

# Communication Networks (0368-3030) / Spring 2011

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A decorative graphic consisting of several horizontal lines of varying lengths and colors (teal, light blue, white) extending from the right side of the slide towards the center.

# Max-Min Fairness – Another Example

Recommended references:

- Computer Networks - Performance and Quality of Service \ Ivan Marsic (available online)
- An Engineering Approach to Computer Networking \ S.Keshav

# Generalization to a graph

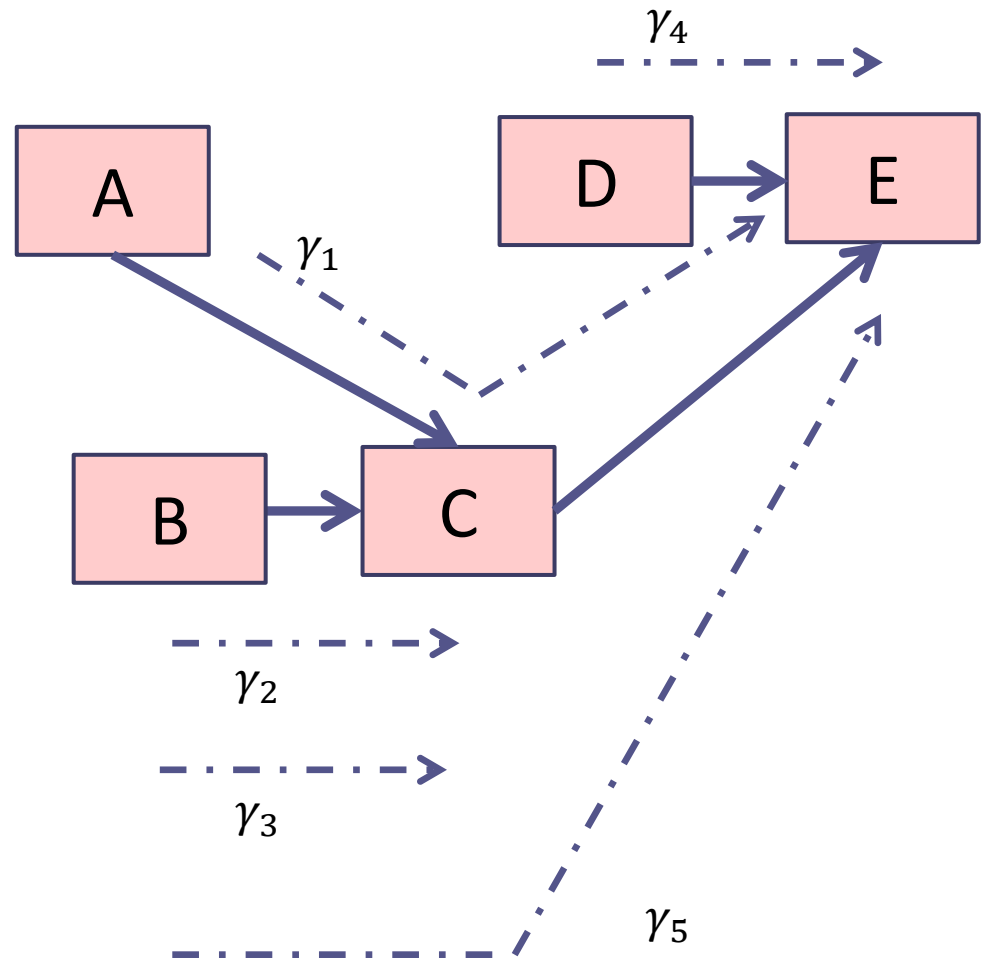
- We have a directed graph  $G = (V, E)$ , with capacity  $c_e$  for each edge  $e$ .
- We have a set of ongoing calls (flows). Each call  $i$  has demand  $r_i$  and a (fixed) path  $p_i$ .
- Algorithm:
  - Increase all flows equally until one link fills.
  - Fix the rate of the bottleneck flows.
  - Continue with the unfixed flows.

# Algorithm – in more detail

1. Assign flow 0 for all calls.
2. Let  $S$  be the set of all calls.
3. Increase the rate equally for all demands in  $S$  until:
  - a) some link is saturated  
– OR –
  - b) until some demand is fulfilled
4. Remove all the calls passing through the saturated links, and all the calls whose demand is fulfilled from  $S$ .
  - (these calls' rates become fixed – these rates will not change anymore)
5. Return to step 3, until there are no more calls left in  $S$ .

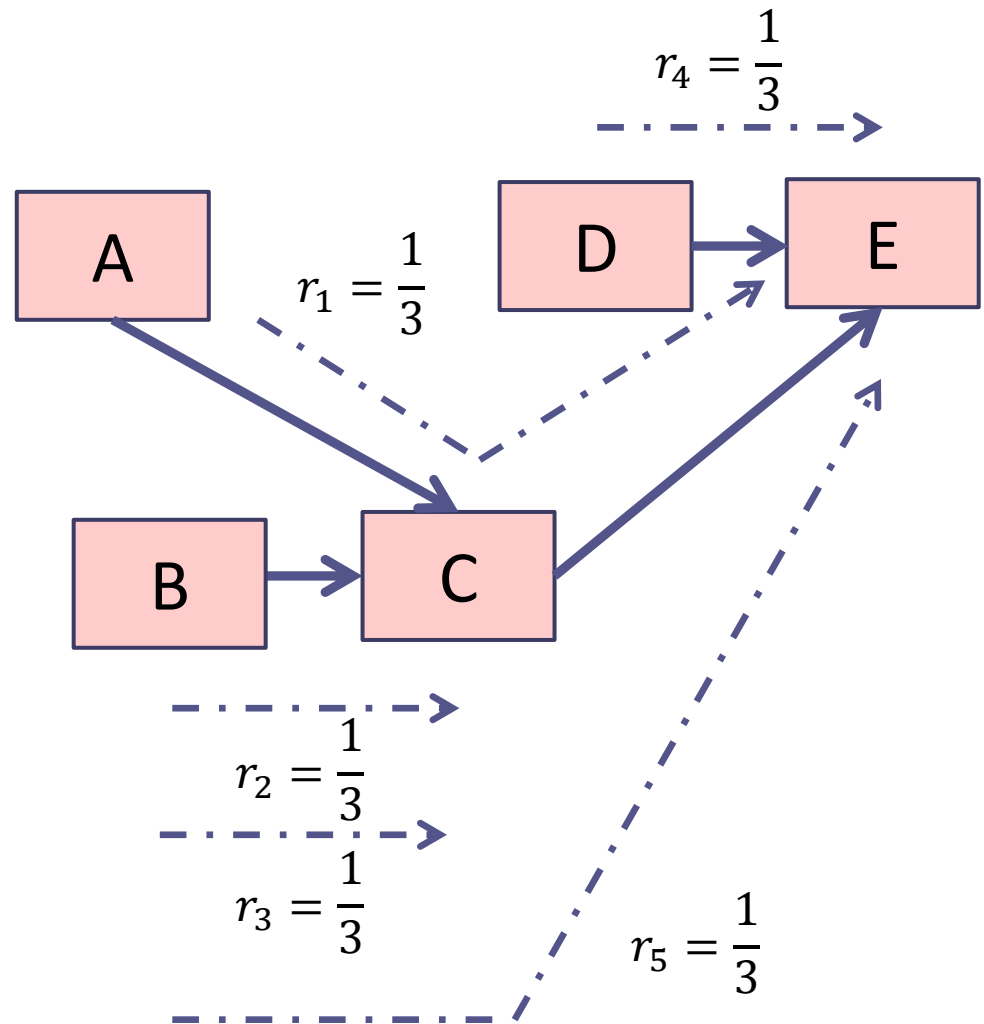
# An example

- A network with 5 nodes, and 5 calls  $\gamma_1, \dots, \gamma_5$ .
- All link capacities are 1.
- All demands are  $\infty$ .
- $S = \{\gamma_1, \dots, \gamma_5\}$ ,  
 $\forall i = 1, \dots, 5. r_i = 0$



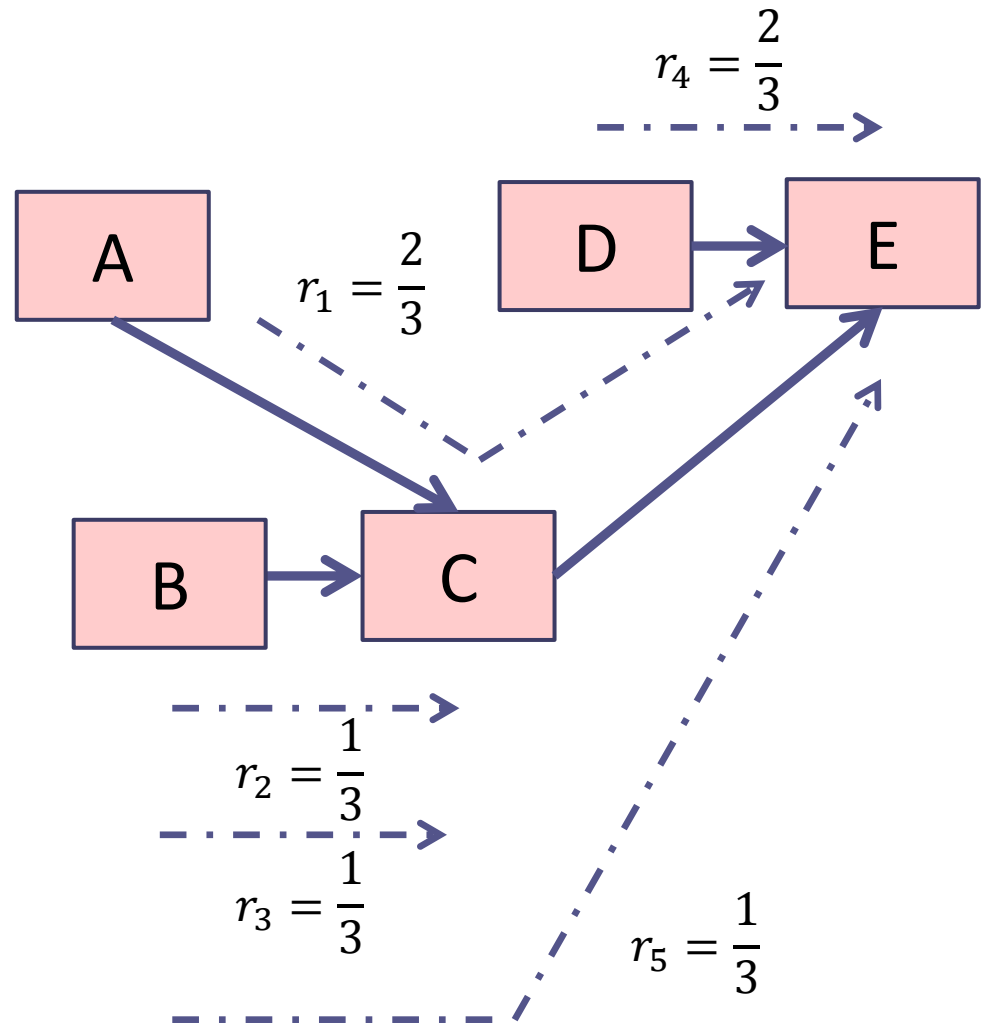
# An example

- Increase all flows in  $S$  equally.
- When all flows get rate  $1/3$  link  $(B, C)$  becomes saturated.
- $\gamma_2, \gamma_3, \gamma_5$  are removed from  $S$ , hence their rates will no longer change.
- $S = \{\gamma_1, \gamma_4\}$



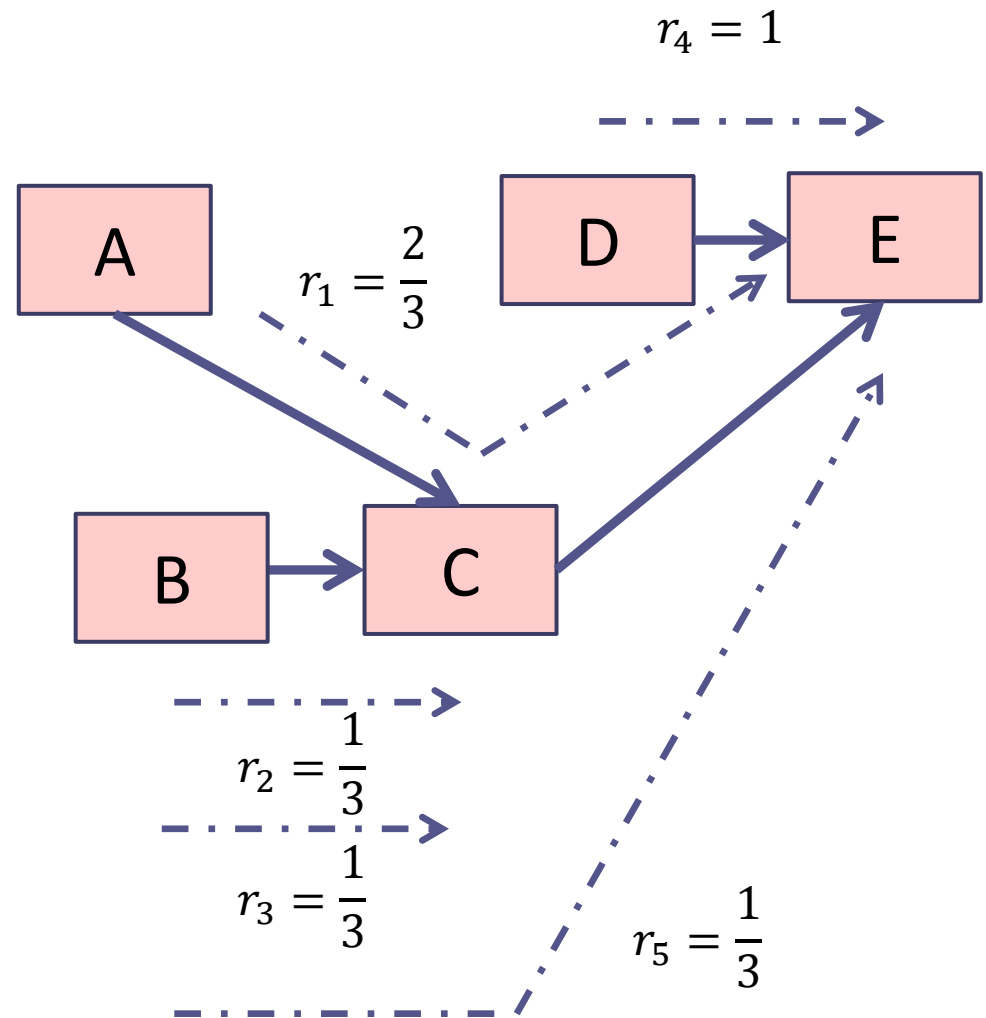
# An example

- Increase all flows in  $S$  equally.
- When you reach  $2/3$  link  $(C, E)$  becomes saturated.
- $S = \{\gamma_4\}$



# An example

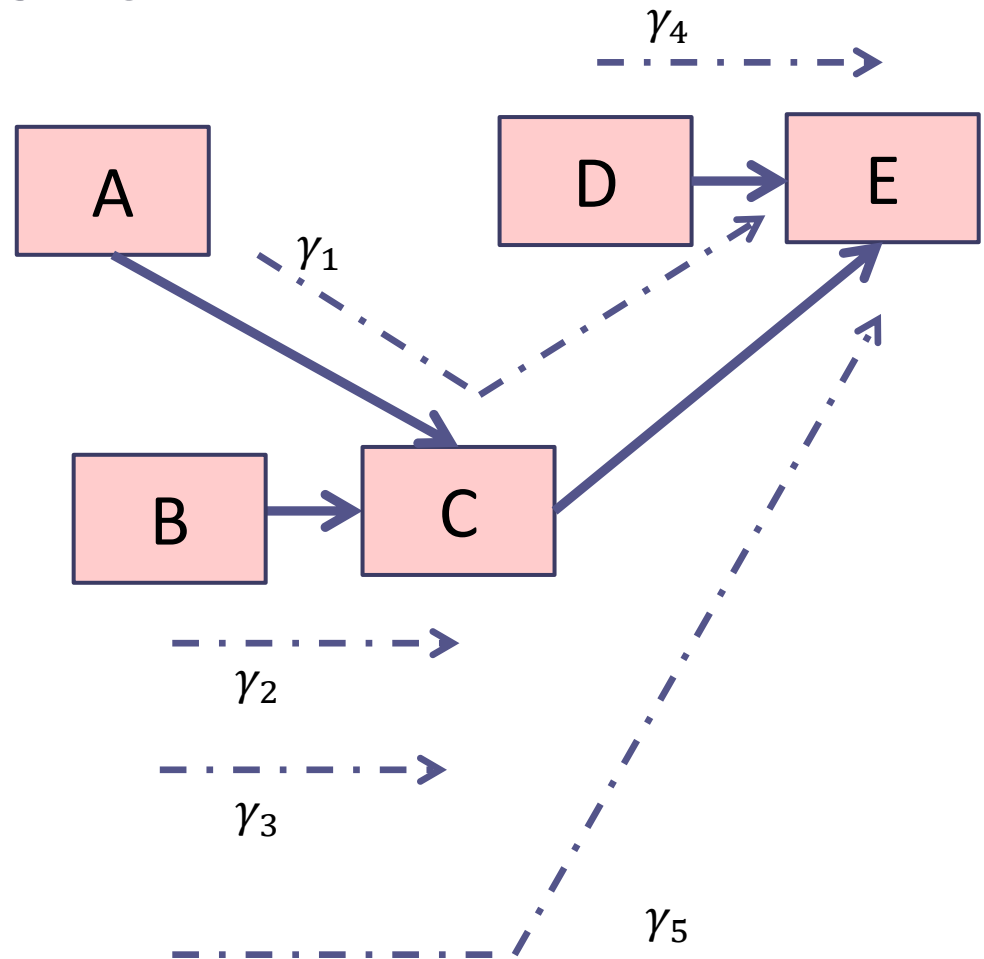
- Increase all flows in  $S$  equally.
- When you reach 1 link  $(D, E)$  becomes saturated
- $S = \Phi$





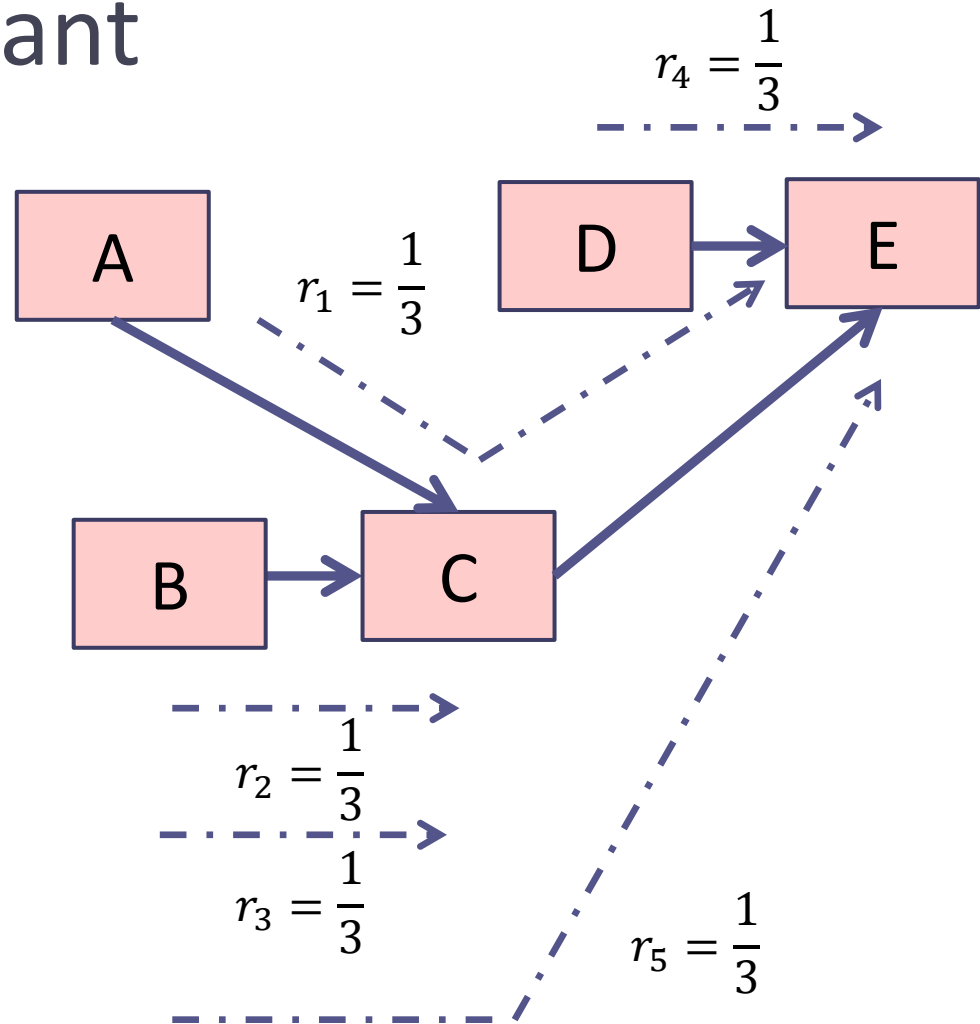
# An example - variant

- A network with 5 nodes, and 5 calls  $\gamma_1, \dots, \gamma_5$ .
- All link capacities are 1.
- All demands are 0.5.
- $S = \{\gamma_1, \dots, \gamma_5\}$ ,  
 $\forall i = 1, \dots, 5. r_i = 0$



# An example - variant

- Increase all flows in  $S$  equally.
- When all flows get rate  $1/3$  link  $(B, C)$  becomes saturated.
- $\gamma_2, \gamma_3, \gamma_5$  are removed from  $S$ , hence their rates will no longer change.
- $S = \{\gamma_1, \gamma_4\}$



# An example - variant

- Increase all flows in  $S$  equally.
- When you reach  $1/2$  the demands  $\gamma_1, \gamma_4$  are fulfilled.
- $\gamma_1, \gamma_4$  are removed from  $S$ .
- $S = \Phi$

