

# Matrix Perturbation Theory and its Applications

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# Prologue

Our hero is the intrepid, yet sensitive matrix  $A$ .  
Our villain is  $E$ , who keeps perturbing  $A$ .  
When  $A$  is perturbed he puts on a crumpled hat:  $\tilde{A} = A + E$ .

G. W. Stewart and J.-G. Sun, Matrix Perturbation Theory (1990)

# Introduction

- In a nutshell : how does a small change in the input affects the output?
- Given a matrix  $A$  and some function  $\phi$  which operates on  $A$ , we are interested in understanding how a small perturbation added to the matrix, affects the behavior of  $\phi$ .
- That is, to understand the relation between  $\phi(A)$  and  $\phi(A + \epsilon)$ , where  $\epsilon$  is a small perturbation (can be noise).
- Interesting  $\phi$  operators : finding the singular values of  $A$ , the eigenvectors  $A$ ,  $\|A\|$  and so on.

# Introduction

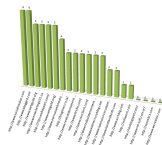
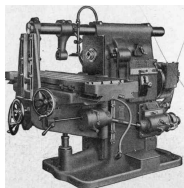
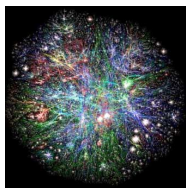
- Example application(1) : Computation of the Google page rank algorithm - how to update the ranks without recomputing the entire algorithm.
- Example application(2) : updating a training set profile with small changes to the input set.
- Books
  - Matrix Analysis - R. Bhatia.
  - Matrix Perturbation Theory - Stewart and Sun.

# Example Application - Google PageRank Calculation I

- Pagerank : The importance of a web page is set by the number of important pages pointing to it.
- $r(P) = \sum_{Q \in B_P} \frac{r(Q)}{|Q|}$  where  $B_P = [\text{all pages pointing to } P]$ ,  $|Q| = [\text{links out of } Q]$ .
- Random walk over the entire web (the probability to reach it).
- Can be calculated by iterating  $\pi_j^T = \pi_{j-1}^T P$
- Here  $P$  is a matrix with  $p_{ij} = \frac{1}{|P_i|}$  if  $P_i$  links to  $P_j$  (0 otherwise)
- The pagerank vector will be  $\pi^T = \lim_{j \rightarrow \infty} \pi_j^T$ . “The stationary probability distribution vector”.

# Example Application - Google PageRank Calculation II

- How to change  $P$  to be stochastic and irreducible (no looped chains)?
- Change the transition probability matrix to be  $\tilde{P} = \alpha P + (1 - \alpha) \frac{1}{n} \mathbf{e} \mathbf{e}^T$ .
- This should run on billions of pages (Takes Google days to run it).
- What if I added a new link to my homepage?



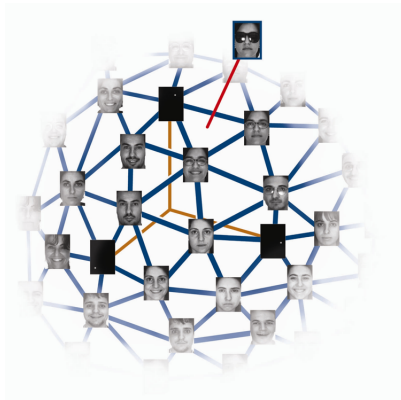
# Another Example - Face Recognition Application

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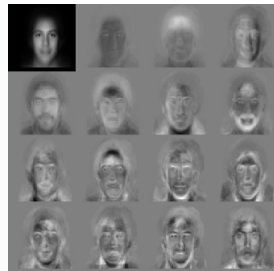
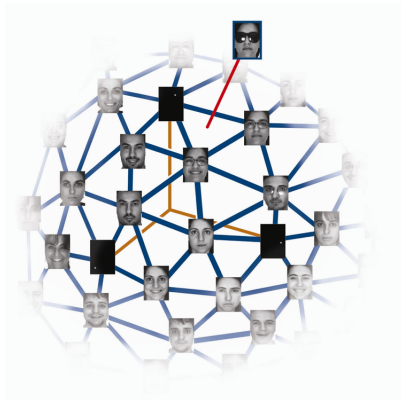




# Another Example - Face Recognition Application



# Another Example - Face Recognition Application



# Eigenpair Approximation

- Using matrix perturbation theory to update the eigenpairs.
- Can update the left principal eigenvector  $\pi$  of a stochastic matrix  $P$  where  $\pi = \pi P$  (stationary distribution of a Markov chain).
- Such methods can accelerate algorithms like Pagerank and HIT that use the stationary distribution values as rating scores.<sup>1 2</sup>
- Suitable for updating the principle eigenvector of the perturbed matrix. eigenvectors.

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<sup>1</sup>Adaptive methods for the computation of PageRank. Kamvar, Haveliwala and Golub, 2004.

<sup>2</sup>Updating Markov chains with an eye on Google's PageRank. Langville and Meyer, 2006.

# Eigenpair Approximation I

$$\begin{pmatrix} \pi_1 \\ \vdots \\ \pi_n \end{pmatrix}^T = \begin{pmatrix} \pi_1 \\ \vdots \\ \pi_n \end{pmatrix}^T \begin{pmatrix} p_{11} & \cdots & p_{1n} \\ \vdots & & \vdots \\ p_{n1} & \cdots & p_{nn} \end{pmatrix}$$

- **Power Iteration Method**

- iterates on  $\phi_{next} = \frac{A\phi}{\|A\phi\|}$
- converges to the (dominant) eigenvector of the largest eigenvalue

- **Adaptive Power Method** uses the fact that most coordinates of the eigenvector become stable within few iterations, and we can compute only ones which have not converged.

# Eigenpair Approximation II

- **Aggregated Power Iteration** reduces the unchanged states of the Markov chain into a single super state, and creates a smaller matrix. This seed eigenvector is used as the guess for each full power iteration.

# Updating A low Dimensional Representation

- We start with a symmetric matrix  $A$  which is the affinity matrix of the dataset.
- That is,  $[A]_{ij}$  is the similarity level between elements  $i$  and  $j$ .  $A$  can be computed in various ways using different kernels and distance metrics.
- A low dimensional embedding for the dataset is computed using the spectral decomposition of  $A$ .

# Updating A low Dimensional Representation

## - Cont.

- We are now given the perturbation matrix  $\tilde{A}$  of the matrix  $A$ .
- We can assume that the perturbations are sufficiently small, that is  $\|\tilde{A} - A\| < \varepsilon$  for some small  $\varepsilon$ .
- We also assume that  $\tilde{A}$  is symmetric since we compute it in the same way as  $A$  using the updated  $\tilde{X}$ .
- We wish to update the perturbed eigenpairs of  $\tilde{A}$  based on  $A$  and its eigenpairs.

# Updating A low Dimensional Representation

## - Cont.

- Given a symmetric  $n \times n$  matrix  $A$  with its  $k$  dominant eigenvalues  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_k$  and eigenvectors  $\phi_1, \phi_2, \dots, \phi_k$ , respectively, and a perturbed matrix  $\tilde{A}$  such that  $\|\tilde{A} - A\| < \varepsilon$ , find the perturbed eigenvalues  $\tilde{\lambda}_1 \geq \tilde{\lambda}_2 \geq \dots \geq \tilde{\lambda}_k$  and its eigenvectors  $\tilde{\phi}_1, \tilde{\phi}_2, \dots, \tilde{\phi}_k$  of  $\tilde{A}$  in the most efficient way.



# Computing the Eigenpairs First Order Approximation I

- Compute the approximation of each eigenpair.
- Given an eigenpair  $(\phi_i, \lambda_i)$  of a symmetric matrix  $A$  where  $A\phi_i = \lambda_i\phi_i$ , we compute the first order approximation of the eigenpair of the perturbed matrix  $\tilde{A} = A + \Delta A$ .
- We assume that the change  $\Delta A$  is sufficiently small, which result in a small perturbation in  $\phi_i$  and  $\lambda_i$ .
- We look for  $\Delta\lambda_i$  and  $\Delta\phi_i$  that satisfy the equation

$$(A + \Delta A)(\phi_i + \Delta\phi_i) = (\lambda_i + \Delta\lambda_i)(\phi_i + \Delta\phi_i). \quad (1)$$

# Eigenpair Approximation

## Theorem

If  $(\phi_i, \lambda_i)$  is an Eigenpair of  $A$  and  $\tilde{A} = A + \Delta A$  then

$$\tilde{\lambda}_i = \lambda_i + \phi_i^T (\tilde{A} - A) \phi_i + o(\|\Delta A\|^2), \quad (2)$$

$$\tilde{\phi}_i = \phi_i + \sum_{j \neq i} \frac{\phi_j^T (\tilde{A} - A) \phi_i}{\lambda_i - \lambda_j} \phi_j + o(\|\Delta A\|^2). \quad (3)$$

# The Recursive Power Iteration (RPI) Algorithm

## Overview

- Power Iteration method has proved to be effective when calculating the principle eigenvector of a matrix.
- In the RPI algorithm the first order approximation of the eigenpairs of  $A$  will be the initial guess for the power iteration method.

# The Recursive Power Iteration (RPI) Algorithm - Cont.

## Overview

- The first order approximation should be close to the actual solution we seek and therefore requires fewer iteration steps to converge.
- Once the eigenvector  $\tilde{\phi}_i$  is obtained in step  $i$ , we transform  $\tilde{A}$  into a matrix that has  $\tilde{\phi}_{i+1}$  as its principle eigenvector. We iterate this step until we recover the  $k$  dominant eigenvectors of  $\tilde{A}$ .

# The Recursive Power Iteration (RPI) Algorithm - Cont.

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## Algorithm 1: Recursive Power Iteration Algorithm

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**Input:** Perturbed symmetric matrix  $\tilde{A}_{n \times n}$ , number of eigenvectors to calculate  $k$ , initial eigenvectors guesses  $\{v_i\}_{i=1}^k$ , admissible error  $err$

**Output:** Approximated eigenvectors  $\{\tilde{\phi}_i\}_{i=1}^k$ , approximated eigenvalues  $\{\tilde{\lambda}_i\}_{i=1}^k$

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# The Recursive Power Iteration (RPI) Algorithm - Cont.

```
1: for  $i = 1 \rightarrow k$  do  
2:    $\phi \leftarrow v_i$   
3:   repeat  
4:      $\phi_{next} \leftarrow \frac{\tilde{A}\phi}{\|\tilde{A}\phi\|}$   
5:      $err_\phi \leftarrow \|\phi - \phi_{next}\|$   
6:      $\phi \leftarrow \phi_{next}$   
7:   until  $err_\phi \leq err$   
8:    $\tilde{\phi}_i \leftarrow \phi$   
9:    $\tilde{\lambda}_i \leftarrow \frac{\tilde{\phi}_i^T \tilde{A} \tilde{\phi}_i}{\tilde{\phi}_i^T \tilde{\phi}_i}$   
10:   $\tilde{A} \leftarrow \tilde{A} - \tilde{\phi}_i \tilde{\lambda}_i \tilde{\phi}_i^T$   
11: end for
```

# Correctness of the RPI Algorithm

- The correctness of the RPI algorithm is proved based on the fact that the power iteration method converges, and on the spectral decomposition properties of  $\tilde{A}$ .
- In step  $i$  we find the  $i$  largest eigenpair using the power method with the first order approximation as the initial guess.
- We then subtract the matrix  $\tilde{\phi}_i \tilde{\lambda}_i \tilde{\phi}_i^T$  from  $\tilde{A}$ . This step force the next eigenpair to become the principal eigenpair which will be found on the next step.
- We use the fact that  $\tilde{A}$  is symmetric and has a spectral decomposition of the form  $\tilde{A} = \sum_{i=1}^n \tilde{\phi}_i \tilde{\lambda}_i \tilde{\phi}_i^T$ , where  $\tilde{\phi}_i, \tilde{\lambda}_i$  are the eigenpairs of  $\tilde{A}$ .

# Matrix Factorization

Determine bounds for the change in the factors of a matrix when the matrix is perturbed.

## Theorem

*Stuart, 1977 If  $A = QR$  and  $A + \Delta A = (Q + \Delta Q)(R + \Delta R)$  are the QR factorizations, then, for sufficiently small  $\Delta A$*

$$\frac{\|\Delta R\|_F}{\|R\|_F} \leq c\kappa(A) \frac{\|\Delta A\|_F}{\|A\|_F}, \quad \|\Delta Q\|_F \leq c\kappa(A) \frac{\|\Delta A\|_F}{\|A\|_F} \quad (4)$$

*where  $c$  is a small constant and  $\kappa(A) = \|A\| \|A^{-1}\|$  is the condition number of  $A$ .*