

# A quantitative Lovász criterion for Property B

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## Abstract

A well known observation of Lovász is that if a hypergraph is not 2-colorable, then at least one pair of its edges intersect at a single vertex. In this short paper we consider the quantitative version of Lovász's criterion. That is, we ask how many pairs of edges intersecting at a single vertex, should belong to a non 2-colorable  $n$ -uniform hypergraph? Our main result is an *exact* answer to this question, which further characterizes all the extremal hypergraphs. The proof combines Bollobás's two families theorem with Pluhar's randomized coloring algorithm.

## 1 Introduction

A hypergraph  $\mathcal{H} = (V, E)$  consists of a vertex set  $V$  and a set of edges  $E$  where each  $X \in E$  is a subset of  $V$ . If all edges of  $\mathcal{H}$  have size  $n$  then  $\mathcal{H}$  is called an  $n$ -uniform hypergraph, or  $n$ -graph for short. A hypergraph is 2-colorable if one can assign each vertex  $v \in V$  one of two colors, say *Red/Blue*, so that each  $X \in E$  contains vertices of both colors. Miller [6], and later Erdős in various papers, referred to this property as *Property B*, after F. Bernstein [2] who introduced it in 1907. Since deciding if a hypergraph is 2-colorable is *NP*-hard one cannot hope to find a simple characterization of all 2-colorable hypergraphs. Instead, one looks for general sufficient/necessary conditions for having this property. For example, a famous result of Seymour [8] states that if  $\mathcal{H}$  is not 2-colorable then  $|E| \geq |V|$ . Probably the most well studied question of this type asks for the smallest number of edges in an  $n$ -graph that is not 2-colorable. The study of this quantity, denote  $m(n)$ , was popularized by Erdős, see [1] for a comprehensive treatment. Despite much effort by many researchers, even the asymptotic value of  $m(n)$  has not been determined yet.

A pair of edges  $X, Y \in E(\mathcal{H})$  is *simple* if  $|X \cap Y| = 1$ . Let  $m_2(\mathcal{H})$  denote the number of ordered simple pairs of edges of  $\mathcal{H}$ . A well known observation of Lovász [5] states that if  $\mathcal{H}$  is not 2-colorable then  $m_2(\mathcal{H}) > 0$ . Despite its simplicity, this observation underlies the best known bounds for  $m(n)$ , see [4, 7]. It is natural to ask if one can obtain a quantitative version of Lovász's observation, that is, estimate how small can  $m_2(\mathcal{H})$  be in an  $n$ -graph not satisfying property *B*? Our main result in this paper states that (somewhat surprisingly), one can give an exact answer to the above extremal question as well as characterize the extremal  $n$ -graphs.

Let  $K_{2n-1}^n$  denote the complete  $n$ -graph on  $2n-1$  vertices. It is easy to see that  $K_{2n-1}^n$  is not 2-colorable and that  $m_2(K_{2n-1}^n) = n \cdot \binom{2n-1}{n}$ . We first observe that this simple upper bound is tight.

**Proposition 1.1.** *If an  $n$ -graph is not 2-colorable then  $m_2(\mathcal{H}) \geq n \cdot \binom{2n-1}{n}$ .*

As with any extremal problem, one would like to know which graphs or hypergraphs are extremal with respect to this problem. For example, Turán's theorem states that among all  $n$ -vertex graphs not containing a complete  $t$ -vertex subgraph, there is only one graph maximizing the number of edges. In the setting of our

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problem, it is easy to see that  $K_{2n-1}^n$  is not the only non 2-colorable  $n$ -graph satisfying  $m_2(\mathcal{H}) = n \cdot \binom{2n-1}{n}$ , since one can take a copy of  $K_{2n-1}^n$  and add to it more vertices and edges without increasing the number of simple pairs. Our main result in this paper characterizes the extremal  $n$ -graphs, by showing that this is in fact the only way to construct an  $n$ -graph meeting the bound of Proposition 1.1.

**Theorem 1.** *If a non 2-colorable  $n$ -graph  $\mathcal{H}$  satisfies  $m_2(\mathcal{H}) = n \cdot \binom{2n-1}{n}$  then it contains a copy of  $K_{2n-1}^n$ .*

While the proof of Proposition 1.1 is implicit in Pluhar's [7] argument for bounding  $m(n)$ , the proof of Theorem 1 is more intricate, relying on Bollobás's two families theorem [3] as well as on a refined analysis of Pluhar's randomized algorithm for 2-coloring  $n$ -graphs.

## 2 Proof of Proposition 1.1

In this section we describe several preliminary observations regarding a coloring algorithm introduced in [7], and use them to derive Proposition 1.1. The algorithm is the following:

**Algorithm**  $\text{Col}(\mathcal{H}, \pi)$ . The input is a hypergraph  $\mathcal{H} = (V, E)$  and an ordering  $\pi : V \mapsto \{1, \dots, |V|\}$  (that is,  $\pi$  is a bijection). The output is a 2-coloring of  $V$  (not necessarily a proper one). The algorithm runs in  $|V|$  steps, where in each time step  $1 \leq i \leq |V|$ , the vertex  $\pi^{-1}(i)$  is being colored *Blue* if this does not form any monochromatic *Blue* edge. Otherwise,  $\pi^{-1}(i)$  is colored *Red*.

We now state an important property of  $\text{Col}(\mathcal{H}, \pi)$ . For two disjoint subsets  $X, Y \subseteq V$ , we use the notation  $\pi(X) < \pi(Y)$  whenever  $\max_{x \in X} \pi(x) < \min_{y \in Y} \pi(y)$ , that is, the elements of  $X$  precede all the elements of  $Y$  in the ordering  $\pi$ . Suppose  $(X, Y)$  is a simple pair of edges in  $\mathcal{H}$  with<sup>1</sup>  $X \cap Y = y$ . We say that  $\pi$  *separates*  $(X, Y)$  if  $\pi(X \setminus y) < \pi(y) < \pi(Y \setminus y)$ .

**Claim 2.1.** *If  $\text{Col}(\mathcal{H}, \pi)$  fails to properly color  $\mathcal{H}$  then  $\pi$  separates at least one pair of simple edges.*

*Proof.* We first observe that (by definition) for every ordering  $\pi$ , the algorithm  $\text{Col}(\mathcal{H}, \pi)$  does not produce monochromatic *Blue* edges. Suppose then it produced a *Red* edge  $Y \in E$ . Let  $y$  be the first vertex of  $Y$  according to the ordering  $\pi$ . If  $y$  was colored red, then there must have been an edge  $X$  so that  $y \in X$ , and all other vertices of  $X$  were already colored *Blue* (otherwise the algorithm would color  $y$  *Blue*). This means  $(X, Y)$  is simple and that  $\pi$  separates it.  $\square$

Note that the claim above already shows that if  $\mathcal{H}$  is not 2-colorable then  $m_2(\mathcal{H}) > 0$ . For the proof of Proposition 1.1 we will also need the following simple fact.

**Claim 2.2.** *A random permutation separates any given simple pair with probability  $1/n \binom{2n-1}{n}$ .*

*Proof.* Let  $(X, Y)$  be a simple pair, and let  $X \cap Y = y$ . A permutation  $\pi$  separates  $(X, Y)$  if and only if  $\pi(X \setminus y) < \pi(y) < \pi(Y \setminus y)$ , and this happens with probability exactly

$$\frac{(n-1)!(n-1)!}{(2n-1)!} = \frac{1}{n \binom{2n-1}{n}}$$

as desired.  $\square$

The above claims suffice for proving Proposition 1.1.

*Proof (of Proposition 1.1):* Assume  $m_2(\mathcal{H}) < n \binom{2n-1}{n}$ . Suppose we pick a uniformly random  $\pi$ . Then by the union bound and Claim 2.2, we infer that with positive probability  $\pi$  does not separate any simple pair edges. Hence, there is a  $\pi$  not separating any simple pair. Claim 2.1 then implies that  $\text{Col}(\mathcal{H}, \pi)$  will produce a legal 2-coloring of  $\mathcal{H}$ .  $\square$

<sup>1</sup>Here, and in what follows, we slightly abuse notation by writing  $y$  instead of the more appropriate  $\{y\}$ .

### 3 Proof of Theorem 1

For the rest of this section fix some non 2-colorable  $n$ -graph  $\mathcal{H} = (V, E)$  satisfying  $m_2(\mathcal{H}) = n \binom{2n-1}{n}$ . We need to show that  $\mathcal{H}$  contains a copy of  $K_{2n-1}^n$ . We start with a few preliminary claims regarding  $\mathcal{H}$ .

First, we show that no  $\pi$  separates more than one simple pair.

**Claim 3.1.** *Every ordering  $\pi$  separates at most one simple pair.*

*Proof.* Suppose  $\pi$  separates two simple pairs. By Claim 2.2, the assumption on  $m_2(\mathcal{H})$ , and by linearity of expectation, the expected number of simple pairs separated by a random permutation is exactly 1. Hence, if  $\pi$  separates 2 simple pairs, then there must exist a permutation  $\sigma$  which separates less than 1, and therefore 0, simple pairs. Therefore, by Claim 2.1 we obtain that  $\text{Col}(\mathcal{H}, \sigma)$  produces a legal 2-coloring of  $\mathcal{H}$ , which is a contradiction to the assumption that  $\mathcal{H}$  is not 2-colorable.  $\square$

**Claim 3.2.** *If  $(X, Y)$  and  $(X', Y)$  are simple pairs, then  $X \cap Y \neq X' \cap Y$ .*

*Proof.* We observe that if  $X \cap Y = X' \cap Y = y$ , then there is a  $\pi$  that separates both  $(X, Y)$  and  $(X', Y)$ , and this will contradict Claim 3.1. Indeed, if  $(X, Y)$  and  $(X', Y)$  are simple pairs and  $X \cap Y = X' \cap Y = y$ , then  $(X \cup X') \setminus y$  and  $Y$  are disjoint. Therefore, any  $\pi$  satisfying

$$\pi((X \cup X') \setminus y) < \pi(y) < \pi(Y \setminus y)$$

separates  $(X, Y)$  and  $(X', Y)$ . This completes the proof.  $\square$

In addition to the above observations about  $\mathcal{H}$ , the last ingredient we will need is the following theorem of Bollobás [3].

**Lemma 3.3.** *Let  $I$  be an index set. For all  $i \in I$ , let  $A_i$  and  $B_i$  be subsets of a set  $V$  of  $p$  elements satisfying the following conditions:*

- i.  $A_i \cap B_i = \emptyset$  for all  $i \in I$ , and*
- ii.  $A_j \not\subseteq A_i \cup B_i$  for all  $i \neq j \in I$ .*

*Then, we have*

$$\sum_{i \in I} \frac{1}{\binom{p-|B_i|}{|A_i|}} \leq 1,$$

*with equality if and only if  $B_i = B$  for all  $i \in I$  and the sets  $A_i$  are all the  $q$ -tuples of the set  $P \setminus B$  for some value of  $q$ .*

Let us now show how to use Lemma 3.3 in order to derive Theorem 1. Recall that  $V$  is the vertex set of  $\mathcal{H}$  and set  $p := |V|$ . Let  $M(\mathcal{H})$  be a collection of simple pairs  $(X, Y)$  defined as follows; out of all the simple pairs  $(X, Y)$  with the same ‘‘second’’ set  $Y$ , put in  $M(\mathcal{H})$  one of these pairs. Observe that by Claim 3.2 each  $Y$  belongs to at most  $|Y| = n$  simple pairs of the form  $(X, Y)$  (i.e, with  $Y$  as the second set), implying that  $t := |M(\mathcal{H})| \geq \frac{1}{n} \cdot m_2(\mathcal{H}) = \binom{2n-1}{n}$ . We now define a collection  $\mathcal{F}$  consisting of pairs of subsets of  $V$  as follows: For every simple pair  $s := (X, Y) \in M(\mathcal{H})$ , define  $A_s = X \setminus Y$  and  $B_s = V \setminus (X \cup Y)$ , and let  $\mathcal{F} = \{(A_s, B_s) : s \in M(\mathcal{H})\}$ . For convenience, let us rename the pairs in  $\mathcal{F}$  as  $(A_i, B_i)$  with  $1 \leq i \leq t$ .

Now we wish to show that  $\mathcal{F}$  satisfies the conditions in Lemma 3.3. Observe that if it does, then since

$$\sum_{i=1}^t \frac{1}{\binom{p-|B_i|}{|A_i|}} = \sum_{i=1}^t \frac{1}{\binom{2n-1}{n-1}} \geq 1,$$

it follows by the first part of Lemma 3.3 that the last inequality is in fact an equality. Therefore, by the second part of Lemma 3.3, we conclude that all the  $B_i$ 's are the same set  $B$ , and the set of all the  $A_i$ 's

consists of all  $n - 1$  subsets of a ground set of size  $2n - 1$ . That is, let  $B = B_i$  and  $U = V \setminus B$ . Then we have that  $|U| = 2n - 1$ , and that the sets  $A_i$  are all the  $n - 1$  subsets of  $U$ . Since by construction we have that  $U \setminus A_i \in E(\mathcal{H})$  for all  $i$ , we conclude that  $\mathcal{H}$  restricted to the set  $U$  is a copy of  $K_{2n-1}^n$  as desired. It thus remains to show the following:

**Claim 3.4.**  $\mathcal{F}$  satisfies the conditions in Lemma 3.3

*Proof.* The first condition  $A_i \cap B_i = \emptyset$  for all  $i$  is trivially satisfied by construction. For the second condition, let  $(A, B)$  and  $(A', B')$  be two elements in  $\mathcal{F}$  coming from simple pairs  $(X, Y)$  and  $(X', Y')$  belonging to  $M(\mathcal{H})$ , respectively. Recall that by the way we defined  $M(\mathcal{H})$  and  $\mathcal{F}$  we have  $Y \neq Y'$ . Let us use  $y$  and  $y'$  to denote the unique elements in  $X \cap Y$  and  $X' \cap Y'$ , respectively. We wish to show that  $A \not\subseteq A' \cup B'$ , which, by construction, is implied by  $(X \setminus y) \cap Y' \neq \emptyset$ . Assuming  $(X \setminus y) \cap Y' = \emptyset$ , we will derive a contradiction to Claim 3.1 by showing that there is a permutation  $\pi$  separating two distinct simple pairs.

Observe that it cannot be that  $y \in Y'$ . Indeed, if it was the case, then together with the assumption that  $(X \setminus y) \cap Y' = \emptyset$  we would infer that  $(X, Y)$  and  $(X, Y')$  are both simple pairs intersecting at  $y$  (and distinct as  $Y \neq Y'$ ), contradicting Claim 3.2. Assume then that  $y \notin Y'$  (so in particular  $y \neq y'$ ). We claim that we can find a  $\pi$  satisfying

$$\pi(X \setminus y) < \pi(y) < \pi((X' \setminus y') \setminus X) < \pi(y') < \pi((Y \cup Y') \setminus (X \cup X')).$$

Indeed, the only thing that needs to be justified is the ability to place  $y'$  as above, which follows from the fact that  $y' \in Y'$  and the assumption  $(X \setminus y) \cap Y' = \emptyset$  which together imply that  $y' \notin X$ . Observe that since  $\pi$  first places  $X \setminus y$  and then  $y$ , the pair  $(X, Y)$  is separated by  $\pi$ . Such a  $\pi$  clearly places  $X' \setminus y'$  before  $y'$  and the assumption  $(X \setminus y) \cap Y' = \emptyset$  together with the fact that  $y \notin Y'$  imply that such a  $\pi$  places all of  $Y' \setminus y'$  after  $y'$ , so it separates  $(X', Y')$  as well, giving us the desired contradiction.  $\square$

This completes the proof of Theorem 1.

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