

Termination

Well-Founded Orderings

$n := 0$

while $x > 0$ **do**

$n := n + 1$

$y := 0$; **while** $y^2 + 2y \leq x$ **do** $y := y + 1$

if $x = y^2$

then $x := y - 1$

else $s := 0$

$r := 0$; **while** $r^2 + 2r \leq x - y^2$ **do** $r := r + 1$

while $x > y^2 + r^2$ **do**

$y := 0$; **while** $y^2 + 2y \leq x$ **do** $y := y + 1$

$s := s + (s + y^2 + y - x)^2$

$x := x - y^2$

$r := 0$; **while** $r^2 + 2r \leq x - y^2$ **do** $r := r + 1$

for $i := 1$ **to** n **do** $x := r^2 + r - 1$

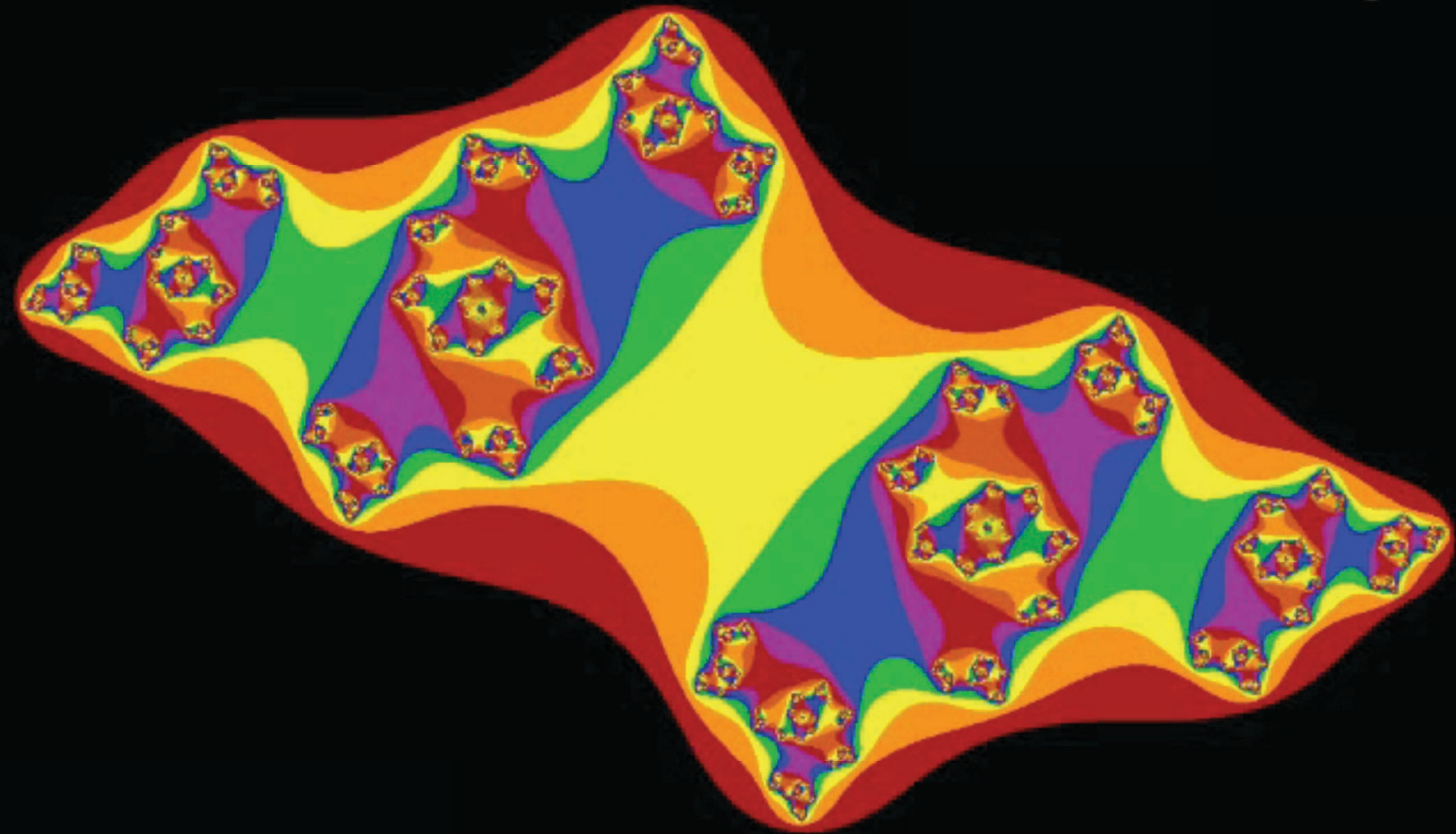
while $s > 0$ **do**

$r := 0$; **while** $r^2 + 2r \leq s$ **do** $r := r + 1$

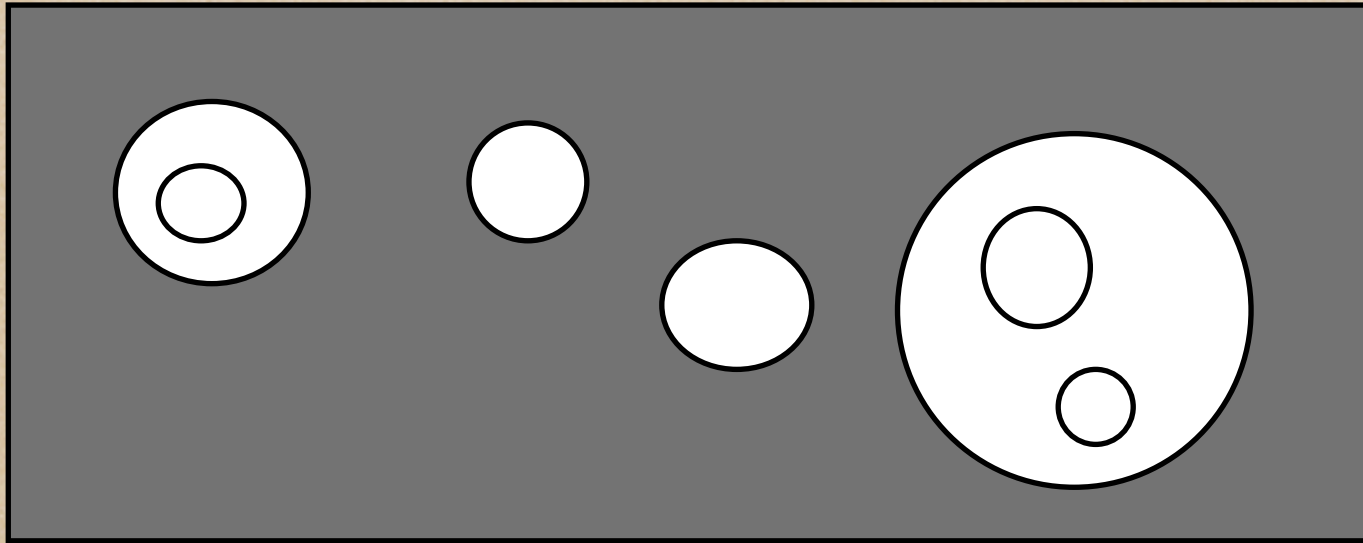
$x := x + (x + r^2 + r - s)^2$

$s := s - r^2$

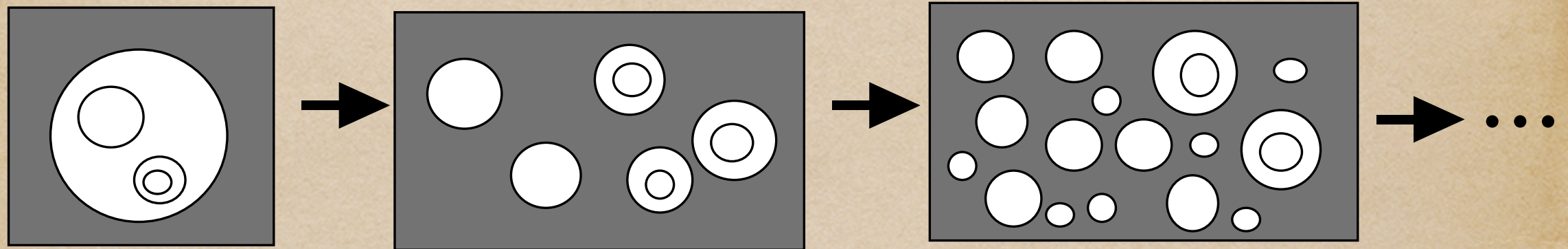


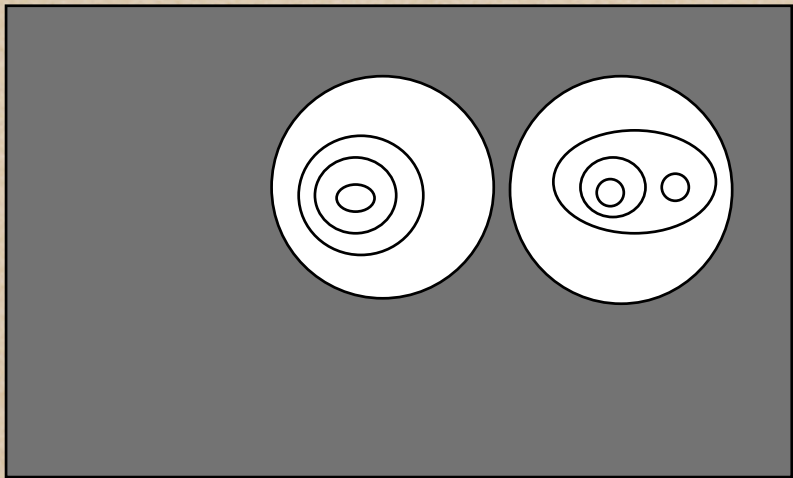


Amoebae

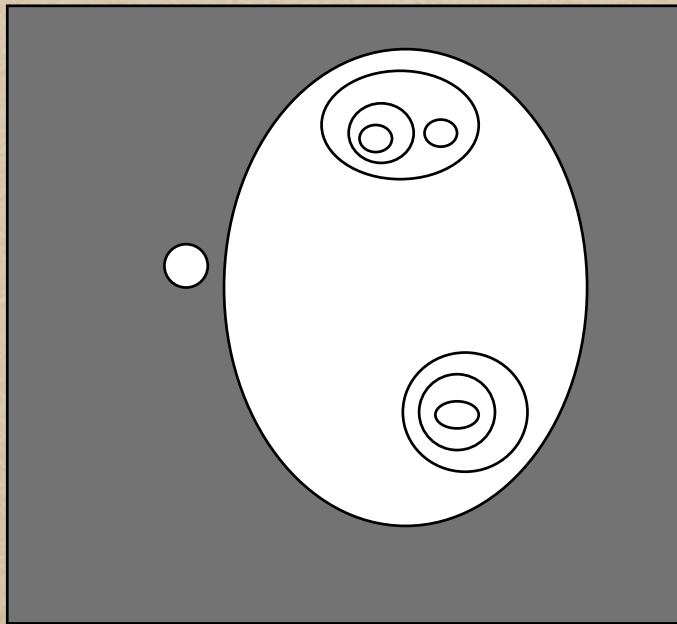
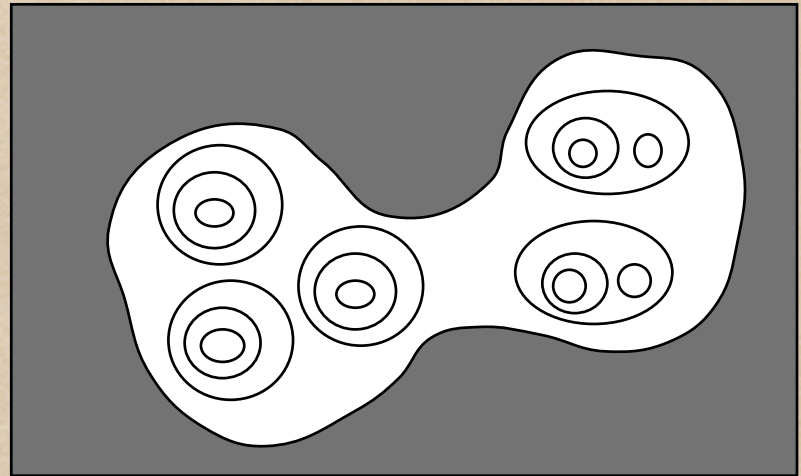


Fission

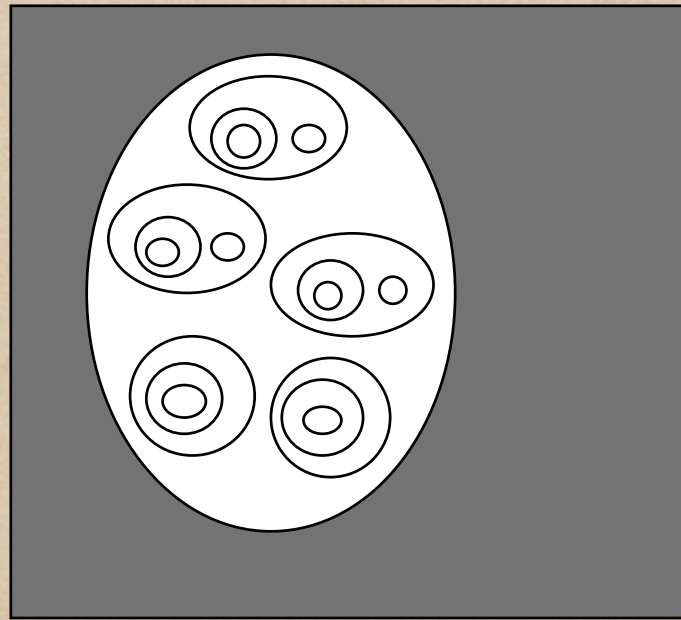




→
fusion



→
fusion



Colony Dies Out

- ◆ $\text{depth}(o) = 0$
- ◆ $\text{depth}(a_1 \dots a_n) = 1 + \max\{\text{depth}(a_i)\}$
- ◆ $\{ (\text{depth}(a), |a|) : \text{subcolony } a \}$
- ◆ fission: depth decreases
- ◆ fusion: size decreases

Big Picture

- ◆ Programs are state-transition systems
- ◆ Choose a well-founded order on states
- ◆ Show that transitions are decreases

Real Picture

- ◆ Programs are state-transition systems
- ◆ Choose a function for “ranking” states
- ◆ Choose a well-founded order on ranks
- ◆ Show that transitions always decrease rank

Imaginary Picture

- ◆ Programs are state-transition systems
- ◆ Choose a function for “ranking” states
- ◆ Choose a well-founded order on ranks
- ◆ Show that transitions eventually decrease rank

Nested Loops

$r := 1$

$u := 1$

loop

$v := u$

until $r \geq n$

$s := 1$

loop $u := u + v$

$s := s + 1$

while $s \leq r$

repeat

$r := r + 1$

repeat

$$(n - r)\omega^2 + (r - s)\omega + k$$

Per Iteration

$r := 1$

$u := 1$

loop

$v := u$

until $r \geq n$

$s := 1$

loop $u := u + v$

$s := s + 1$

while $s \leq r$

repeat

$r := r + 1$

repeat

$\omega(n-r) + r + 1 - s$

Lexicographic

$r := 1$

$u := 1$

loop

$v := u$

until $r \geq n$

$(n-r, r+1-s)$

$s := 1$

loop $u := u+v$

$s := s+1$

while $s \leq r$

repeat

$r := r+1$

repeat

Invariants

$r := 1$

$u := 1$

loop

$v := u$

until $r \geq n$

$s := 1$

loop $u := u + v$

$s := s + 1$

while $s \leq r$

repeat

$r := r + 1$

repeat

$1 \leq r \leq n$

$1 \leq s \leq r + 1$

Well-Founded Orderings

- ◆ No infinite descending sequences

$$x_1 > x_2 > x_3 > \dots$$

Well-Founded Induction

$>$ is a wfo of X

$$\frac{\forall x \in X. [\forall y < x. P(y)] \Rightarrow P(x)}{\forall x \in X. P(x)}$$

Why?

David Gries

- ◆ Under the reasonable assumption that nondeterminism is bounded, the two methods are equivalent.... In this situation, we prefer using strong termination.

All-Purpose Ranks

$$0 < 1 < 2 < \dots$$

$$< \omega < \omega+1 < \omega+2 < \dots$$

$$< \omega^2 < \omega^2+1 < \dots < \omega^3 < \dots < \omega^4 < \dots$$

$$< \omega^2 < \omega^2+1 < \dots < \omega^2+\omega < \omega^2+\omega+1 < \dots$$

$$< \omega^3 < \omega^3+1 < \dots < \omega^4 < \dots < \omega^5 < \dots$$

$$< \omega^\omega < \dots < \omega^{\omega^\omega} < \dots < \omega^{\omega^{\omega^\omega}} < \dots$$

Ordinals

$0, 1, 2, \dots,$

$\omega, \omega+1, \omega+2, \dots,$

$2\omega, 2\omega+1, \dots, 3\omega, \dots,$

$\omega^2, \dots, \omega^2+2\omega+3, \dots, \omega^3, \dots,$

$\omega^\omega, \dots, \omega^{\omega^\omega}, \dots,$

$\varepsilon_0, \varepsilon_0+1, \dots, 2\varepsilon_0+\omega^\omega+2\omega+3, \dots,$

$\varepsilon_1, \dots, \varepsilon_{\varepsilon_0}, \dots,$

\dots

THE GROWTH

$0, 1, 2, 3, \dots, \omega, \omega+1, \omega+2, \omega+\omega, \dots, \omega+\omega = \omega^2, \omega^2+1, \omega^2+2, \omega^2+3, \dots, \omega^3, \omega^3+1, \omega^3+2, \omega^3+3, \dots$
 $\omega\omega = \omega^2, \omega^2+1, \omega^2+2, \dots, \omega^2+\omega, \dots, \omega^2+\omega^2, \dots, \omega^2+\omega^3, \dots, \omega^2_2, \dots, \omega^2_3, \dots$
 $\dots, \omega^{\omega^2}, \dots, \omega^{\omega^3}, \dots, \omega^{\omega^{\omega^2}}, \dots, \omega^{\omega^{\omega^3}}, \dots, \omega^{\omega^{\omega^{\omega^2}}}, \dots, \omega^{\omega^{\omega^{\omega^3}}}, \dots, \omega^{\omega^{\omega^{\omega^{\omega^2}}}}, \dots, \omega^{\omega^{\omega^{\omega^{\omega^3}}}}, \dots$
 $\epsilon_0 + \omega^2, \dots, \epsilon_0 + \omega^{\omega^2}, \dots, \epsilon_0 + \omega^{\omega}, \dots, \epsilon_0 + \epsilon_0 = \epsilon_0^2, \dots, \epsilon_0^3, \dots, \epsilon_0^{\omega^2}, \dots, \epsilon_0^{\omega^2}, \dots$
 $\epsilon_0 \epsilon_0 = \epsilon_0^{\omega}, \dots, \epsilon_0^3, \dots, \epsilon_0^{\omega}, \dots, \epsilon_0^{\omega^{\omega}}, \dots, \epsilon_0^{\epsilon_0}, \dots, \epsilon_0^{\epsilon_0^{\omega}} = \epsilon_1, \dots, \epsilon_2, \dots, \epsilon_3, \dots, \epsilon_{\omega}$
 $\epsilon_{\omega^{\omega}}, \dots, \epsilon_{\epsilon_0}, \dots, \epsilon_{\epsilon_1}, \dots, \epsilon_{\epsilon_{\omega}}, \dots, \epsilon_{\epsilon_{\omega^{\omega}}}, \dots, \epsilon_{\epsilon_{\epsilon_0}}, \dots, \epsilon_{\epsilon_{\epsilon_1}} = \eta_0, \dots, \eta_1, \dots$
 $\eta_{\epsilon_0}, \dots, \eta_{\epsilon_{\omega}}, \dots, \eta_{\epsilon_{\epsilon_0}}, \dots, \eta_{\epsilon_{\epsilon_{\epsilon_0}}} = \eta_{\eta_0}, \dots, \eta_{\eta_{\eta_0}} = \zeta_0, \dots, \dots$

Transition System

State

Transition

Well-Founded Method

- ◆ States Q
- ◆ Algorithm $R \subseteq Q \times Q$
- ◆ Well-founded order $>$ on Q
- ◆ $R \subseteq >$

All-Purpose Ranking

◆ $r : Q \rightarrow \text{Ord}$

◆ $r(x) = \sup \{ r(y) + 1 : x \rightarrow y \}$

Computation



Abstraction



Frank Ramsey



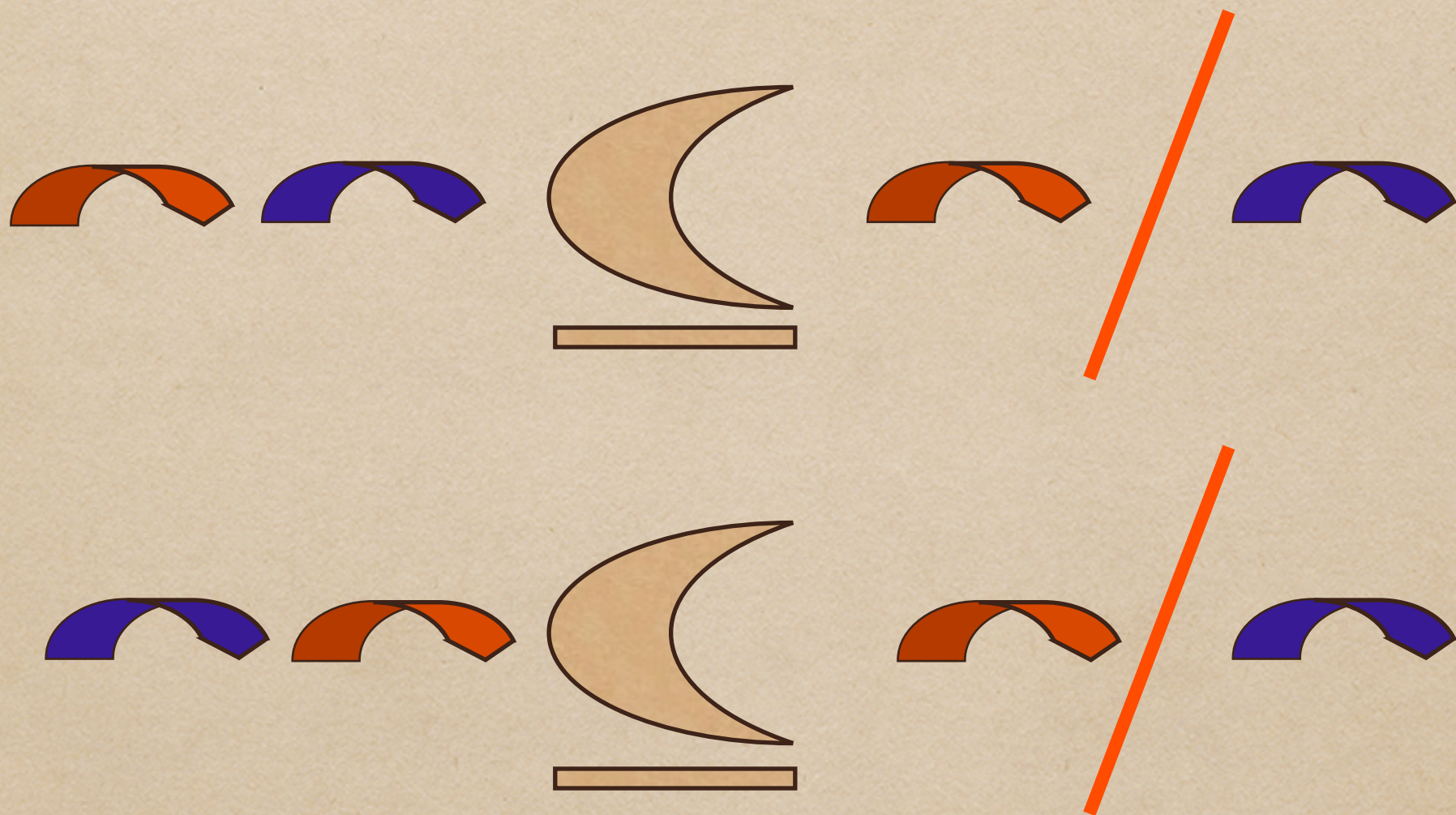
Ramsey's Theorem

Infinite complete graph

Finitely colored edges

Monochrome infinite clique

Closure



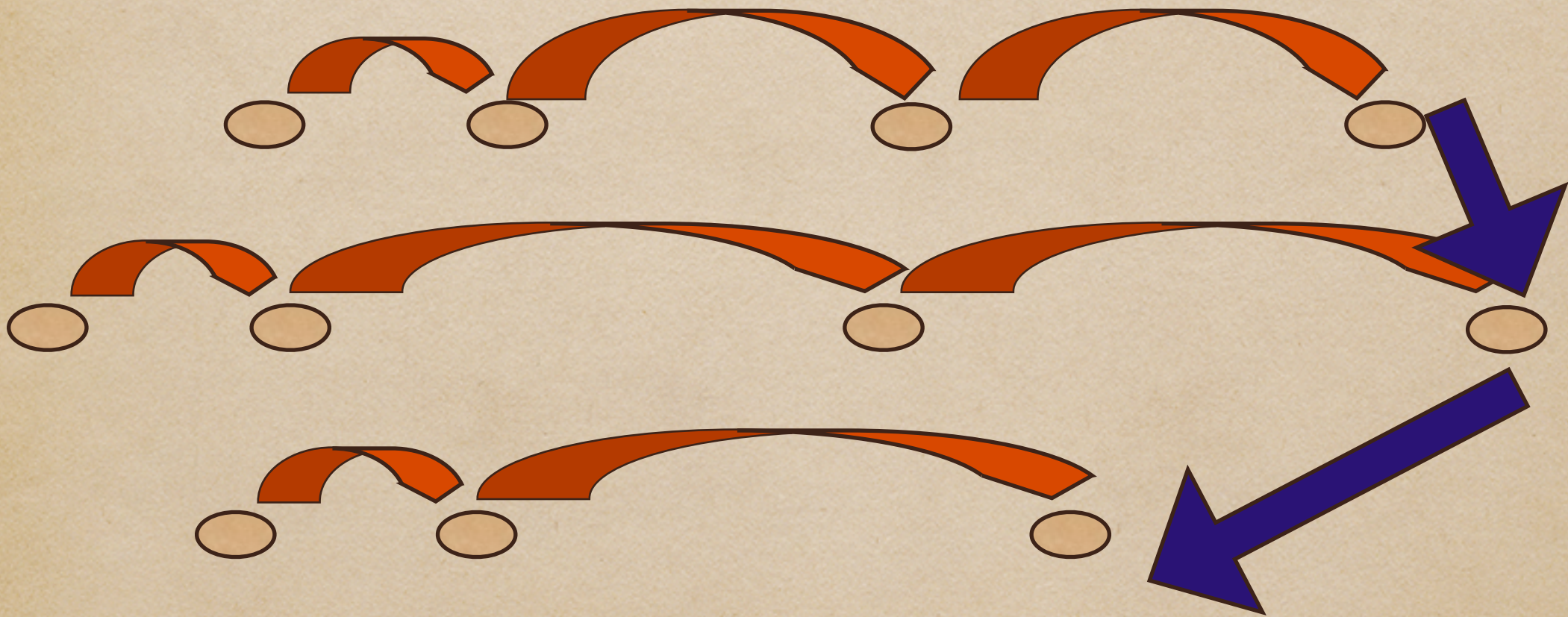
Proof



Proof



Proof



Disjunctive Orders

- ◆ States Q
- ◆ Algorithm $R \subseteq Q \times Q$
- ◆ Transitive closure R^+
- ◆ Well-founded orders $>$ and \sqsupset on Q
- ◆ $R^+ \subseteq > \cup \sqsupset$

Ranking Method

- ◆ States Q
- ◆ Algorithm $R \subseteq Q \times Q$
- ◆ Well-founded order \succ on W
- ◆ Ranking function $r : Q \rightarrow W$
- ◆ Define $X > Y$ if $r(X) \succ r(Y)$
- ◆ $R \subseteq >$

Invariants

- ◆ States Q
- ◆ Algorithm $R \subseteq Q \times Q$
- ◆ Well-founded order \succ on W
- ◆ Ranking function $r : Q \rightarrow W$
- ◆ Define $X > Y$ if $r(X) \succ r(Y)$
- ◆ $R \subseteq >$

Algorithmic System

State

P
r
o
g
r
a
m

Transition

Classical Algorithms

- ◆ Every algorithm can be expressed precisely as a set of conditional assignments, executed in parallel repeatedly.

if c then $f(s_1, \dots, s_n) := t$

if c then $f(s_1, \dots, s_n) := t$

if c then $f(s_1, \dots, s_n) := t$

Practical Method

- ◆ States Q
- ◆ Algorithm $R \subseteq Q \times Q$
- ◆ Well-founded order \succ on W
- ◆ Ranking function $r : Q \rightarrow W$
- ◆ Define $X > Y$ if $r(X) \succ r(Y)$
- ◆ $R \subseteq >$

The image shows a cardboard cutout of a circular logo. The logo consists of a teal-colored ring with a white horizontal bar across its center. The letters 'DLR' are printed in white on the blue bar. The cutout is placed on a light brown, textured surface.

DLR



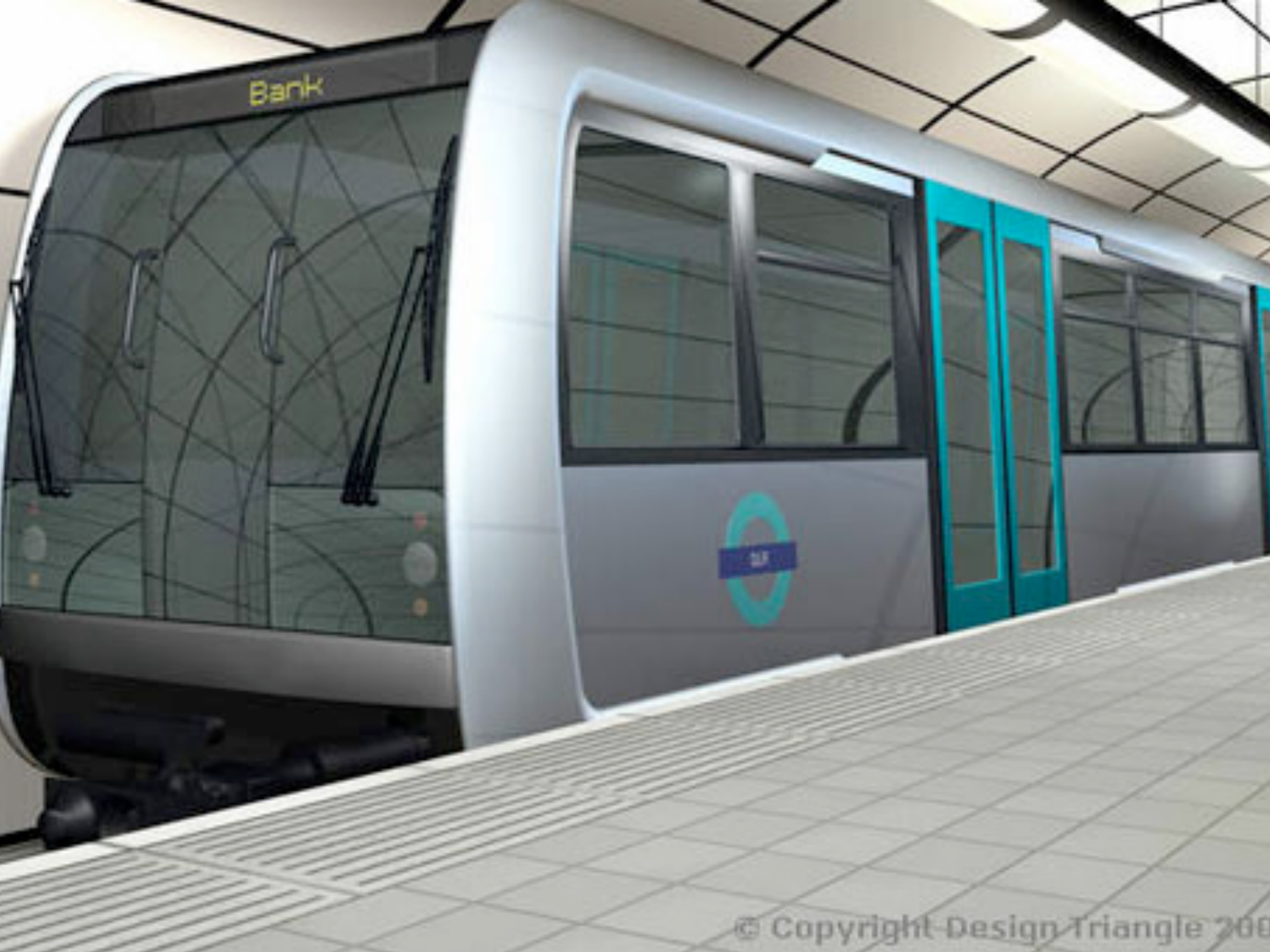
Town Centre
Woolwich New Rd
National Rail

Board

www.dlr.co.uk



95



Bank



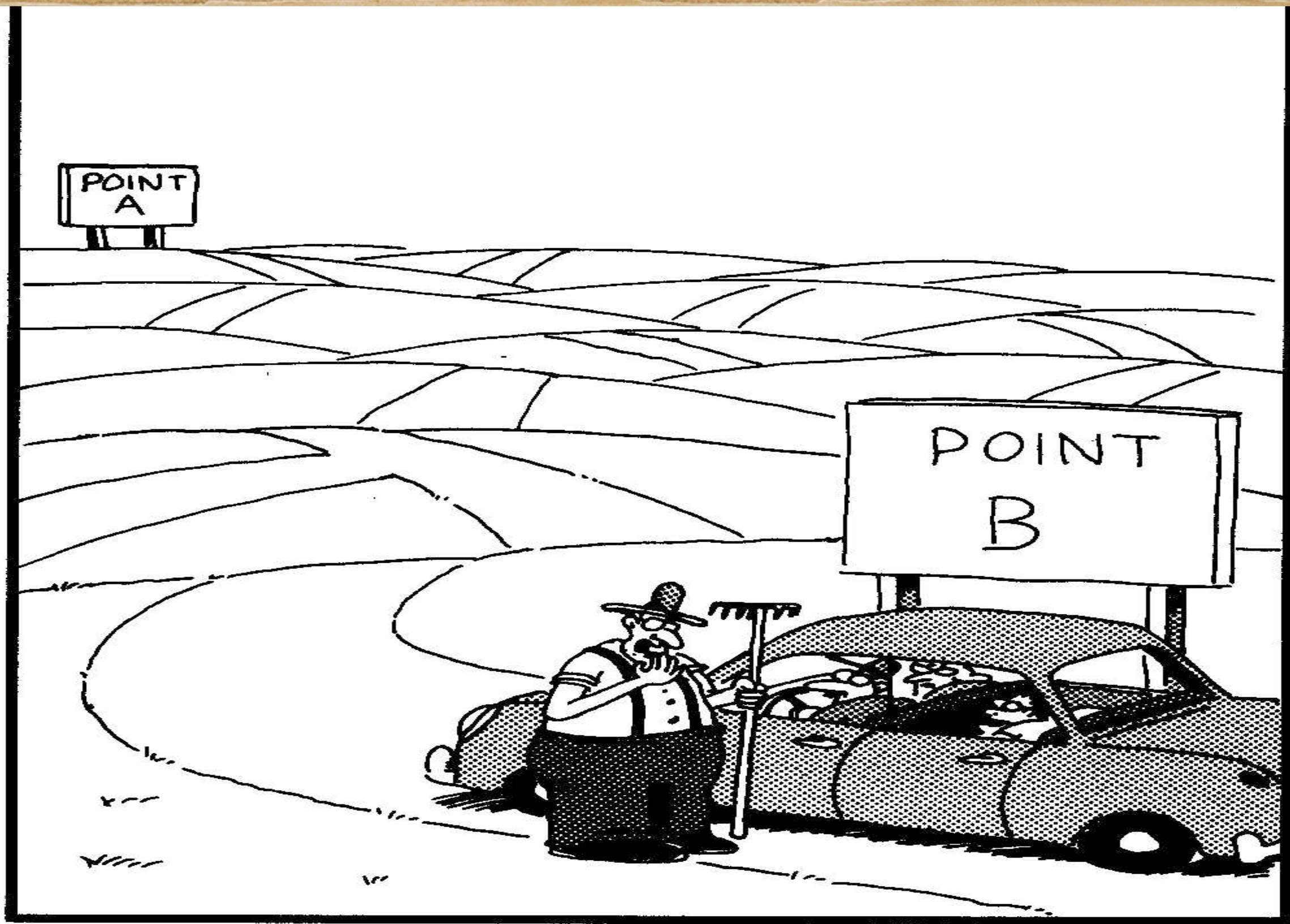
Color Code

Bordeaux



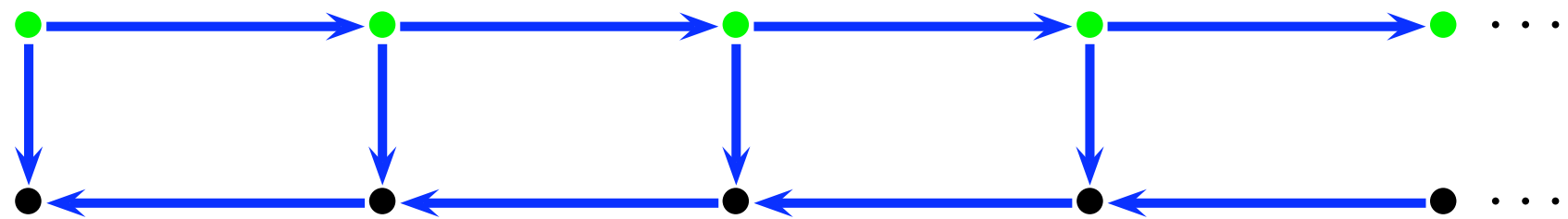
Azure



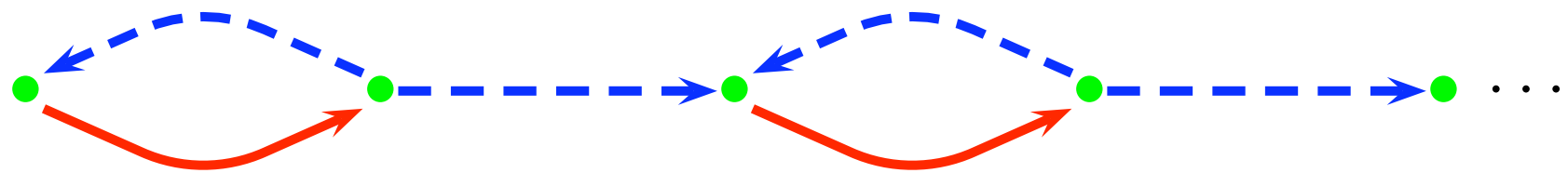


“Well, lemme think. ... You’ve stumped me, son. Most folks only wanna know how to go the other way.”

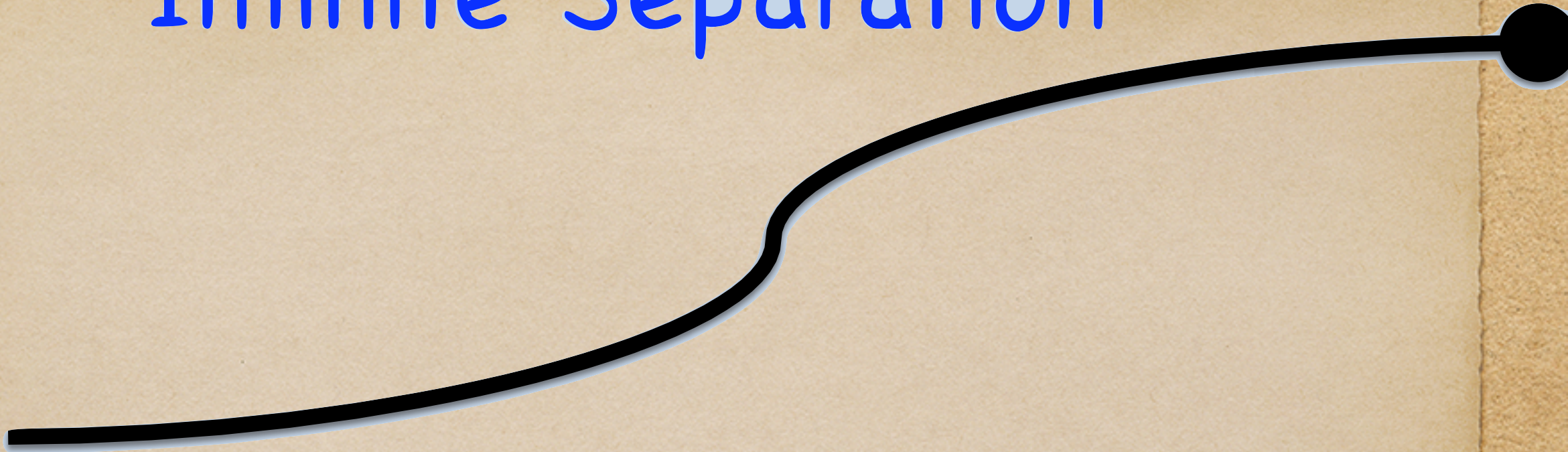
Mortal (black) nodes on bottom and immortal (green) nodes on top



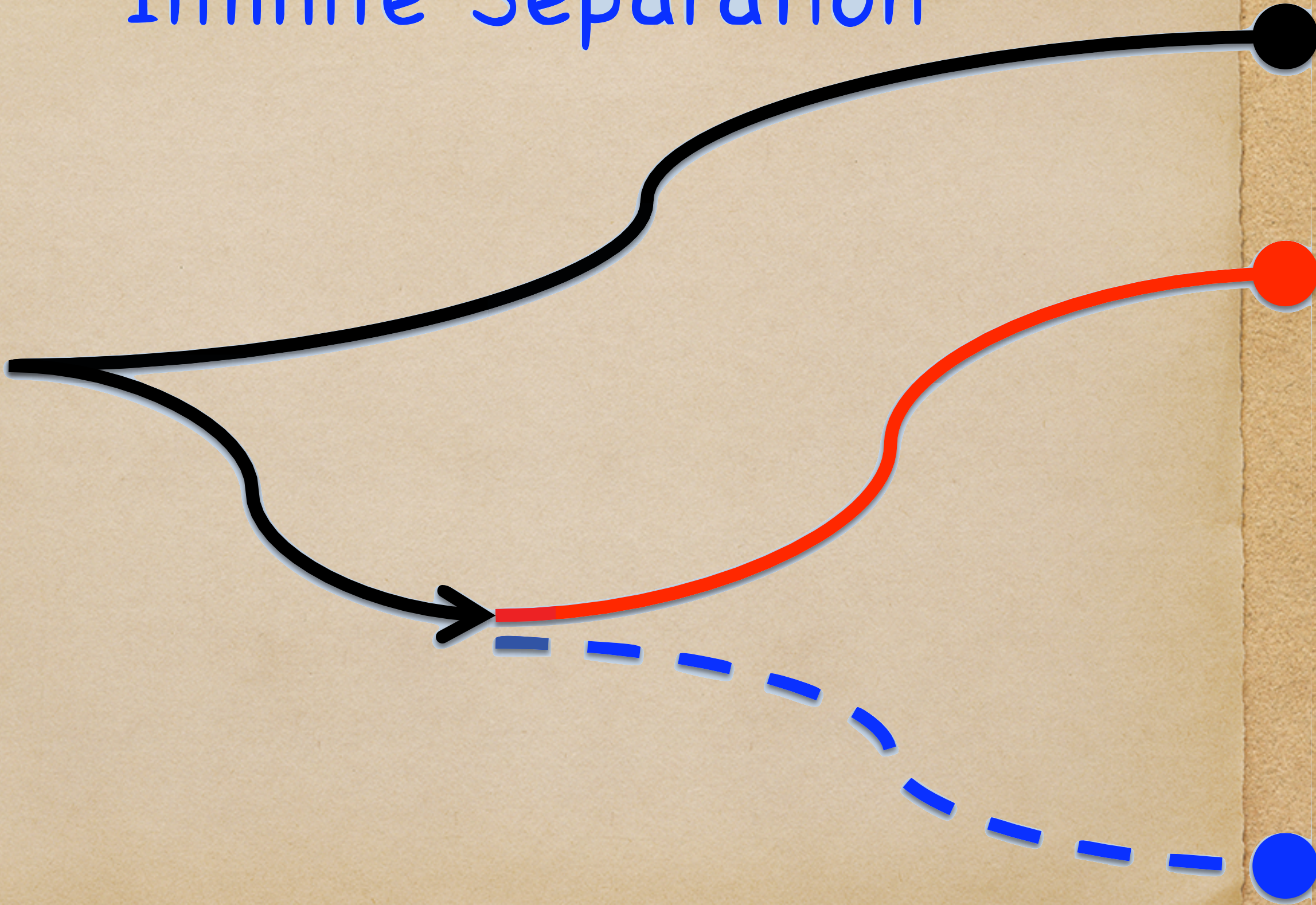
Mortal in each alone (dashed **Azure** or solid **Bordeaux**),
but immortal in their union



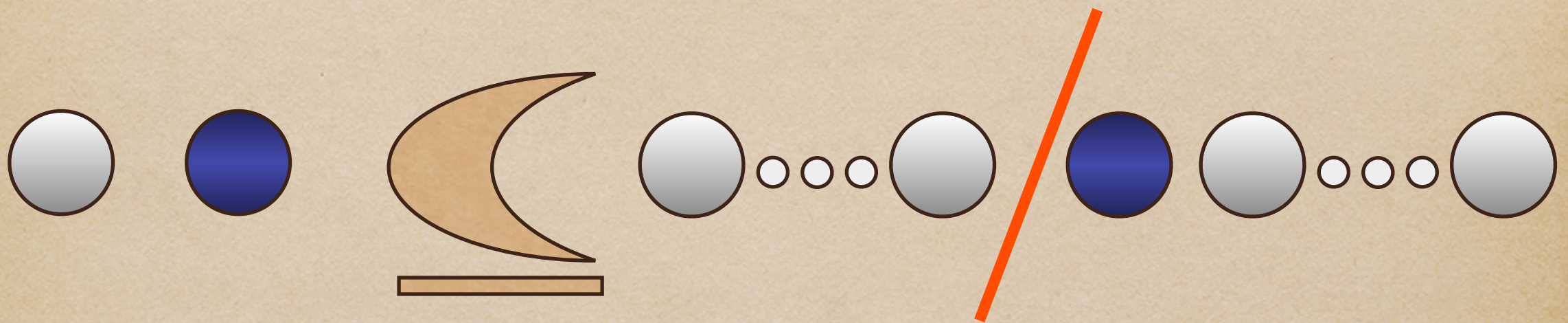
Infnite Separation



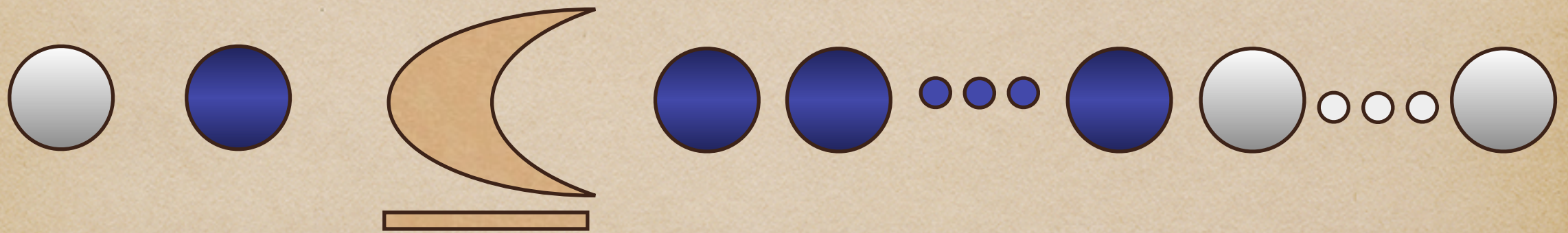
Infinite Separation



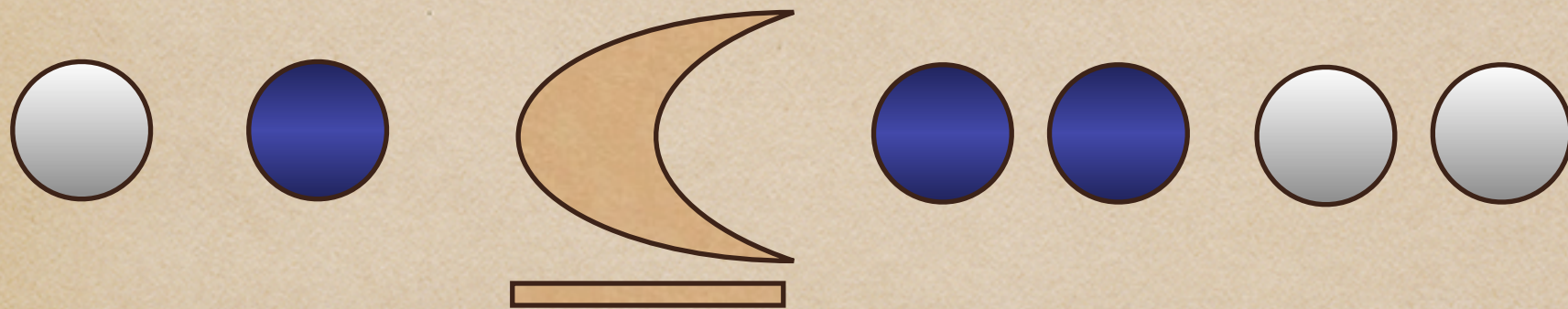
Enough?



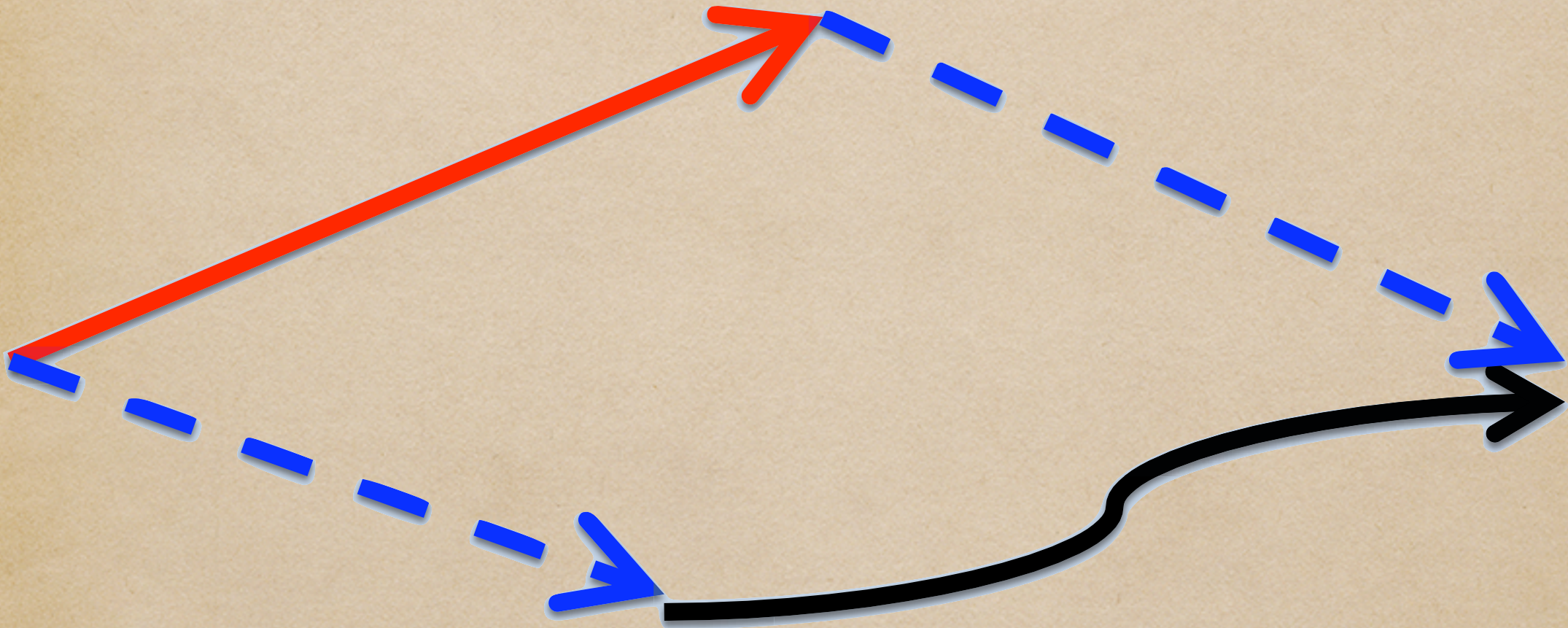
Enough?



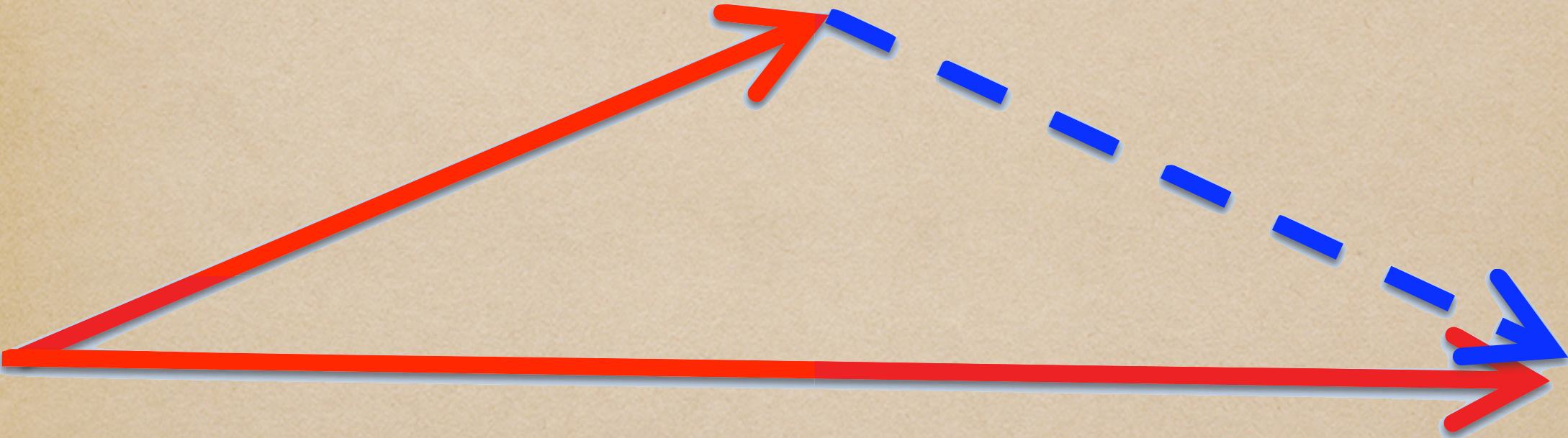
Enough?



Jumping

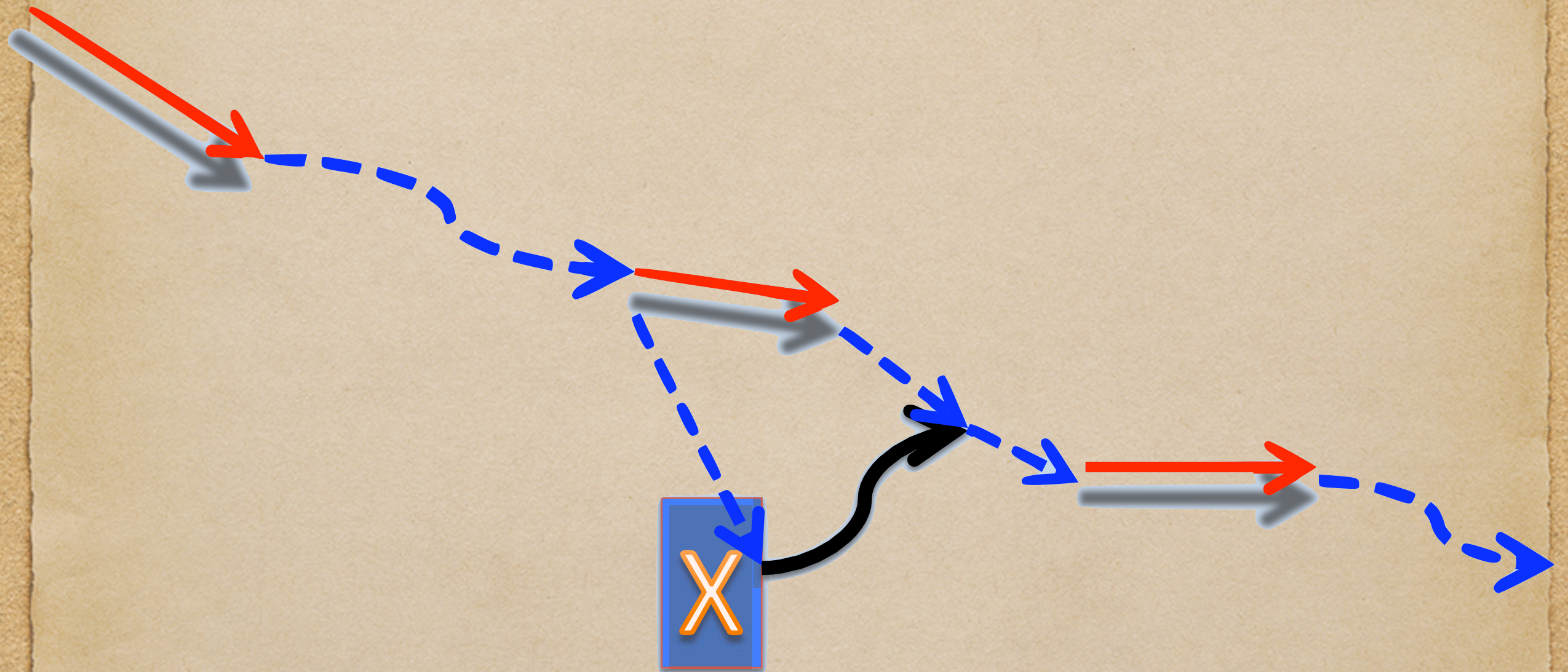


Jumping

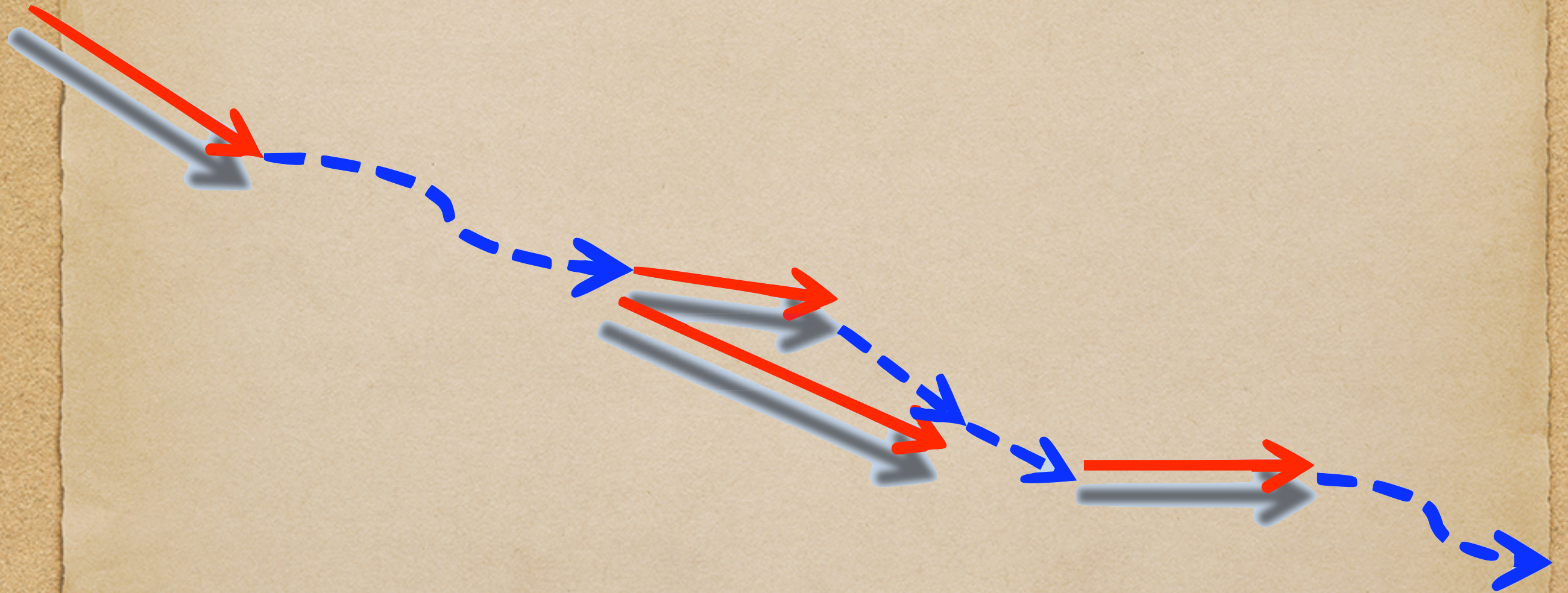




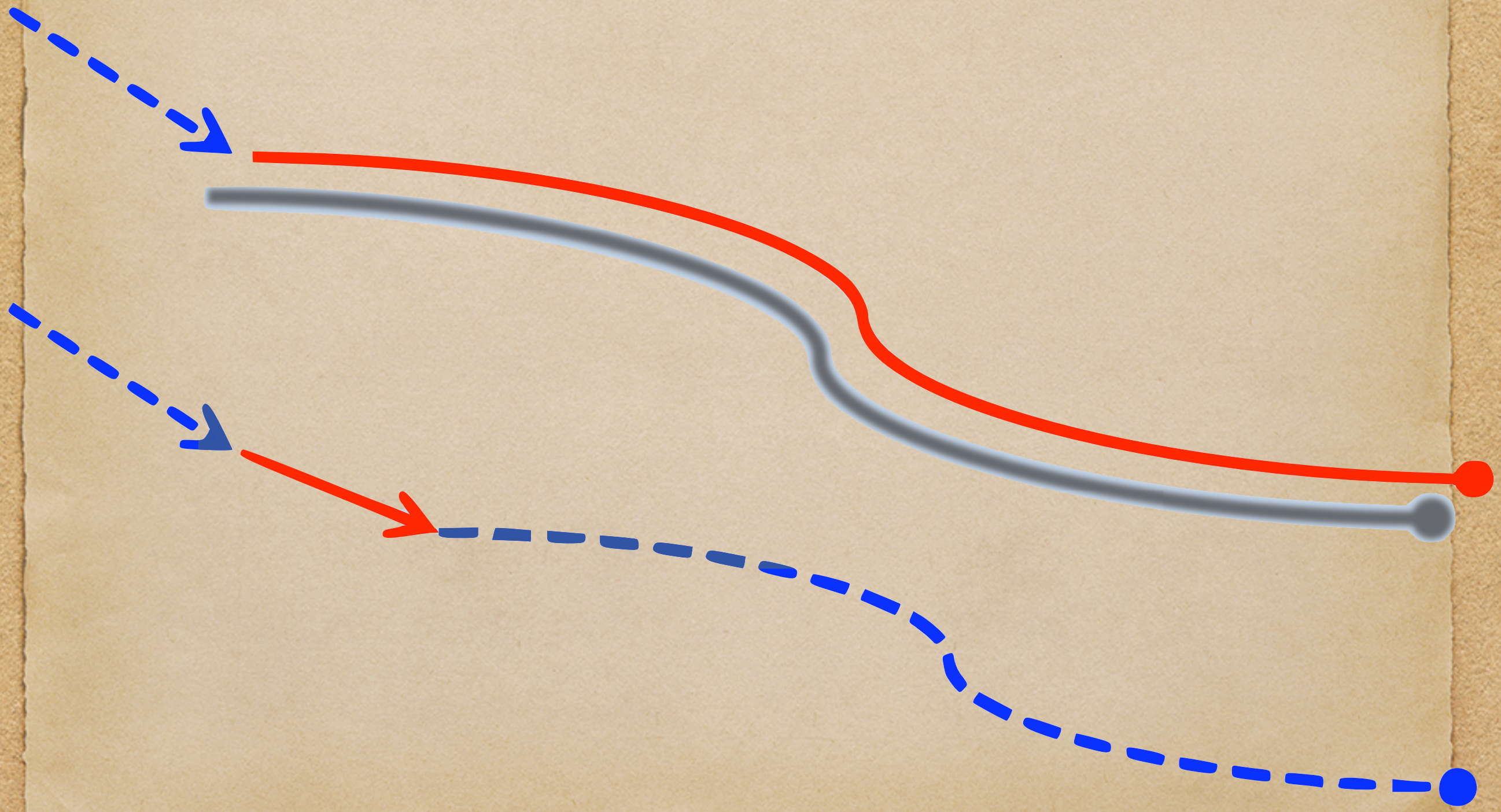
Constriction + Jumping



Constriction + Jumping



Constriction + Jumping



Gremlins



Gremlins

