Compiler Construction

Intermediate Representation II

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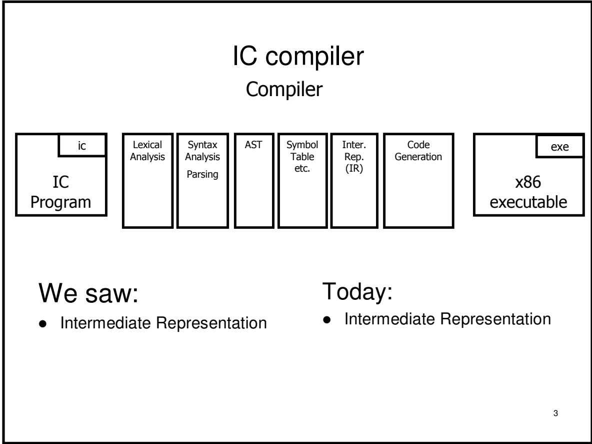
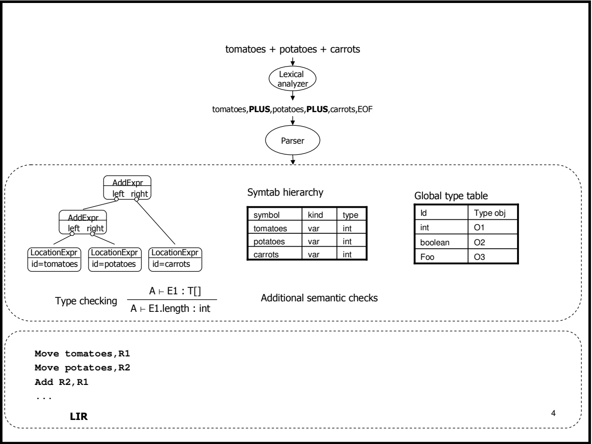
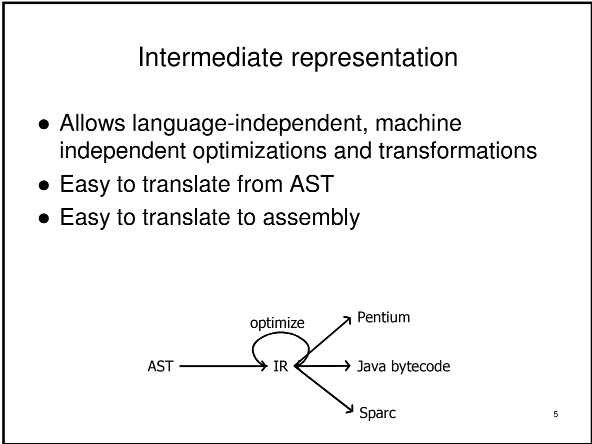
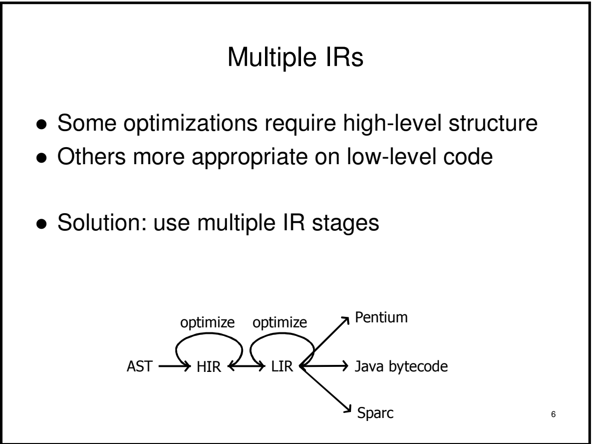
Administration

- PA3 submission deadline was extended to Dec 30
- Submit TA1 in Paz Grimberg's box by March 20









IR

- HIR = AST
- LIR
  - An abstract machine language
    - Not specific to a particular machine
  - Low-level language constructs
    - No looping structures, only jumps/conditional jumps
  - Two-operand instructions
  - OP a b
    - b = b OP a

LIR instructions

Instruction	Meaning
Move c,Rn	Rn = c
Move x,Rn	Rn = x
Move Rn,x	x = Rn
Add Rm,Rn	Rn = Rn + Rm
Sub Rm,Rn	Rn = Rn – Rm
Mul Rm,Rn	Rn = Rn * Rm
...	

Immediate (constant)

Memory (variable)

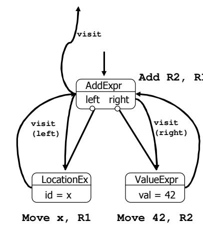
Note 1: rightmost operand = operation destination
Note 2: two register instr - second operand doubles as source and destination

Translating Expressions - Example

```
TR[x + 42]      Move x, R1
                Move 42, R2
                Add R2, R1
```


Translating expressions – example

```
TR[x + 42]
```



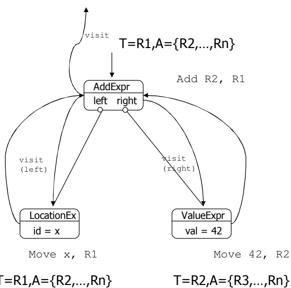
Move x, R1

Move 42, R2

Add R2, R1

Translating Expressions – Example (Cont'd)

```
TR[x + 42,T,A]
```



Move x, R1

Move 42, R2

Add R2, R1

Translating if-then-else

```
TR[if (e)          R1 := TR[e]
  then s1          Compare 0,R1
  else s2]         JumpTrue _false_label
                   TR[s1]
                   Jump _end_label
                   _false_label:
                   TR[s2]
                   _end_label:
```

Fresh labels generated during translation

Translating call/return

`TR[e1.f]`	`R1 := TR[e1]` `MoveField R1.`c_f`,R3`

`TR[C.foo(e1,…,en)]`	`R1 := TR[e1]`
	`…`
	`Rn := TR[en]`
	`StaticCall C_foo(x1=R1,…,xn=Rn),R`

`TR[e1.foo(e2)]`	`R1 := TR[e1]`
	`R2 := TR[e2]`
	`VirtualCall R1.`c_{foo}`(x=R2),R`

Runtime organization

- Representation of basic types
- Representation of allocated objects
  - Class instances
    - Dispatch vectors
  - Strings
  - Arrays
- Procedures
- Runtime checks

Representing basic types in IC

- int, boolean, string
  - Simplified representation: 32 bit for all types
  - **`boolean`** type could be implemented with single byte
- Arithmetic operations
  - Addition, subtraction, multiplication, division, remainder
- Mapped directly to target language types and operations
  - Exception: string concatenation implemented using library function **`__stringCat`**

Pointer types

- Represent addresses of source language data structures
- Usually implemented as an unsigned integer (4 bytes)
- Pointer dereferencing – retrieves pointed value
- May produce an error
  - Null pointer dereference

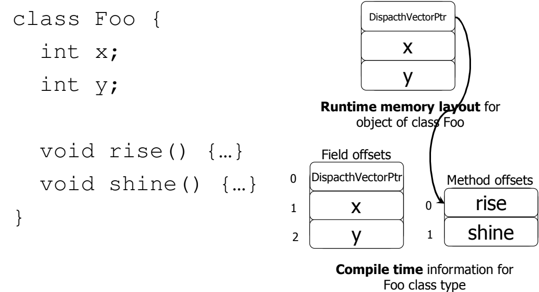
Object types

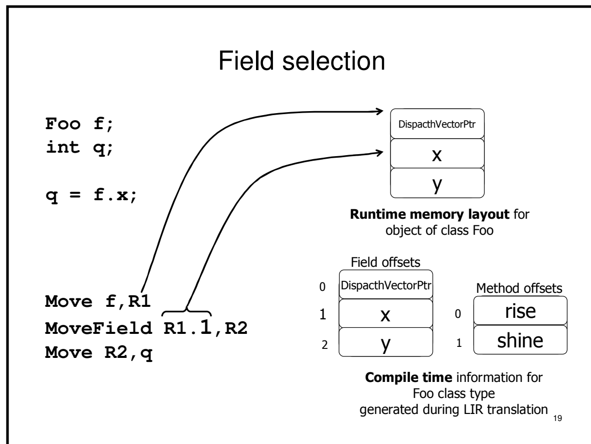
- Basic operations
  - Field selection
    - computing address of field, dereferencing address
  - Method invocation
    - Identifying method to be called, calling it

Object types

```
class Foo {
  int x;
  int y;

  void rise() {…}
  void shine() {…}
}
```





Implementation

- Store map for each `ClassType`
  - From field to offset
    - Note that 0 reserved for DispatchVectorPtr
  - From method to offset

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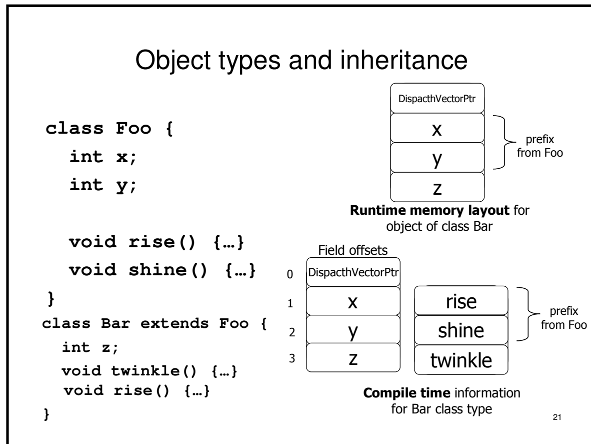
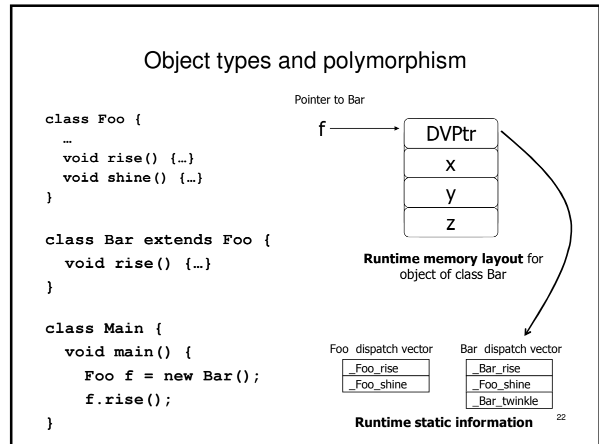
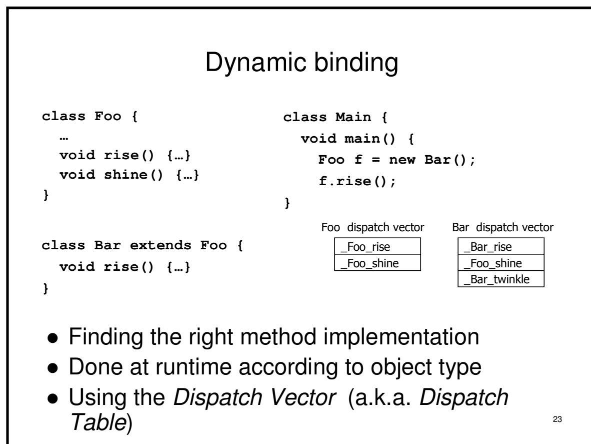
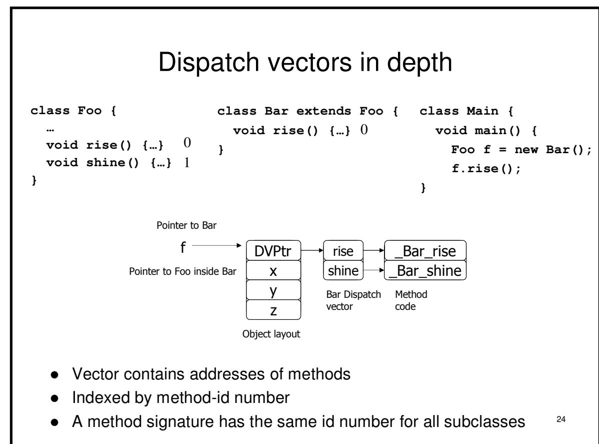


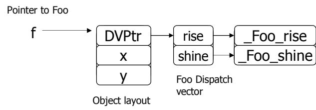






Dispatch vectors in depth

```
class Foo {
  …
  void rise() {…}  0
  void shine() {…} 1
}

class Bar extends Foo {
  void rise() {…} 0
}

class Main {
  void main() {
    Foo f = new Foo();
    f.rise();
}
```



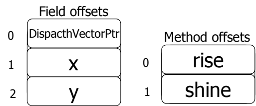
Object creation

```
Foo f = new Bar();
```

|Bar| = |x|+|y|+|z|+|DVPtr| =1+1+1 +1= 4 (16 bytes)

```
Library __allocateObject(16),R1
MoveField _Bar_DV,R1.0
Move R1,f
```

Label generated for class type `Bar` during LIR translation



Compile time information for Foo class type generated during LIR translation

LIR translation example

```
class A {
  int x;
  string s;
  int foo(int y) {
    int z=y+1;
    return z;
  }
  static void main(string[] args) {
    A p = new B();
    p.foo(5);
  }
}

class B extends A {
  int z;
  int foo(int y) {
    s = "y=" + Library.itos(y);
    Library.println(s);
    int[] sarr = Library.stoa(s);
    int l = sarr.length;
    Library.printi(l);
    return l;
  }
}
```

Translating virtual functions (dispatch tables)

Translating the main function

Translating virtual function calls

Translation for literal strings

Translating `.length` operator

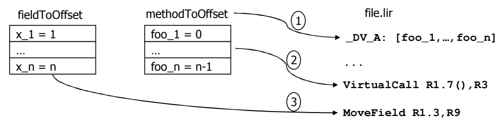
LIR program (manual trans.)

```
str1: "y="                               # Literal string in program
_DV_A: [_A_foo]                          # dispatch table for class A
_DV_B: [_B_foo]                          # dispatch table for class B

_A_foo:                                  # int foo(int y)
Move y,R1                                # int z=y+1;
Add 1,R1
Move R1,z
Return z                                 # return z;

_B_foo:                                  # int foo(int y)
Library __itos(y),R1                     # Library.itos(y);
Library __stringCat(str1,R1),R2          # "y=" + Library.itos(y);
Move this,R3                             # this.s = "y=" + Library.itos(y);
MoveField R2,R3.2
MoveField R3.2,R4
Library __println(R4),Rdummy             # Library.println(s);
Library __stoa(R4),R5                    # int[] sarr = Library.stoa(s);
Move R5,sarr
ArrayLength sarr,R6                      # int l = sarr.length;
Move R6,l
Library __printi(l),Rdummy               # Library.printi(l)
Return l                                 # return l;

# main in A
_ic_main:                                # A {static void main(string[] args)…}
Library __allocateObject(16),R1          # A p = new B();
MoveField _DV_B,R1.0                     # Update DVPtr of new object
VirtualCall R1.0(y=5),Rdummy             # p.foo(5)
```


Class layout implementation

```
class A {
  int x_1;
  ...
  boolean x_n;
  void foo_1(…) {…}
  ...
  int foo_n(…) {…}
}
```

```
class ClassLayout {
  Map<Method,Integer> methodToOffset;
  // DVPtr = 0
  Map<Field,Integer> fieldToOffset;
}
```



LIR optimizations

- Aim to reduce number of LIR registers and number of instructions
- Avoid storing variables and constants in registers
- Use accumulator registers
- Reuse "dead" registers
- Weighted register allocation

Avoid storing constants and variables in registers

- Don't allocate target register for each instruction
  - TR[5] = Move 5,Rj
  - TR[x] = Move x,Rk
- For a constant TR[5] = 5
- For a variable TR[x] = x
- TR[x+5] = Move 5,R1
  Add x,R1
  - Assign to register if **both operands** non-registers

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Accumulator registers

- Use same register for sub-expression and result

```
TR[e1 OP e2]

R1 := TR[e1]
R2 := TR[e2]
R1 := R1 OP R2
```

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Accumulator registers

```
TR[e1 OP e2]
a+(b*c)
```

```
R1 := TR[e1]
R2 := TR[e2]
R3 := R1 OP R2

Move a,R1
Move b,R2
Mul R1,R2
Move R2,R3
Move c,R4
Add R3,R4
Move R4,R5
```

```
R1 := TR[e1]
R2 := TR[e2]
R1 := R1 OP R2

Move b,R1
Mul c,R1
Add a,R1
```

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Accumulator registers cont.

- For instruction with N registers dedicate one register for accumulation
- Accumulating instructions, use:
  - **MoveArray R1[R2],R1**
  - **MoveField R1.7,R1**
  - **StaticCall _foo(R1,…),R1**
  - …

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Reuse registers

- Registers have very-limited lifetime
- TR[e1 OP e2] = R1:=TR[e1]
  R2:=TR[e2]
  R1:=R1 OP R2

  Registers from TR[e1] can be reused in TR[e2]
- Solution:
  - Use a stack of LIR registers
  - Stack corresponds to recursive invocations of t := TR[e]
  - All the temporaries on the stack are alive

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Weighted register allocation

- Sethi & Ullman algorithm
  - Two expression e1, e2 and an operation OP
  - e1,e2 without side-effects
    - function calls
  - TR[e1 OP e2] = TR[e2 OP e1]
  - Weighted register allocation
    - translate heavier sub-tree first

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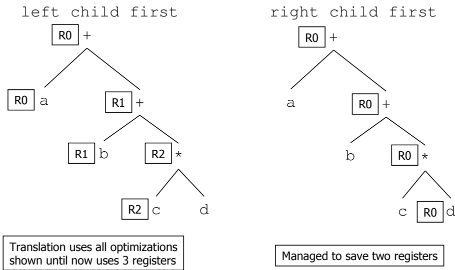
Example

R0 := TR[a+(b+(c*d))]



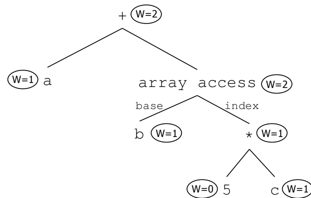
Weighted register allocation

- Can save registers by re-ordering subtree computations
- Label each node with its weight
  - Weight = number of registers needed
  - Leaf weight known
  - Internal node weight
    - w(left) > w(right) then w = left
    - w(right) > w(left) then w = right
    - w(right) = w(left) then w = left + 1
- Choose heavier child as first to be translated
- Have to check that no side-effects exist

Weighted reg. alloc. example

R0 := TR[a+b[5*c]]

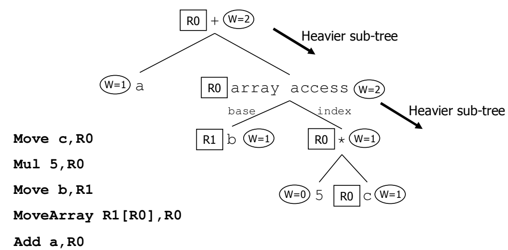
Phase 1: - check absence of side-effects in expression tree
- assign weight to each AST node



Weighted reg. alloc. example

R0 := TR[a+b[5*c]]

Phase 2: use weights to decide on order of translation




```
Move c,R0
Mul 5,R0
Move b,R1
MoveArray R1[R0],R0
Add a,R0
```


PA4

- Translate AST to LIR (file.ic -> file.lir)
  - Dispatch table for each class
  - Literal strings (all literal strings in file.ic)
  - Instruction list for every function
    - Leading label for each function _CLASS_FUNC
  - Label of main function should be **`_ic_main`**
- Maintain internally for each function
  - List of LIR instructions
  - Reference to method AST node
    - Needed to generate frame information in PA5
- Maintain for each call instruction
  - Reference to method AST
    - Needed to generate call sequence in PA5
- Optimizations
  - After assignment works
  - Keep optimized and non-optimized translations separately

Tips for PA4

- Keep **list** of LIR instructions for each translated method
- Keep **`ClassLayout`** information for each class
  - Field offsets
  - Method offsets
  - Don't forget to take superclass fields and methods into account
- May be useful to keep reference in each LIR instruction to AST node from which it was generated
- Two AST passes:
  - Pass 1:
    - Collect and name strings literals (**`Map<ASTStringLiteral,String>`**)
    - Create **`ClassLayout`** for each class
  - Pass 2: use literals and field/method offsets to translate method bodies
- Finally: print string literals, dispatch tables, print translation list of each method body

microLIR simulator

- Written by Roman Manevich
- Java application
  - Accepts file.lir (your translation)
  - Executes program
- Use it to test your translation
  - Checks correct syntax
  - Performs lightweight semantic checks
  - Runtime semantic checks
  - Debug modes (-verbose:1|2)
  - Prints program statistics (#registers, #labels, etc.)
- Comes with sample inputs
- Read manual

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