Create a new namespace
- Every file `myFile.ml` with contents *D* is essentially wrapped in a module definition
`module` MyFile = `struct` *D* `end`
- *Modules can be opened locally to save writing*

```
module M = struct let x = 42 end
//module M : sig val x : int end
let fortytwo = M.x
// val fortytwo : int = 42
```


Stack Module

```
module Stack = struct
  let empty = []
  let is_empty s = s = []
  let push x s = x :: s
  let pop s = match s with
    [] -> failwith "Empty"
  | x::xs -> (x,xs)
end
//module Stack :
// sig
//   val empty : 'a list
//   val is_empty : 'a list -> bool
//   val push : 'a -> 'a list -> 'a list
//   val pop : 'a list -> 'a * 'a list
// end
```

```
fst (Stack.pop
(Stack.push 1 Stack.empty))
//- : int = 1
```

Might Seem Backwards...

Java:

```
s = new Stack();  
s.push(1);  
s.pop()
```

OCaml:

```
let s = Stack.empty in  
  let s' = Stack.push 1 s in  
    let (one, _) = Stack.pop s'
```

Abstraction

- Forgetting Information
- Treating different things as identical

Abstraction

- Programming language **predefined abstractions**
 - Data structures like list
 - Library functions like map and fold
- Programming languages enable to define **new abstractions**
 - Procedural abstractions
 - Data abstraction
 - Iterator abstraction

Procedural Abstraction

- Abstract implementation details
 - `sqrt : float -> float`
- `List.sort : ('a -> 'a -> int) -> 'a list -> 'a list`
- Abstracts how the functions are implemented
 - Both the implementation and the usage should obey the type contract
 - The implementation can assume the right type
 - Important for composing functions

Data Abstraction

- Abstract from details of organizing data
 - stacks, symbol tables, environments, bank accounts, polynomials, matrices, dictionaries, ...
- Abstract from implementation of organization:
 - Actual code used to add elements (e.g.) isn't Important
 - But types of operations and assumptions about what they do and what they require are important

Ocaml Advanced Modularity Features

- Functors and Signatures
- Functions from Modules to Modules
- Permit
 - Dependency injection
 - Swap implementations
 - Advanced testing

Stack Abstract Data Type

```
module type STACK = sig
  val empty : 'a list
  val is_empty : 'a list -> bool
  val push : 'a -> 'a list -> 'a list
  val pop : 'a list -> 'a * 'a list
end
module Stack : STACK = struct
  ... (* as before *)
end
```


Stack with Abstract Data Types

```
module type STACK = sig
  type t
  val empty : t
  val is_empty : t -> bool
  val push : int -> t -> t
  val pop : t -> int * t
end
module Stack : STACK = struct
  type t = int list
  let empty = [ ]
  let is_empty s = s = [ ]
  let push x s = x :: s
  let pop s = match s with
    [ ] -> failwith "Empty"
  | x::xs -> (x,xs)
end
```

Summary Modularity

- ML provides flexible mechanisms for modularity
- Guarantees type safety

Side-Effects

- But sometimes side-effects are necessary
- The whole purpose of programming is to conduct side-effects
 - Input/Output
- Sometimes sharing is essential for functionality
- ML provides mechanisms to capture side-effects
 - Enable efficient handling of code with little side effects

Input/Output

```
print_string: String -> Unit
```

```
let x = 3 in  
  let () = print_string ("Value of x is " ^ (string_of_int x)) in  
  x + 1  
;;value of x is 3- : int = 4
```

```
e ::= ... | ( e1; ... ; en )
```

```
let x = 3 in  
  (print_string ("Value of x is " ^ (string_of_int x));  
   x + 1)
```

Refs and Arrays

- Two built-in data-structures for implementing shared objects

```
module type REF =  
  sig  
    type 'a ref  
    (* ref(x) creates a new ref containing x *)  
    val ref : 'a -> 'a ref  
    (* !x is the contents of the ref cell x *)  
    val (!) : 'a ref -> 'a  
    (* Effects: x := y updates the contents of x  
     * so it contains y. *)  
    val (:=) : 'a ref -> 'a -> unit  
  end
```

Simple Ref Examples

```
let x : int ref = ref 3
in
  let y : int = !x
  in
    (x := !x + 1);
    y + !x
// - : int = 7
```

More Examples of Imperative Programming

- Create cell and change contents

```
val x = ref "Bob";  
x := "Bill";
```



- Create cell and increment

```
val y = ref 0;  
y := !y + 1;
```



- While loop

```
val i = ref 0;  
while !i < 10 do i := !i + 1;  
i;
```

Summary References

- Provide an escape for imperative programming
- But insures type safety
 - No dangling references
 - No (double) free
 - No null dereferences
- Relies on automatic memory management

Functional Programming Languages

PL	types	evaluation	Side-effect
scheme	Weakly typed	Eager	yes
ML OCAML F#	Polymorphic strongly typed	Eager	References
Haskell	Polymorphic strongly typed	Lazy	None

Things to Notice

- Pure functions are easy to test

```
prop_RevRev l = reverse(reverse l) == l
```

- In an imperative or OO language, you have to
 - set up the state of the object and the external state it reads or writes
 - make the call
 - inspect the state of the object and the external state
 - perhaps copy part of the object or global state, so that you can use it in the post condition

Things to Notice

Types are everywhere.

```
reverse :: [w] -> [w]
```

- Usual static-typing panegyric omitted...
- In ML, **types express high-level design**, in the same way that UML diagrams do, with the advantage that the type signatures are machine-checked
- Types are (almost always) optional: type inference fills them in if you leave them out

Information from Type Inference

- Consider this function...

```
let reverse ls = match ls with  
  [] -> []  
  | x :: xs -> reverse xs
```

... and its most general type:

```
:- reverse :: list 't_1 -> list 't_2
```

- What does this type mean?

Reversing a list should not change its type, so there must be an error in the definition of reverse!

Recommended ML Textbooks

- L. C. PAULSON: ML for the Working Programmer
- J. Ullman: Elements of ML Programming
- R. Harper: Programming in Standard ML

Recommended Ocaml Textbooks

- Xavier Leroy: The OCaml system release 4.02
 - Part I: Introduction
- Jason Hickey: Introduction to Objective Caml
- Yaron Minsky, Anil Madhavapeddy, Jason Hickey: Real World Ocaml

Summary

- Functional programs provide concise coding
- Compiled code compares with C code
- Successfully used in some commercial applications
 - F#, ERLANG, Jane Street
- Ideas used in imperative programs
- Good conceptual tool
- Less popular than imperative programs