

Improved Orientations  
of  
Physical Networks

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HAIFA

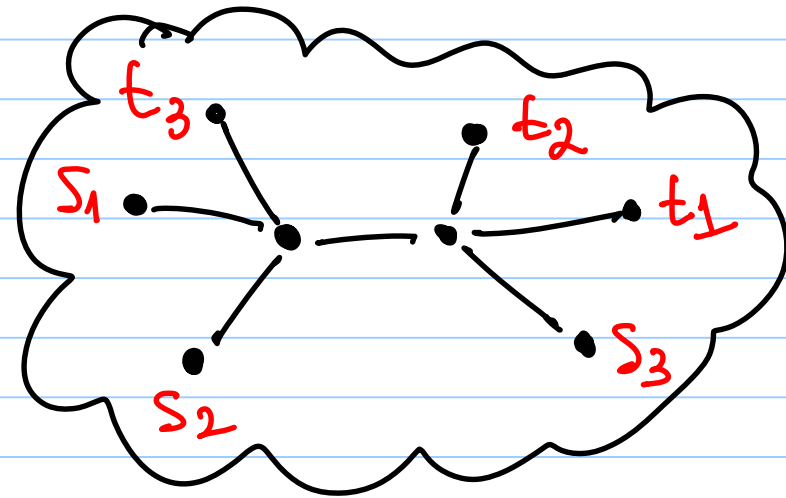
Roded Shafan  
TAU

## MOTIVATION

- Protein-protein interactions propagate **directional signals** in the cell
- Current measurement techniques capture only the presence of an interaction but **not its direction**
- Indirect directionality information can be obtained from knockout experiments that expose **cause-effect pairs**
- Goal: use pair information to **infer interaction directions**

## MAXIMUM GRAPH ORIENTATION

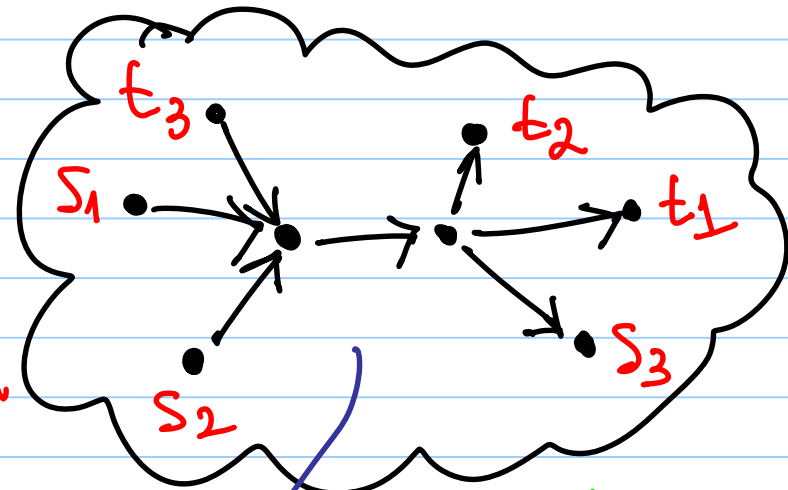
- Input: undirected graph  $G=(V,E)$  with  $n$  vertices  
source-target pairs  $(s_1, t_1), \dots, (s_k, t_k)$



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in which the number of  
connected pairs is maximized

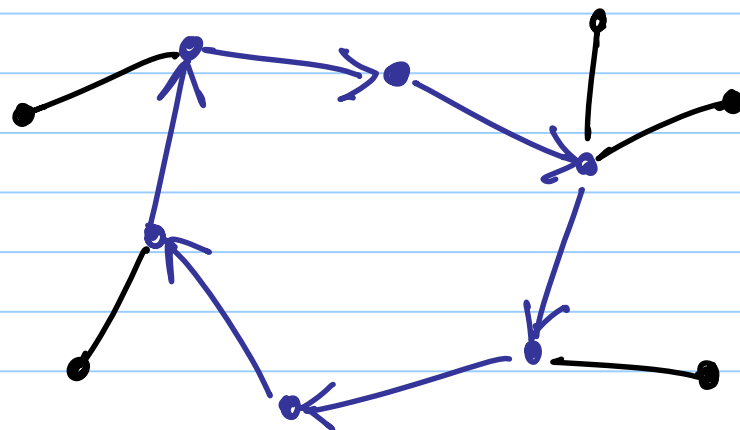


- $(s_1, t_1)$  ✓
- $(s_2, t_2)$  ✓
- $(s_3, t_3)$  ✗

## MAXIMUM GRAPH ORIENTATION

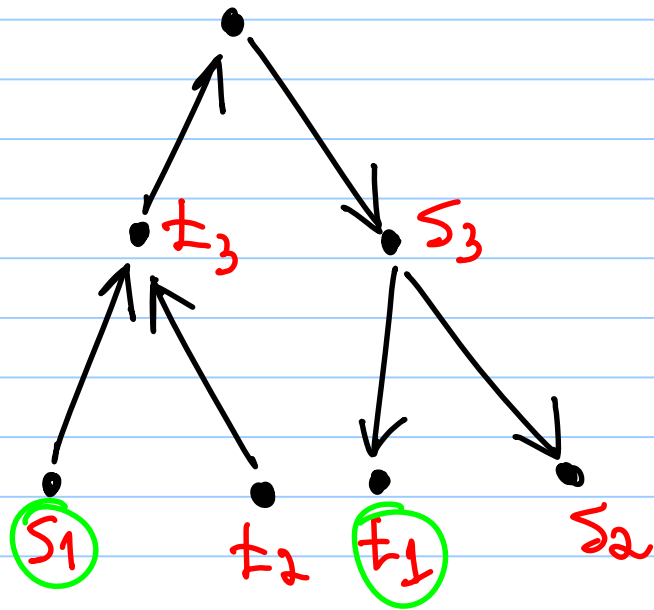
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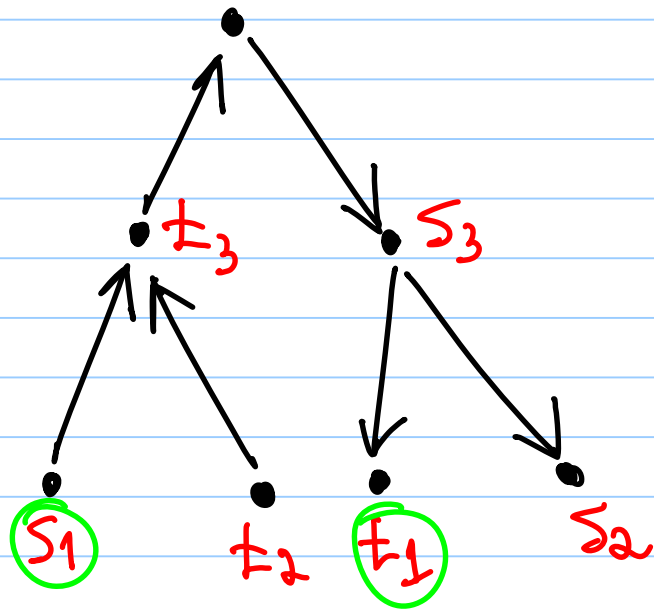
- remark: we may assume that the underlying graph is a tree

# EXAMPLE

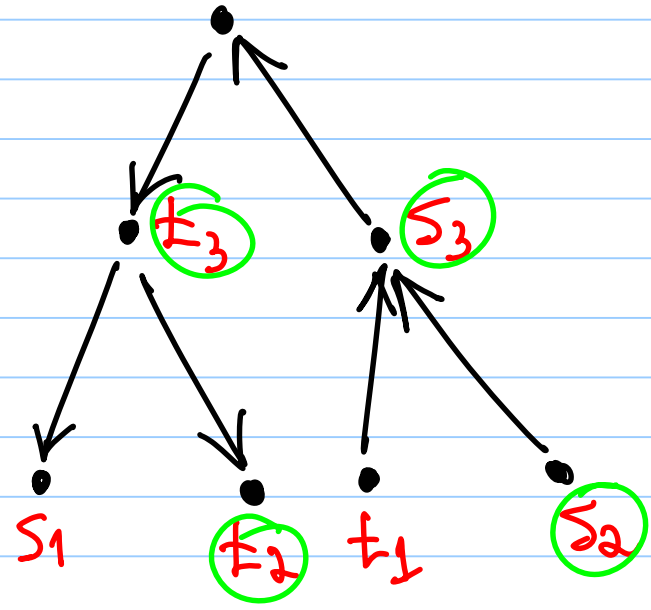


# connected: 1

# EXAMPLE



# connected: 1



# connected: 2

## PREVIOUS WORK

Medvedovsky, Bafna, Zwick, and Sharan [WABI '08]:

- can be solved to **optimality** on **paths**

- **star networks**  $\iff$  **max directed cut**

lower bounds: **0.916**, unless  **$P=NP$**  (Hastad, [JACM '01])

**0.878**, **UGC** (Khot et al. [SIComp '07])

upper bound: **0.874** (Feige and Goemans '95)

- **$O(\log n)$ -approx** for **arbitrary trees**  $\leftarrow$

## MAIN RESULT

deterministic algorithm with ratio

$$O(\log n / \log \log n)$$

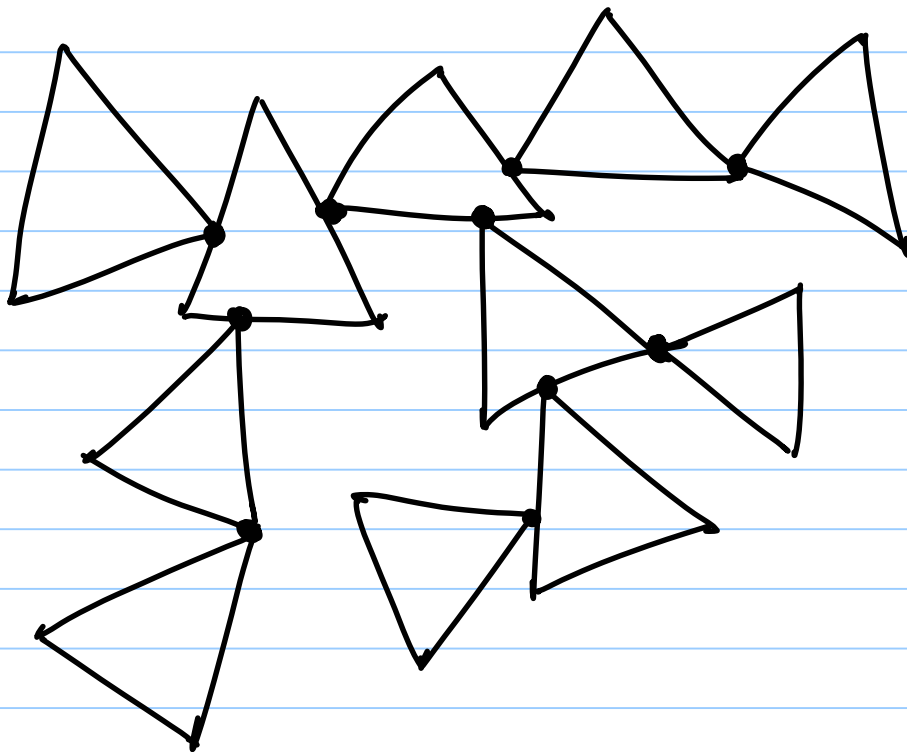
for maximum tree orientation

THE ALGORITHM:

GENERAL IDEA

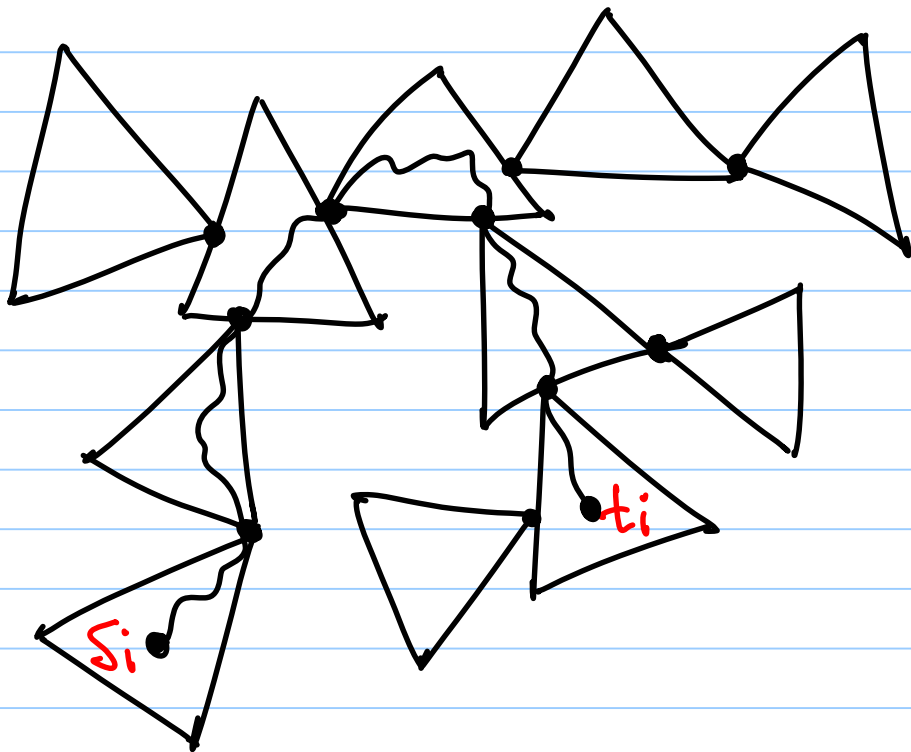
## ALMOST BALANCED $k$ -DECOMPOSITIONS

- any tree  $T = (V, E)$  can be decomposed into  $k$  edge-disjoint subtrees, each containing  $O(|E|/k)$  edges



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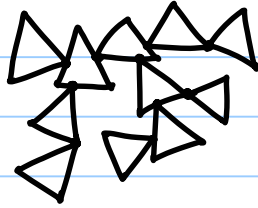
- any tree  $T = (V, E)$  can be decomposed into  $k$  edge-disjoint subtrees, each containing  $\Theta(|E|/k)$  edges



- pair  $(s_i, t_i)$  separated when  $s_i$  and  $t_i$  reside in different subtrees

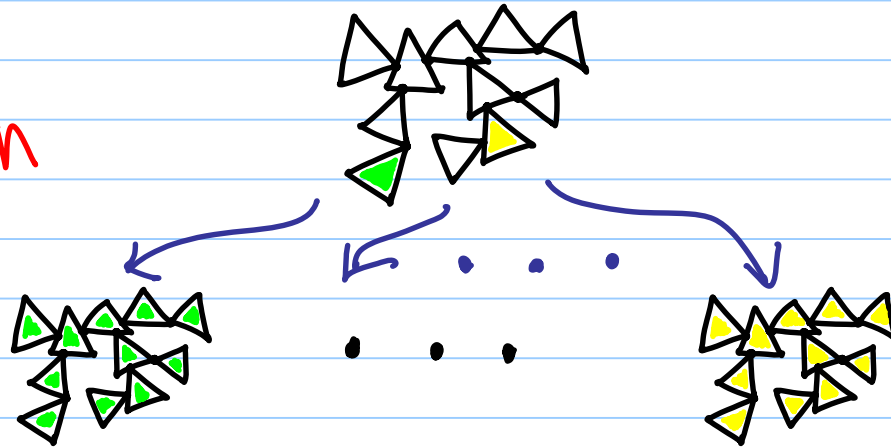
# REPEATED $k$ -DECOMPOSITIONS

pick  $k = \log n$



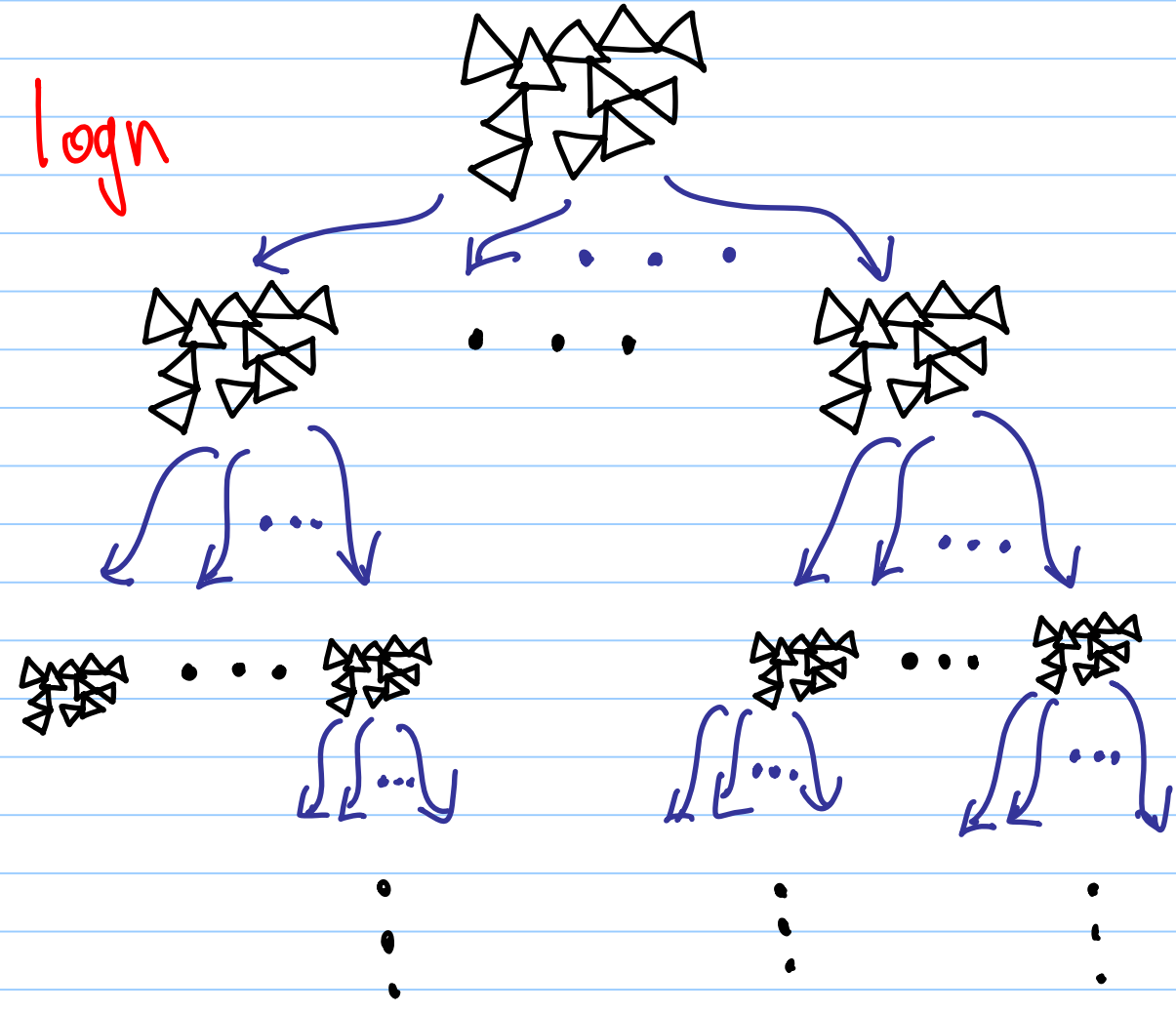
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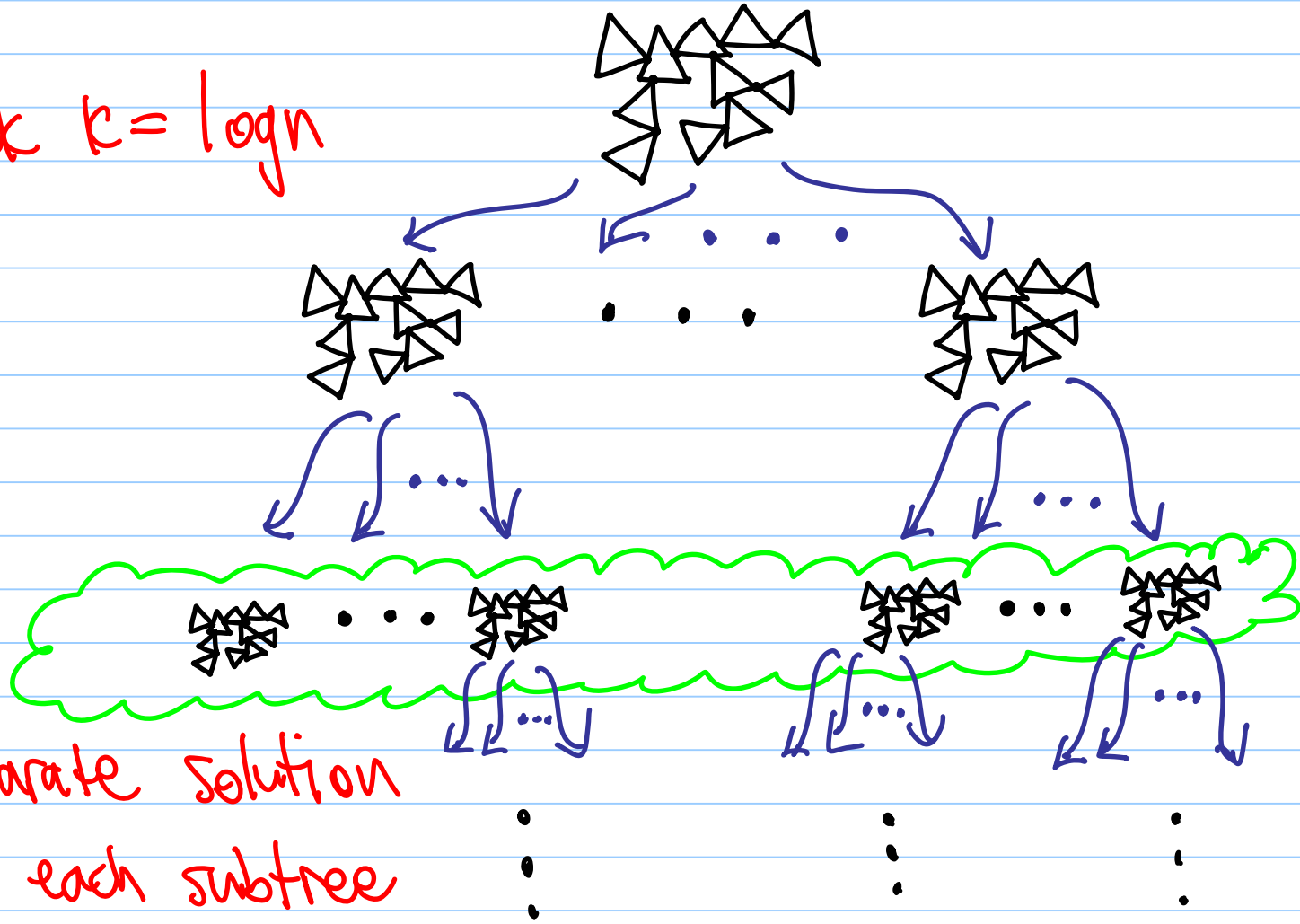
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$O\left(\frac{\log n}{\log \log n}\right)$   
levels

# REPEATED K-DECOMPOSITIONS

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separate solution  
for each subtree

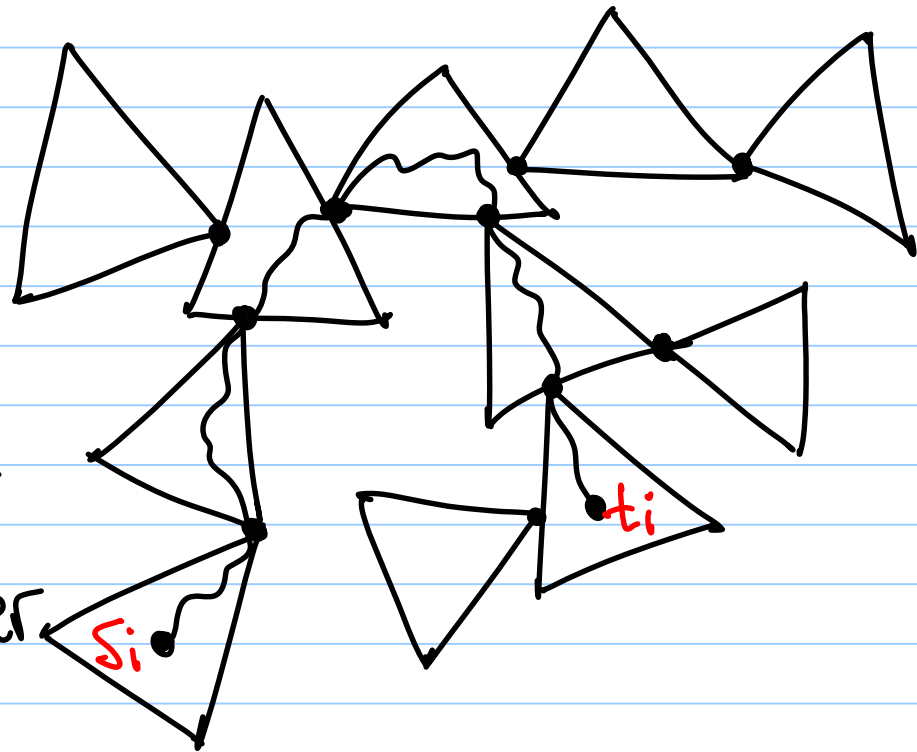
$O\left(\frac{\log n}{\log \log n}\right)$   
levels

NEED  $O(1)$ -APPROX FOR...

- Input: tree  $T=(V,E)$ , decomposed into  $T_1, \dots, T_k$   
source-target pairs  $(s_1, t_1), \dots, (s_k, t_k)$

- additional structure:  
all pairs are separated

- goal: find orientation that  
maximizes the number  
of connected pairs



CONSTANT-FACTOR APPROXIMATION

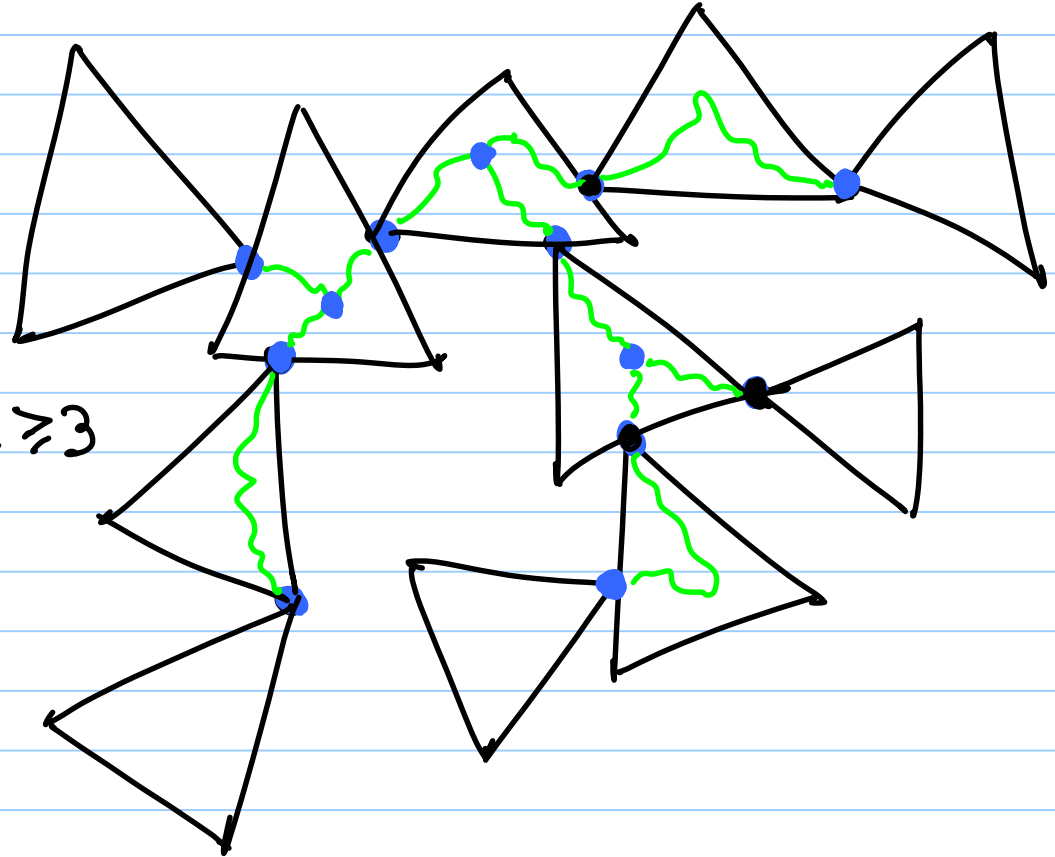
FOR

K-DECOMPOSITIONS

## NOTATION AND TERMINOLOGY

- **skeleton  $\mathcal{S}$**  = minimal subtree that spans all intersection vertices

- **special vertices** =  
intersection vertices +  
vertices in  $\mathcal{S}$  with degree  $\geq 3$

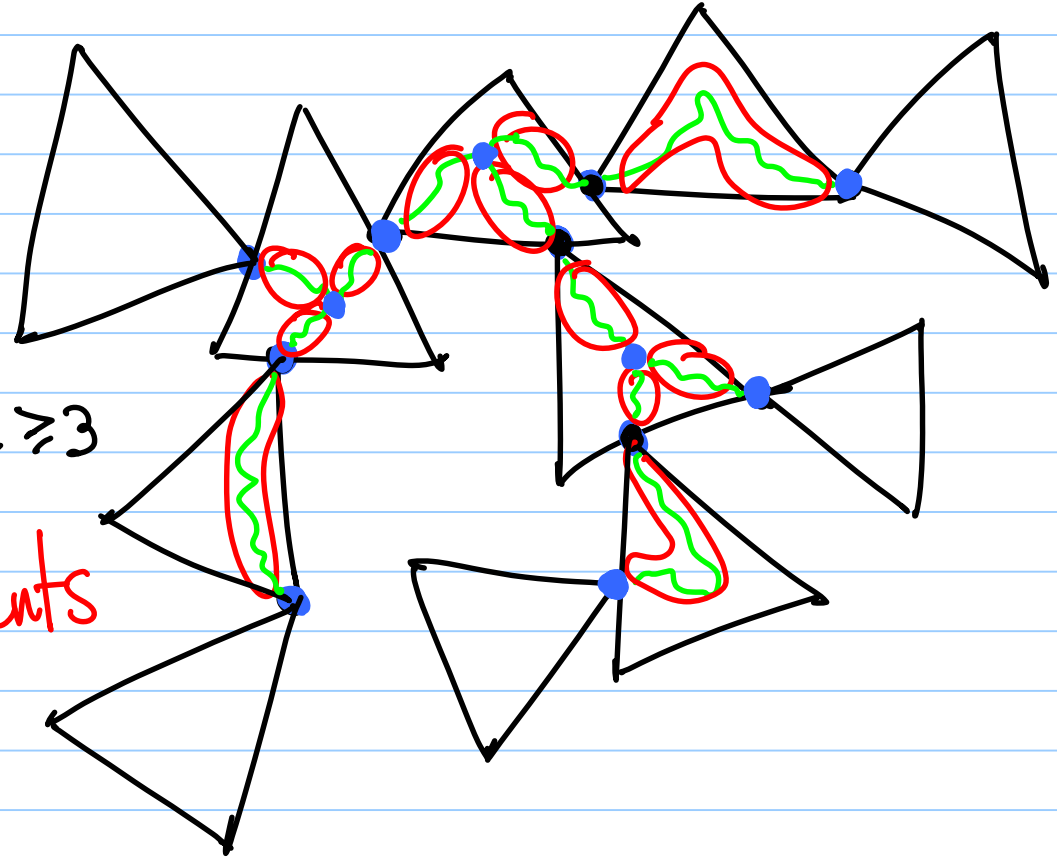


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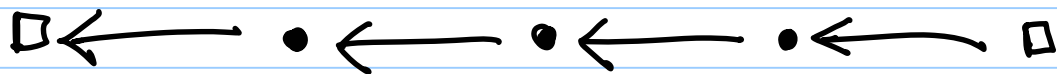
- $\Sigma(\mathcal{S})$  = set of segments




## SEGMENT GUESSING

In optimal orientation, each segment  $\alpha \in \Sigma(\delta)$  can be

(1) right-oriented: 

(2) left-oriented: 

(3) mixed: 

## SEGMENT GUESSING

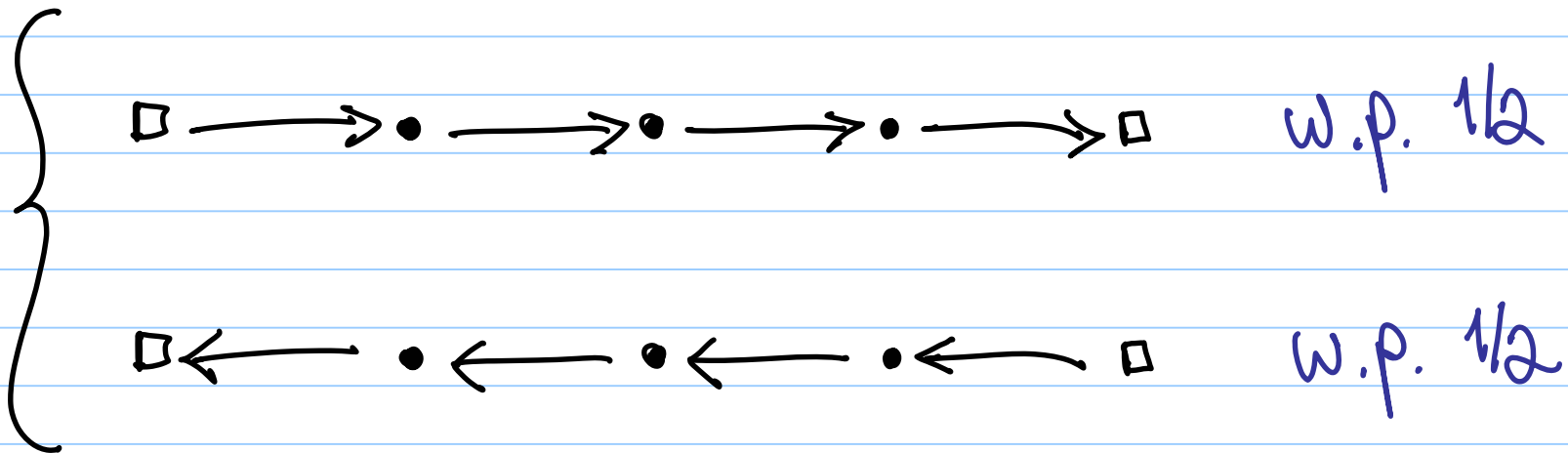
- assume that we know the optimal orientation of all skeleton segments  $\sigma \in \Sigma(\mathcal{S}')$
- number of guesses is:

$$\begin{aligned} 3^{|\Sigma(\mathcal{S}')|} &= 3^{o(k)} \\ &= 3^{o(\log n)} \\ &= 2^{o(n)} \end{aligned}$$

## THE RANDOMIZED ORIENTATION

(1) for **skeleton segments**:

- if right or left oriented, set accordingly
- if **mixed**, pick **direction at random**



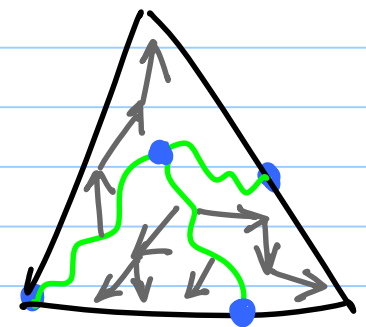
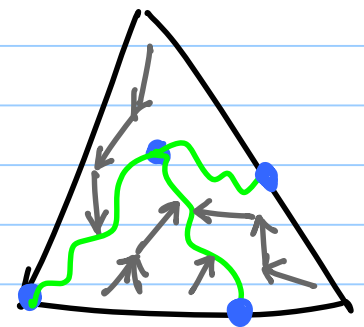
## THE RANDOMIZED ORIENTATION

(2) for each of the decomposition trees:

- randomly pick whether **SENDER** or **RECEIVER**

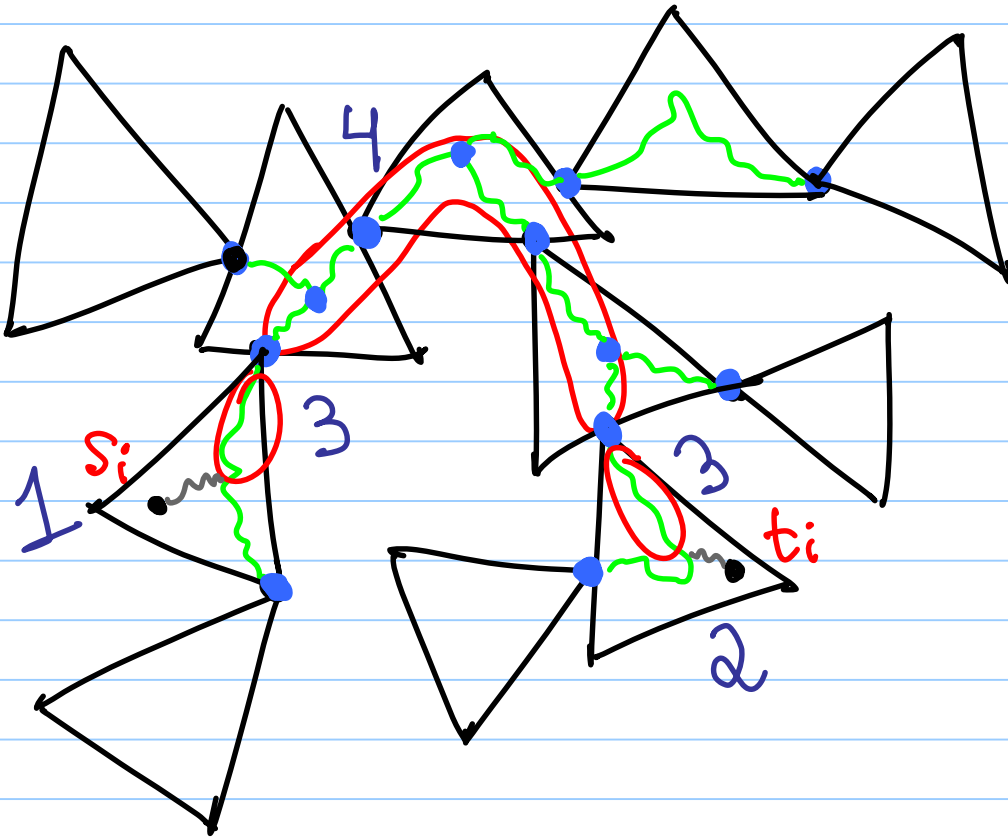
- in **RECEIVERS**: all non-skeleton edges are directed toward the skeleton

- in **SENDERS**: all non-skeleton edges are directed away from the skeleton



## ANALYSIS

Claim: each pair connected by OPT is connected by the random orientation with probability  $\geq 1/16$



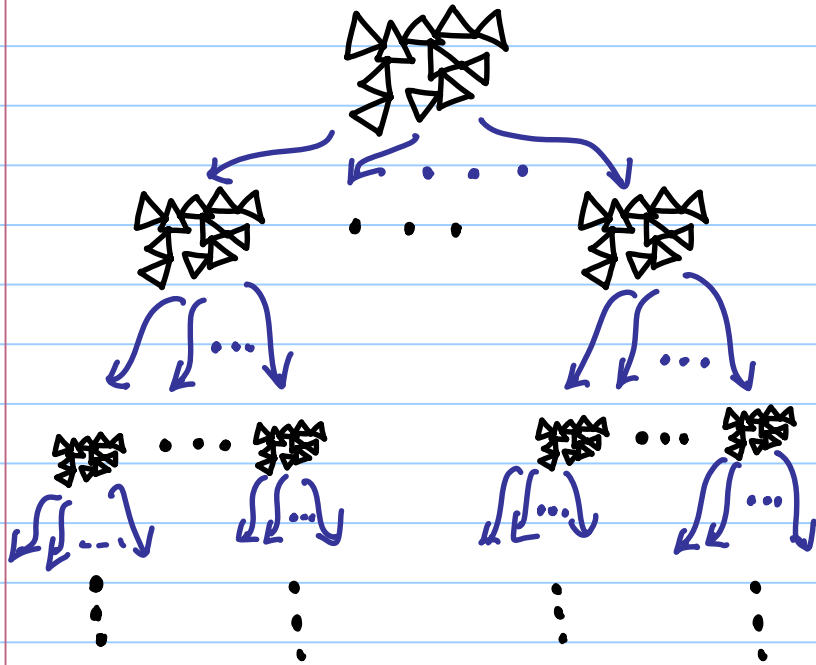
(1) SENDER w.p.  $1/2$

(2) RECEIVER w.p.  $1/2$

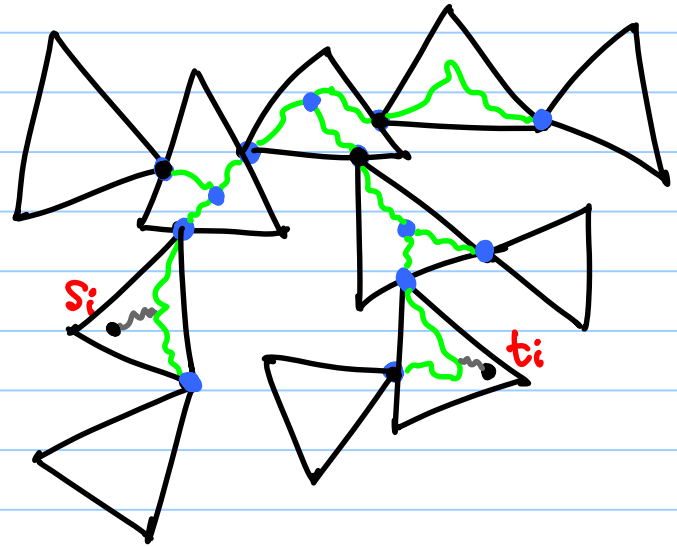
(3) as in OPT w.p.  $\geq 1/2$

(4) as in OPT w.p. 1

# QUICK RECAP

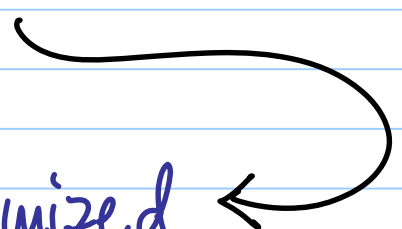


$$O\left(\frac{\log n}{\log \log n}\right)$$



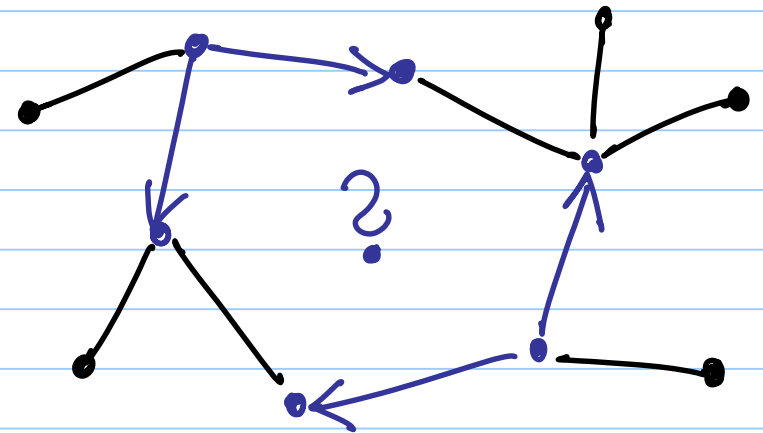
$$O(i)$$

can be derandomized



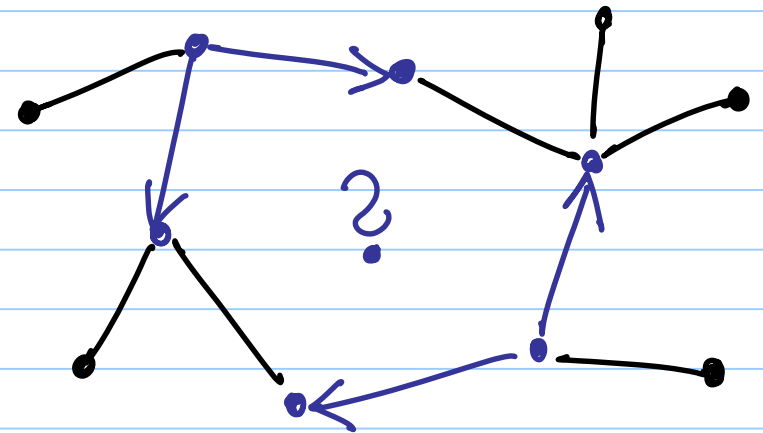
## OPEN QUESTIONS

- improved approximation factor?
- predetermined edge directions (mixed graphs)?
  - reduction to trees **does not work**
  - $O(\log^M n)$ -approx, where  
 $M \approx$  number of PDI segments  
in the underlying pathways



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THANK YOU