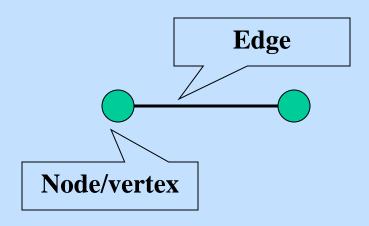
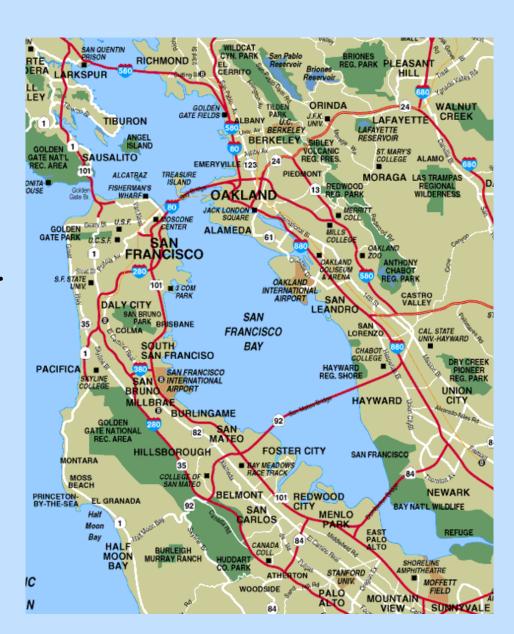
### Networks & Modules

- 1. Networks
- 2. Protein complex identification
- 3. Pathway identification

## Networks

- Represent relations between elements.
- *Nodes* elements (towns).
- *Edges* relations (roads).





## It's a Small World

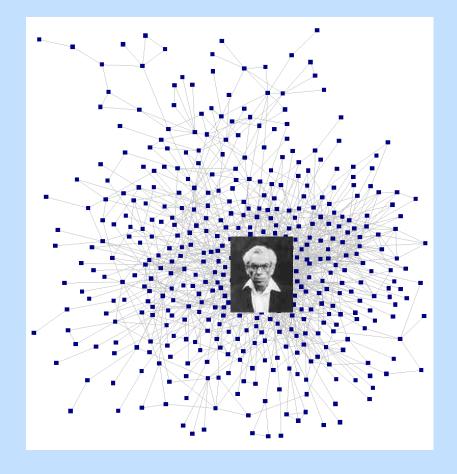


Milgram'67: six degrees of separation.

## Collaboration Networks

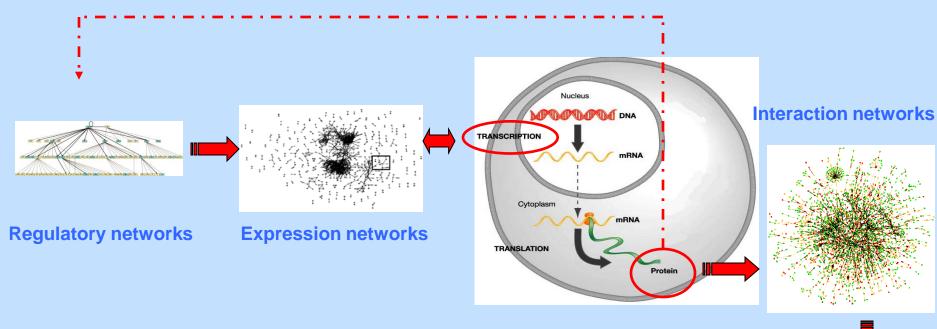
Nodes – collaborators
(scientists)

Edges – acts of collaboration
(joint articles)



http://www.orgnet.com/Erdos.html

## Molecular Networks



*Nodes* – molecules

*Edges* – interactions / similarity

Metabolic networks

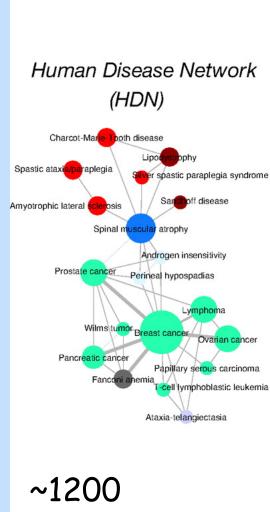


#### The human disease network

Kwang-Il Goh\*<sup>†‡§</sup>, Michael E. Cusick<sup>†‡¶</sup>, David Valle<sup>||</sup>, Barton Childs<sup>|</sup>, Marc Vidal<sup>†‡¶</sup>\*\*, and Albert-László Barabási\*<sup>†‡</sup>\*\*

DISEASOME

Fanconi anemia



diseases

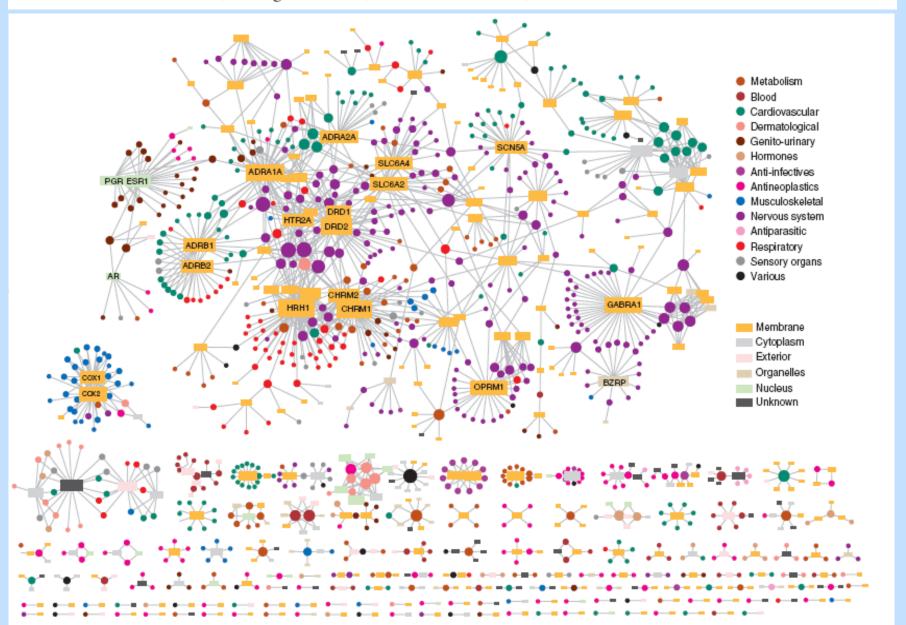
#### disease phenome disease genome Ataxia-telangiectasia Perineal hypospadias ATM Androgen insensitivity T-cell lymphoblastic leukemia BRCA1 Papillary serous carcinoma BRCA2 Prostate cancer CDH1 Ovarian cancer GARS HEXB Lymphoma KRAS Breast cancer LMNA MSH2 Pancreatic cancer PIK3CA Wilms tumor TP53 Spinal muscular atrophy MAD1L1 Sandhalldisease RAD54L VAPB Charcot-Marie-Tooth disease CHEK2 Amyotrophic lateral sclerosis BSCL2 Silver spastic paraplegia syndrome Spastic ataxia/paraplegia BRIP1

Disease Gene Network (DGN) LMNA BSCL2 VAPB GARS ATM BRCA2 BRIP1 BRCA1 KRAS RAD54L TP53 MAD1L1 CHEK2 PIK3CA CDH1 ~1800

genes

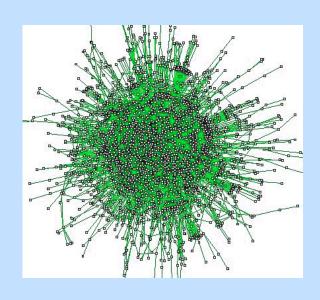
#### Drug-target network

Muhammed A Yıldırım<sup>1,2,3</sup>, Kwang-Il Goh<sup>1,4,5</sup>, Michael E Cusick<sup>1,2</sup>, Albert-László Barabási<sup>1,4,6</sup> & Marc Vidal<sup>1,2</sup>

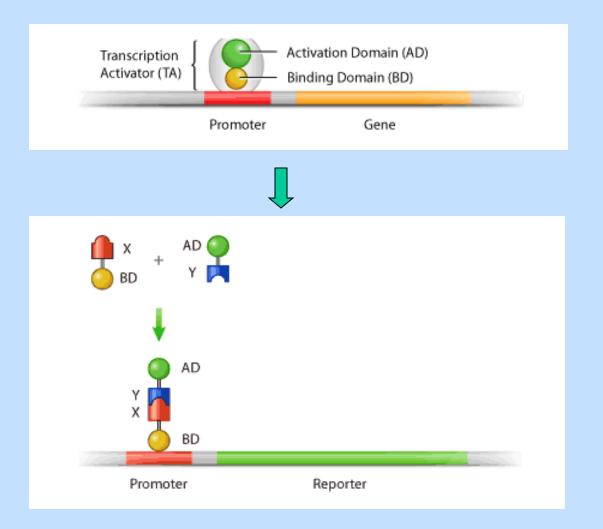


# Protein-protein Interaction Networks

- *Nodes* proteins (6K).
- Edges interactions (40K).
- Reflect the cell's machinery and signlaing pathways.
- Measured by high-throughput technologies:
  - yeast two-hybrid
  - co-immunoprecipitation



## Yeast Two-Hybrid

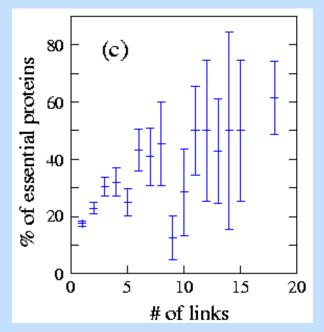


http://www.bioteach.ubc.ca/MolecularBiology/AYeastTwoHybridAssay/

# Network properties: Degree

# Why is degree important?

- Degree: #neighbors.
- Local characterization of a node.
- Indicates its centrality in the network.

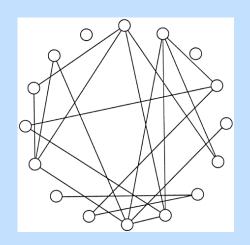


## Degree Distribution

- Degree distribution P(k): probability that a node has degree k.
- For directed graphs, two distributions: in-degree and out-degree.
- Average degree:  $d = \sum_{k>0} kP(k)$
- Number of edges: Nd/2.

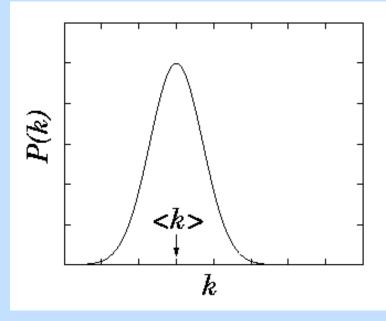
# Random Networks (Erdös/Rényi)

- N nodes.
- Every pair of nodes is connected with probability *p*.



- Mean degree:  $d=(N-1)p\sim Np$ .
- Degree distribution is binomial, asymptotically Poisson:

$$P(k) = \frac{e^{-d}d^k}{k!}$$

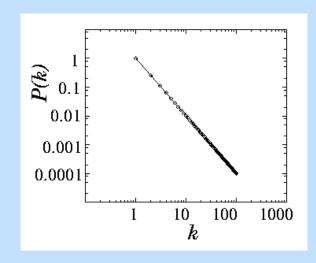


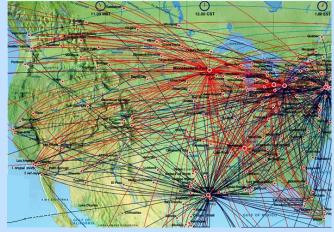
# Scale-Free Networks (Barabasi&Albert'99)

Power-law degree distribution

$$P(k) \propto k^{-c}, k \neq 0, c > 1$$

• Characterized by a small number of highly connected nodes, known as *hubs*.





### Scale-Free Distribution

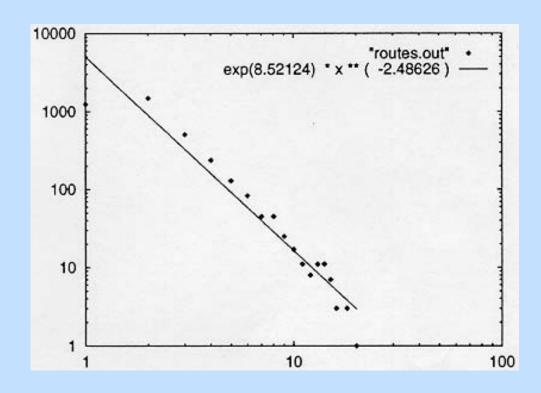
- A distribution p(x) that is scale-invariant, i.e.: p(ax)=g(a)p(x)
- It can be shown that the only scale free distributions are power-law distributions!!!

# Are Real Networks Random or Scale-Free?

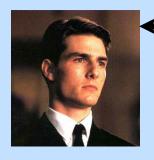
### The Internet

*Nodes* – routers. *Edges* – physical links.

 $P(k) \sim k^{-2.5}$  (Faloutsos et al.'99)



## Film Actors

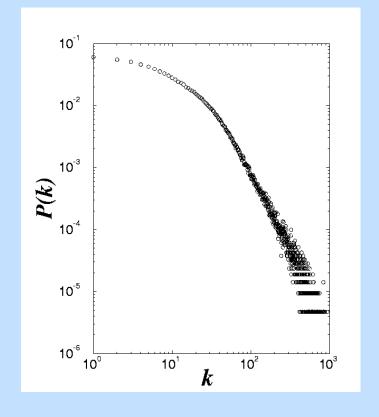


Days of Thunder (1990) Far and Away (1992) Eyes Wide Shut (1999)



Nodes – actors. Edges – joint movies.

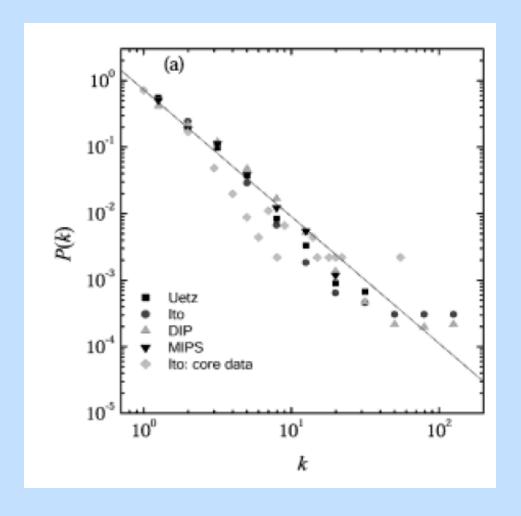
 $P(k) \sim k^{-2.3}$  (Barabasi&Albert'99)



## Protein Interaction Networks

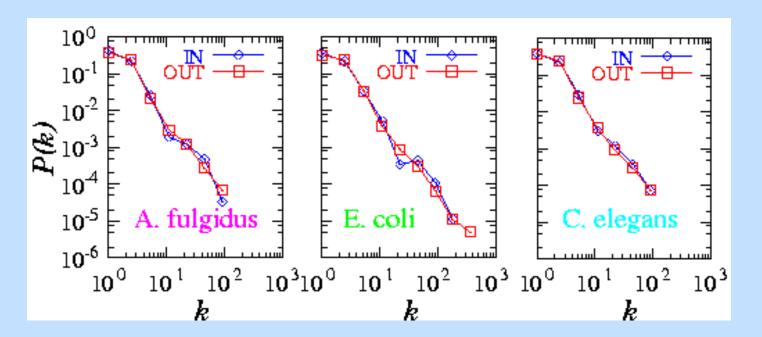
- *Nodes* proteins.
- *Edges* interactions.

 $P(k) \sim k^{-2.5}$  (Yook et al. '04)



### Metabolic Networks

- *Nodes* metabolites.
- Edges biochemichal reactions.



Metabolic networks from all kingdoms of life are scale-free  $c=2.2\pm0.2$  (Jeong et al.'00)

# Why Are Real Networks Scale-Free?

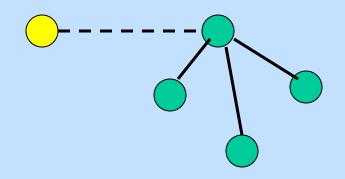
# Scale-Free Model (Barábasi & Albert)

- Growth: nodes are constantly added.
- **Preferential attachment:** the probability that a new node connects to existing ones is proportional to their degree.

In resulting network:

$$P(k) \approx k^{-3}$$

Relevance to biology?



# Clustering

# Clustering Coefficient (Watts & Strogatz)

• Characterizes tendency of nodes to cluster

$$C(v) = \frac{\#\{\text{pairs of connected neighbors of } v\}}{d(v)(d(v)-1)/2}$$

$$C = \frac{1}{N} \sum_{v} C(v)$$

(if d(v)=0,1 then C(v) is defined to be 0)

• Lies in [0,1].

#### • What is C for random graphs?

Table 1: Clustering coefficients, C, for a number of different networks; n is the number of node, z is the mean degree. Taken from [146].

Network	n	z	C	C for
			measured	random graph
Internet [153]	6,374	3.8	0.24	0.00060
World Wide Web (sites) [2]	153,127	35.2	0.11	0.00023
power grid [192]	4,941	2.7	0.080	0.00054
biology collaborations [140]	1,520,251	15.5	0.081	0.000010
mathematics collaborations [141]	253,339	3.9	0.15	0.000015
film actor collaborations [149]	449,913	113.4	0.20	0.00025
company directors [149]	7,673	14.4	0.59	0.0019
word co-occurrence [90]	460,902	70.1	0.44	0.00015
neural network [192]	282	14.0	0.28	0.049
metabolic network [69]	315	28.3	0.59	0.090
food web [138]	134	8.7	0.22	0.065

## Shortest Paths

### Small World

- What is the avg. distance in a random network?
- Fact: a random network is locally tree-like (exponential growth of neighbors with distance)
- $d^i$  vertices on avg. are at distance i or closer from a vertex.
- Since  $N\sim d^l$  we have  $l\sim \ln N/\ln d small$  world effect.
- Implies fast spread of information.

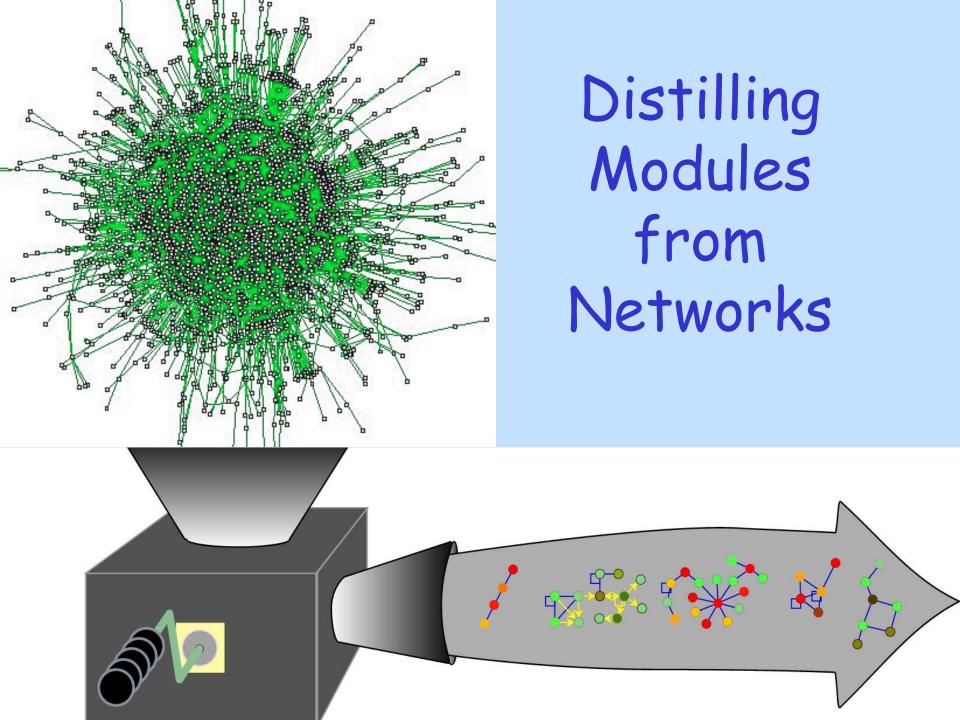
	network	type	n	m	z	$\ell$
	film actors	undirected	449 913	25516482	113.43	3.48
	company directors	undirected	7673	55392	14.44	4.60
	math coauthorship	undirected	253339	496489	3.92	7.57
	physics coauthorship	undirected	52 909	245300	9.27	6.19
social	biology coauthorship	undirected	1520251	11803064	15.53	4.92
SOC	telephone call graph	undirected	47 000 000	80 000 000	3.16	
	email messages	directed	59912	86 300	1.44	4.95
	email address books	directed	16881	57029	3.38	5.22
	student relationships	undirected	573	477	1.66	16.01
	sexual contacts	undirected	2810			
п	WWW nd.edu	directed	269 504	1497135	5.55	11.27
tio	WWW Altavista	directed	203 549 046	2130000000	10.46	16.18
information	citation network	directed	783 339	6716198	8.57	
nfoı	Roget's Thesaurus	directed	1022	5103	4.99	4.87
·:	word co-occurrence	${ m undirected}$	460902	17000000	70.13	

(Taken from Newman'03)

## Module Identification

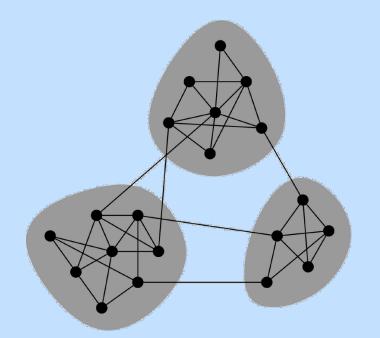
### Gene/Protein Modules

- A *module* is a set of genes/proteins performing a distinct biological function.
- Characterized by a coherent behavior of its genes w.r.t. a certain biological property.
- Examples:
  - transcriptional module: a set of co-expressed genes sharing a common function.
  - protein complex: assembly of proteins that build up some cellular machinery.
  - *signaling pathway*: a chain of interacting proteins propagating a signal in the cell.



# Modularity and Community Structure in Networks

M.E.J Newman, PNAS 2006



# Modularity of a division (Q)

Q = #(edges within groups) - E(#(edges within groups in a RANDOM graph with same node degrees))

Trivial division: all vertices in one group

==> Q(trivial division) = 0

k<sub>i</sub> = degree of node i

 $M = \sum k_i = 2|E|$ 

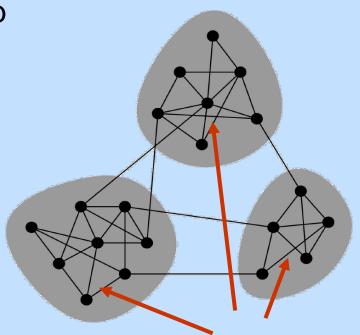
Aij = 1 if  $(i,j) \in E$ , 0 otherwise

Eij = expected number of edges

between i and j in a random graph with

same node degrees.

<u>Lemma</u>: Eij  $\approx k_i^* k_j / M$ 



Edges within groups

 $Q = \sum (Aij - ki*kj/M \mid i,j \text{ in the same group})$ 

## Division into two groups

## $Q = \sum (Aij - ki*kj/M \mid i,j \text{ in the same group})$

- Suppose we have n vertices {1,...,n}
- $s \{\pm 1\}$  vector of size n. Represent a 2-division:
  - si == sj iff i and j are in the same group
  - $-\frac{1}{2}$  (si\*sj+1) = 1 if si==sj, 0 otherwise

• ==> 
$$Q = \frac{1}{2} \sum_{i,j} (A_{ij} - \frac{k_i k_j}{M})(s_i s_j + 1)$$

## Division into two groups (2)

$$Q = \frac{1}{2} \sum_{i,j} (A_{ij} - \frac{k_i k_j}{M})(s_i s_j + 1)$$



Since 
$$\sum_{i,j} A_{ij} = \sum_i k_i = M$$

$$Q = \frac{1}{2} \sum_{i,j} (A_{ij} - \frac{k_i k_j}{M}) s_i s_j$$



**B** = the modularity matrix

- symmetric

$$Q = rac{1}{2} oldsymbol{s}^T oldsymbol{B} oldsymbol{s}^T$$

$$B = rac{1}{2} oldsymbol{s}^T oldsymbol{B} oldsymbol{s}^T$$
 where  $B_{ij} = A_{ij} - rac{k_i k_j}{M}$ 

## Division into two groups (3)

**B** is symmetric  $\Rightarrow$  **B** is diagonalizable (real eigenvalues)

**B**'s eigenvalues

B's orthonormal eigenvectors

$$\beta_1 \geq \beta_2 \geq \cdots \geq \beta_n \left[ \boldsymbol{u}_1, \boldsymbol{u}_2, \dots \boldsymbol{u}_n \right] \left[ \boldsymbol{B} \boldsymbol{u}_{\mathrm{i}} = \beta_{\mathrm{i}} \boldsymbol{u}_{\mathrm{i}} \right]$$

$$Q = \frac{1}{2} \mathbf{s}^T \mathbf{B} \mathbf{s}$$
 $Q = \sum_i a_i \mathbf{u}_i$ 
 $Q = \sum_i a_i \mathbf{u}_i$ 

- · Which vector s maximizes Q?
  - clearly  $\mathbf{s} \sim \mathbf{u1}$  maximizes Q, but  $\mathbf{u1}$  may not be  $\{\pm 1\}$  vector
  - Heuristic: maximize the projection of s on u1 ( $a_1$ ): choose si= +1 if u1<sub>i</sub>>0, si=-1 otherwise

# Identifying protein pathways

## Finding Simple Paths

Problem: Given a graph G=(V,E) and a parameter k, find a simple path of length k in G.

- NPC by reduction from Hamiltonian path.
- Trivial algorithm runs in  $O(n^k)$ .
- We will be interested in a *fixed parameter* algorithm (Downey & Fellows '92) i.e., time is exponential in k but polynomial in n.

## Color Coding [AYZ'95]

Problem: Given a graph G=(V,E) and a parameter k, find a simple path with k vertices (length k-1) in G.

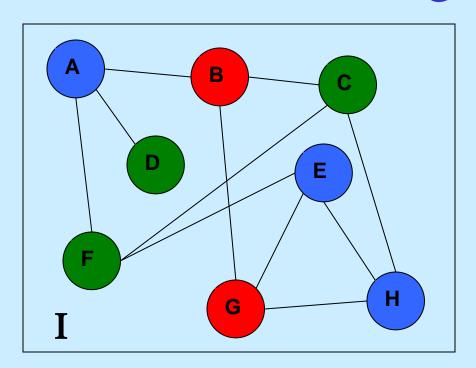
Algorithm: Randomly color vertices with *k* colors, and find a *colorful* path (distinct colors).

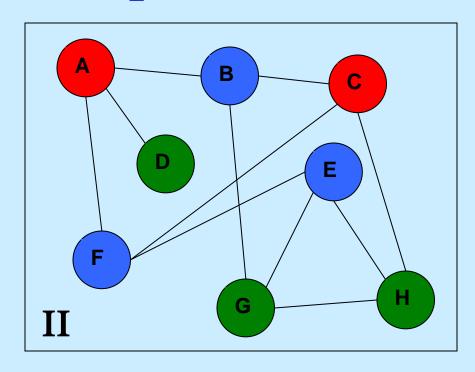
$$c: V \to [1, k]; S \in 2^{[1,k]}$$

$$P(v, S) = \max_{u:(u,v)\in E, c(u)\in S-\{c(v)\}} P(u, S - \{c(v)\})$$

Main idea: only  $2^k$  color subsets vs.  $n^k$  node subsets.

## Coloring Example





- Two different colorings on toy graph, k=3
- In coloring I, P(A,RGB) is built C->BC->ABC
- In coloring II, P(A,RGB) is built G->BG->ABG
- ABC is not colorful in coloring **II**

## Randomization Analysis

- A colorful path is simple, but a simple path may not be colorful *under a given coloring*
- Solution: run multiple independent trials.
- After one trial:

```
\Pr(Success) = \frac{k!}{k^k} \ge \frac{1}{e^k}
```

## Color Coding [AYZ'95]

#### **Complexity:**

- Space complexity is  $O(2^k n)$ .
- Colorful path found by DP in  $O(km2^k)$ .
- $-O(e^k)$  iterations are sufficient.
- Overall time is  $2^{O(k)}m$ .
- Note that the exponential part involves the parameter only, that is, the problem is *fixed* parameter tractable.

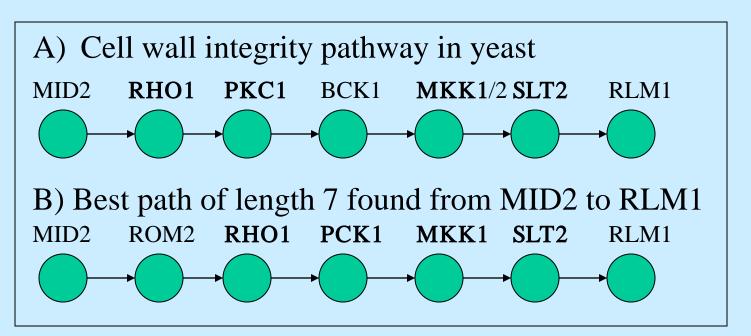
## Comparison of Running Times

Path length	Color coding	Exhaustive
8	435	866
9	2,149	15,120
10	11,650	

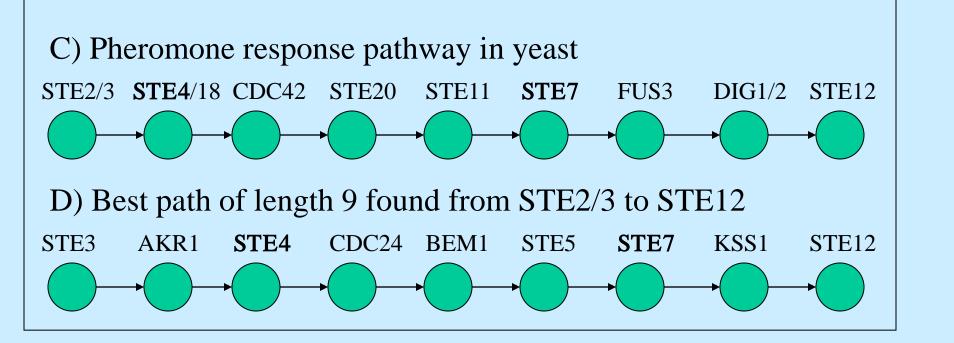
• ~4500 vertices, ~14500 edges.

## Biologically-Motivated Constraints

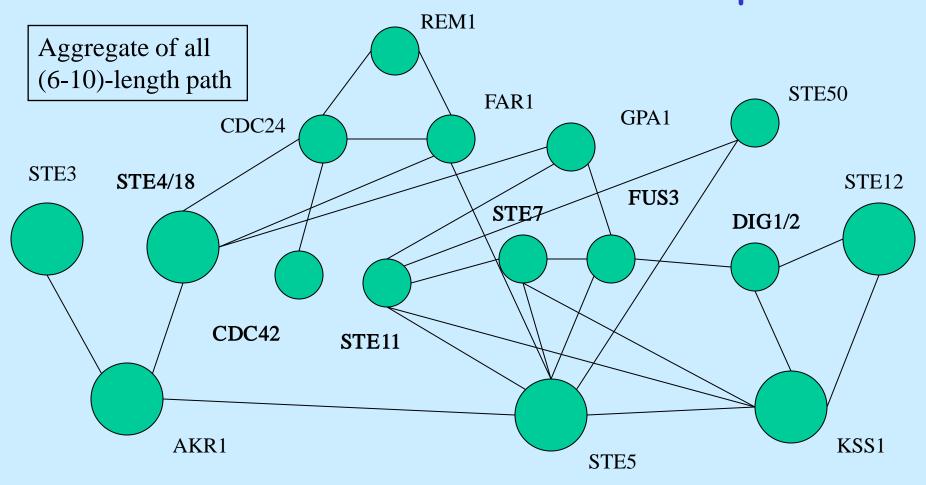
- Color-Coding gives an algorithmic basis, now introduce biologically motivated extensions.
- Can introduce edge weights (confidence).
- Can constrain the start or end of a path by type.
  - Steffen et al. '02: pathways from membrane to TF.
- Can force the inclusion of a specific protein on the path by ...



Appl. to yeast



## A Closer Look at Pheromone Response



The real pathway (main chain):

